

## Mood as illocutionary centering

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### Abstract

By this point, we have developed articulated analyses of top-level temporal anaphora, including temporal quantification, in languages with grammatical tense and/or aspect systems, represented by English, Polish, and Mandarin. But it is still not clear how this approach might extend to temporal anaphora in a language such as Kalaallisut, which has neither grammatical tense nor grammatical aspect, but instead marks only grammatical mood and person. Most theories of mood and modal reference either ignore temporal reference entirely (e.g. Hamblin 1973, Stalnaker 1975, Karttunen 1977, Kratzer 1981, Groenendijk and Stokhof 1984, Roberts 1987, Farkas 1992, Portner 2004, a.m.o.) or analyze modal and temporal reference as independent phenomena (e.g. Muskens 1995, Stone 1997, Mastop 2005). Such theories provide no inkling how temporal reference in a tenseless mood-based language can be as precise as in English (see e.g. Bittner 2005, 2007, 2009, 2011).

More promising is the performative approach, which relates modal and temporal reference in a way that might shed light on this puzzle (Stenius 1967, etc). For example, Lewis (1972) proposes “to analyze all sentences, declarative and non-declarative, into two components: a *sentence radical* that specifies a state of affairs, and a *mood* that determines whether the speaker is declaring that the state of affairs holds, commanding that it hold, asking whether it holds, or what.”

Lewis’s own implementation is problematic. For example, it wrongly predicts that all sentences, including non-declaratives, have truth values. It also fails to extend to embedded moods. In Kalaallisut, however, this failure is, in fact, a plus, because illocutionary moods do not embed. Moreover, recent theories of *at-issue* and *not-at-issue* assertions (see AnderBois *et al.* 2010, Murray 2010) make it possible to distinguish declaratives, which make at-issue assertions and therefore have truth values, from non-declaratives, which do not. Building on these recent advances, I propose a new performative theory of illocutionary mood that avoids the problems of Lewis’s (1972) implementation and extends the proposed *TAMP* universals to grammatical mood.

### Outline

1. Declaratives with(out) reportative recentering
2. Optatives with(out) reportative recentering
3. Kalaallisut verbal inflections as top-level anaphora
4. Why only declaratives have truth conditions

## 1 DECLARATIVES WITH(OUT) REPORTATIVE RECENTERING

- (1) *Ole ani-sima-pu-q* (#*uanga=li tamanna upiri-nngit-la-ra*).  
 Ole go.out-prf-DEC<sub>T</sub>-3SG (1SG=but that believe-not-DEC-1SG.3SG)  
 Ole has gone out (#, but I don't believe it).
- (2) *Ole=guuq ani-sima-pu-q* (*uanga=li tamanna upiri-nngit-la-ra*).  
 Ole=*RPT* go.out-prf-DEC<sub>T</sub>-3SG (1SG=but that believe-not-DEC-1SG.3SG)  
 Ole has gone out, I'm told (, but I don't believe it).
- (3)  ${}^{st}e_a := \{\langle \langle w_a, p_a, e_a \rangle, \langle \rangle \rangle \mid w_a \in p_a \subseteq \{v \mid \langle e_a, \uparrow e_a, \downarrow e_a \rangle \in \llbracket \text{spk.to} \rrbracket(v)\}\}$
- (4)  $\langle \langle w_0, p_{012}, e_a \rangle, \langle \rangle \rangle$  Notation: *minimal info-state*  
 $\langle \langle w_1, p_{012}, e_a \rangle, \langle \rangle \rangle$   $p_{012} = \{w_0, w_1, w_2\}$   
 $\langle \langle w_2, p_{012}, e_a \rangle, \langle \rangle \rangle$   $p_{nm} = \{w_n, w_m\}$   
 $\langle \langle w_0, p_{01}, e_a \rangle, \langle \rangle \rangle$   $p_n = \{w_n\}$   
 $\langle \langle w_1, p_{01}, e_a \rangle, \langle \rangle \rangle$   
 $\langle \langle w_0, p_{02}, e_a \rangle, \langle \rangle \rangle$   
 $\langle \langle w_2, p_{02}, e_a \rangle, \langle \rangle \rangle$   
 $\langle \langle w_1, p_{12}, e_a \rangle, \langle \rangle \rangle$   
 $\langle \langle w_2, p_{12}, e_a \rangle, \langle \rangle \rangle$   
 $\langle \langle w_0, p_0, e_a \rangle, \langle \rangle \rangle$   
 $\langle \langle w_1, p_1, e_a \rangle, \langle \rangle \rangle$   
 $\langle \langle w_2, p_2, e_a \rangle, \langle \rangle \rangle$
- (1') Ole  ${}_{[DEC}$  go.out-prf-  $]-DEC_T-3SG$   
 ${}^T[x \mid x =_i ole]; (\partial[\text{spk.to}_{T\Omega} \langle \uparrow \varepsilon, \uparrow \varepsilon, \downarrow \varepsilon \rangle]; [go.out_{T\omega} \langle \uparrow \delta \rangle]; {}^T[p \mid p =_I \uparrow \omega]_{T\Omega})$
- (1'')  $\{\langle \langle p'_a, \mathbf{o}, w_a, p_a, e_a \rangle, \langle \rangle \rangle \mid \mathbf{o} = \llbracket ole \rrbracket$   
 $\& p_a \subseteq \{v \mid \langle e_a, \uparrow e_a, \downarrow e_a \rangle \in \llbracket \text{spk.to} \rrbracket(v)\}$   
 $\& w_a \in p'_a = \{v \in p_a \mid \mathbf{o} \in \llbracket \text{gone.out} \rrbracket(v)\}\}$
- ( $\overset{\leftarrow}{\ })$   $e' = \overset{\leftarrow}{e} \Rightarrow e, e' \in \mathcal{D}_\varepsilon \& \uparrow e' \neq \uparrow e$  ( $\overset{\leftarrow}{e}$ , source of  $e$ )
- (2') Ole=*RPT*  
 ${}^T[x \mid x =_i ole]; \partial[\text{spk.to}_{T\Omega} \langle \uparrow \varepsilon, \uparrow \varepsilon, \downarrow \varepsilon \rangle]; {}^T[p \mid \text{spk.to}_p \langle \overset{\leftarrow}{\uparrow} \varepsilon, \uparrow \overset{\leftarrow}{\uparrow} \varepsilon, \downarrow \overset{\leftarrow}{\uparrow} \varepsilon \rangle];$   
 ${}^T[w \mid w \in_i \uparrow \Omega]{}^T; \dots$   
 ${}_{[DEC}$  go.out-prf-  $]-DEC_T-3SG$   
 $\partial[\text{spk.to}_{T\Omega} \langle \overset{\leftarrow}{\uparrow} \varepsilon, \uparrow \overset{\leftarrow}{\uparrow} \varepsilon, \downarrow \overset{\leftarrow}{\uparrow} \varepsilon \rangle]; [gone.out_{T\omega} \langle \uparrow \delta \rangle]; {}^T[p \mid p =_I \uparrow \omega]_{T\Omega}$
- (2'')  $\{\langle \langle p'_{<b}, w_{<b}, p_{<b}, \mathbf{o}, w_b, p_b, e_b \rangle, \langle \rangle \rangle \mid \mathbf{o} = \llbracket ole \rrbracket$   
 $\& w_b \in p_b \subseteq \{v \in p_{<b} \mid \langle e_b, \uparrow e_b, \downarrow e_b \rangle \in \llbracket \text{spk.to} \rrbracket(v)\}$   
 $\& p_{<b} \subseteq \{v \mid \langle \overset{\leftarrow}{e}_b, \uparrow \overset{\leftarrow}{e}_b, \downarrow \overset{\leftarrow}{e}_b \rangle \in \llbracket \text{spk.to} \rrbracket(v)\}$   
 $\& w_{<b} \in p'_{<b} = \{v \in p_{<b} \mid \mathbf{o} \in \llbracket \text{gone.out} \rrbracket(v)\}\}$

## 2 OPTATIVES WITH(OUT) REPORTATIVE RECENTERING

Figure 1. Attitudes and ideals (cf. Lewis 1981)

$A_w e := \lambda p(A_w(e, \uparrow e, p))$   $A \in \{\underline{des}, \underline{dir}, \dots\} \subseteq Con_{\omega \varepsilon \delta \Omega t}$

$w <_Q v := \lambda p(p \in Q \wedge w \in p) \subset \lambda p(p \in Q \wedge v \in p)$

$OPT(p, Q) := \lambda w(w \in p \wedge \neg \exists v(v \in p \wedge w <_Q v))$

(5) *Ole pulaarli!*

*Ole pulaar-li*

Ole visit<sub>T-OPT.3SG</sub>

Let Ole come and visit!

(5') Ole<sup>T</sup>...

<sup>T</sup>[ $x | x =_i ole$ ]<sup>T</sup>; ...

visit<sub>T-OPT-3SG</sub>

$\partial[spk.to_{T\Omega} \langle T\varepsilon, \uparrow T\varepsilon, \downarrow T\varepsilon \rangle]; [w | visit_w \langle T\delta, \uparrow T\varepsilon \rangle]; [p | p =_I \perp \omega];$

$[OPT \langle T\Omega, \underline{des}_{T\omega} T\varepsilon \rangle \subseteq_i \perp \Omega]; [T\Omega =_I T\omega |_{T\Omega}]$

(5'')  $\{ \langle \mathbf{o}, w_a, p_a, e_a \rangle, \langle q_a, v_a \rangle \mid \mathbf{o} = \llbracket ole \rrbracket$

$\& v_a \in q_a = \{v \mid \langle \mathbf{o}, \uparrow e_a \rangle \in \llbracket visit \rrbracket(v)\}$

$\& w_a \in p_a \subseteq \{v \mid \langle e_a, \uparrow e_a, \downarrow e_a \rangle \in \llbracket spk.to \rrbracket(v) \&$

$\exists Q_v: Q_v = \{q \mid \langle e_a, \uparrow e_a, q \rangle \in \llbracket des \rrbracket(v)\} \& OPT(p_a, Q_v) \subseteq q_a \}$

(6) *Olegooq pulaarli!*

*Ole=guuq pulaar-li*

Ole=<sub>RPT</sub> visit<sub>T-OPT.3SG</sub>

RPT<sup>></sup>. Say for me [ $\uparrow e_a$ ], let Ole come and visit [me!] (see bk, Ch. 6)

RPT<sup><</sup>. I was to say [for  $\uparrow e_r$ ], let Ole come and visit [ $\uparrow e_r$ ]! ( $e_a = e_r$ )

(6'<) Ole<sup>T=RPT<</sup>...

<sup>T</sup>[ $x | x =_i ole$ ];  $\partial[spk.to_{T\Omega} \langle T\varepsilon, \uparrow T\varepsilon, \downarrow T\varepsilon \rangle];$  <sup>T</sup>[ $p | spk.to_p \langle T\varepsilon, \uparrow T\varepsilon, \downarrow T\varepsilon \rangle$ ];

<sup>T</sup>[ $w | w \in_i T\Omega$ ]<sup>T</sup>; ...

visit<sub>T-OPT-3SG</sub>

$\partial[spk.to_{T\Omega} \langle T\varepsilon, \uparrow T\varepsilon, \downarrow T\varepsilon \rangle]; [w | visit_w \langle T\delta, \uparrow T\varepsilon \rangle]; [p | p =_I \perp \omega];$

$[OPT \langle T\Omega, \underline{des}_{T\omega} T\varepsilon \rangle \subseteq_i \perp \Omega]; [T\Omega =_I T\omega |_{T\Omega}]$

(6''<)  $\{ \langle w_{<r}, p_{<r}, \mathbf{o}, w_r, p_r, e_r \rangle, \langle q_a, v_a \rangle \mid \mathbf{o} = \llbracket ole \rrbracket$

$\& w_r \in p_r \subseteq \{v \mid \langle e_r, \uparrow e_r, \downarrow e_r \rangle \in \llbracket spk.to \rrbracket(v)\}$

$\& v_a \in q_a = \{v \mid \langle \mathbf{o}, \uparrow e_r \rangle \in \llbracket visit \rrbracket(v)\}$

$\& w_{<r} \in p_{<r} \subseteq \{v \mid \langle e_r, \uparrow e_r, \downarrow e_r \rangle \in \llbracket spk.to \rrbracket(v) \&$

$\exists Q_v: Q_v = \{q \mid \langle e_r, \uparrow e_r, q \rangle \in \llbracket des \rrbracket(v)\} \& OPT(p_{<r}, Q_v) \subseteq q_a \}$

## 3 KALAALLISUT INFLECTIONS AS TOP-LEVEL ANAPHORA

(7) [Yesterday the school kids had a dog sled race.]

*Ole-p ikinngut-ni ajugaa-mm-at nuannaar-utigi-pa-a.*Ole-ERG<sup>T</sup> friend-3SG<sub>T</sub>.<sup>⊥</sup> win-FCT<sub>⊥</sub>-3SG<sub>⊥</sub> happy-about-DEC<sub>T</sub>.<sup>⊥</sup>-3SG.3SGOle<sup>T</sup>, whose<sub>T</sub> friend<sup>⊥</sup> won,<sup>⊥</sup> was happy about it<sub>⊥</sub>.(7') *Model:*

- $e_a$  is a speech event in  $\{w_1, w_2, w_3\} =: p_{123}$
- Ole (**o**) has friend **b** in  $\{w_1, w_2\}$ , friend **d** in  $\{w_2\}$
- **b** wins in  $\{w_1, w_2\} =: p_{12}$ , **d** wins in  $\{w_3\} =: p_3$
- Ole (**o**) is happy about inhabiting  $p_{12}$  in  $\{w_2\} =: p_2$

Ole-ERG <sup>T</sup>	friend-3SG <sub>T</sub> . <sup>⊥</sup>	[ <sub>FCT</sub> win-	]-FCT <sub>⊥</sub> -3SG <sub>⊥</sub>
<sup>T</sup> [ $x   x =_i ole$ ];	[ $x   frn_{T\omega} \langle x, T\delta \rangle$ ];	[ $w   win_w \langle \perp \delta \rangle$ ];	[ $T\Omega \subseteq_I \perp \omega   _{\perp \delta}$ ];
$C_2$		$C_3$	$C_4$
$\langle \langle \mathbf{o}, w_1, p_{12}, e_a \rangle, \langle \mathbf{b} \rangle \rangle$		$\langle \langle \mathbf{o}, w_1, p_{12}, e_a \rangle, \langle w_1, \mathbf{b} \rangle \rangle$	$\langle \langle \mathbf{o}, w_1, p_{12}, e_a \rangle, \langle w_1, \mathbf{b} \rangle \rangle$
		$\langle \langle \mathbf{o}, w_1, p_{12}, e_a \rangle, \langle w_2, \mathbf{b} \rangle \rangle$	$\langle \langle \mathbf{o}, w_1, p_{12}, e_a \rangle, \langle w_2, \mathbf{b} \rangle \rangle$
$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle \mathbf{b} \rangle \rangle$		$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_1, \mathbf{b} \rangle \rangle$	$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_1, \mathbf{b} \rangle \rangle$
		$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_2, \mathbf{b} \rangle \rangle$	$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_2, \mathbf{b} \rangle \rangle$
$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle \mathbf{d} \rangle \rangle$		$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_3, \mathbf{d} \rangle \rangle$	
: (other $T\Omega$ -values)		: (other $T\Omega$ -values)	: (other $T\Omega$ -values)
[ <sub>DEC</sub>	happy-	]-DEC <sub>T</sub> . <sup>⊥</sup> -3SG.3SG	
$\partial [spk.to_{T\Omega} \langle T\epsilon, \uparrow T\epsilon, \downarrow T\epsilon \rangle$ ];	[ $happy.abt_{T\omega} \{ T\delta, \perp \omega   _{\perp \delta} \}$ ];	<sup>T</sup> [ $p   p =_I T\omega   _{T\Omega}$ ]	
$\checkmark C_4$	$C_5$	$C_6$	
	$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_1, \mathbf{b} \rangle \rangle$	$\langle \langle p_2, \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_1, \mathbf{b} \rangle \rangle$	
	$\langle \langle \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_2, \mathbf{b} \rangle \rangle$	$\langle \langle p_2, \mathbf{o}, w_2, p_{12}, e_a \rangle, \langle w_2, \mathbf{b} \rangle \rangle$	
	: (other $T\Omega$ -values)	: (other $T\Omega$ -values)	

(7'')  $\{ \langle \langle p'_a, \mathbf{o}, w_a, p_a, e_a \rangle, \langle w_x, x \rangle \rangle \mid \mathbf{o} = [[ole]] \ \&$  $\exists q_x: w_x \in q_x = \{v \mid x \in [[win]](v)\}$  $\ \& p_a \subseteq \{v \in q_x \mid \langle e_a, \uparrow e_a, \downarrow e_a \rangle \in [[spk.to]](v) \ \& \langle x, \mathbf{o} \rangle \in [[friend]](v)\}$  $\ \& w_a \in p'_a = \{v \in p_a \mid \langle \mathbf{o}, q_x \rangle \in [[happy.abt]](v)\}$ (8) *Naamik nuannaar-utigi-nngit-la-a*

no happy-about-not-DEC-3SG.3SG

No, he<sub>T</sub> was not happy about it<sub>⊥</sub>.

(8') START-UP (no) ...

<sup>T</sup>[ $e | \uparrow e =_i \downarrow T\epsilon, \downarrow e =_i \uparrow T\epsilon$ ]; <sup>T</sup>[ $p | spk.to_p \langle T\epsilon, \uparrow T\epsilon, \downarrow T\epsilon \rangle$ ]; [ $T\Omega \subseteq_I \perp \omega ||_{\perp \delta}$ ];<sup>T</sup>[ $w | w \in_i T\Omega$ ]; ...[<sub>DEC</sub> happy-not- ]-DEC<sub>T</sub>.<sup>⊥</sup>-3SG.3SG  
 $\partial [spk.to_{T\Omega} \langle T\epsilon, \uparrow T\epsilon, \downarrow T\epsilon \rangle$ ];  $\sim [happy.abt_{T\omega} \{ T\delta, \perp \omega ||_{\perp \delta} \}]$ ; <sup>T</sup>[ $p | p =_I T\omega ||_{T\Omega}$ ]

- (8'')  $\{\langle \langle p'_b, w_b, p_b, e_b, p'_a, \mathbf{o}, w_a, p_a, e_a \rangle, \langle w_x, x \rangle \rangle |$   
 $\langle \langle p'_a, \mathbf{o}, w_a, p_a, e_a \rangle, \langle w_x, x \rangle \rangle \in (6) \ \& \ \uparrow e_b = \downarrow e_a \ \& \ \downarrow e_b = \uparrow e_a \ \&$   
 $\exists q_x: q_x = \{v \mid x \in \llbracket win \rrbracket(v)\} \ \& \ p_b \subseteq \{v \in q_x \mid \langle e_b, \uparrow e_b, \downarrow e_b \rangle \in \llbracket spk.to \rrbracket(v)\}$   
 $\ \& \ w_b \in p'_b = \{v \in p_b \mid \langle \mathbf{o}, q_x \rangle \notin \llbracket happy.abt \rrbracket(v)\}$
- (9) *Aap, pilluaqqu-pa-a*  
 yes congratulate-DEC<sub>τ⊥</sub>-3SG.3SG  
 Yes, he<sub>τ</sub> congratulated him<sub>⊥</sub>.
- (9') START-UP (yes) ...  
 $\uparrow [e \mid \uparrow e =_i \downarrow \uparrow \varepsilon, \downarrow e =_i \uparrow \uparrow \varepsilon]; [\text{spk.to}_{\uparrow \Omega} \langle \uparrow \varepsilon, \uparrow \uparrow \varepsilon, \downarrow \uparrow \varepsilon \rangle]; [\uparrow \omega \in_i \uparrow \Omega]; \dots$   
 $\left[ \begin{array}{ccc} \text{DEC} & \text{congratulate-} & \text{]-DEC}_{\tau \perp} \text{-3SG.3SG} \end{array} \right]$   
 $\partial [\text{spk.to}_{\uparrow \Omega} \langle \uparrow \varepsilon, \uparrow \uparrow \varepsilon, \downarrow \uparrow \varepsilon \rangle]; [\text{congratulate}_{\uparrow \omega} \langle \uparrow \delta, \perp \delta \rangle]; \uparrow [p \mid p =_I \uparrow \omega]_{\uparrow \Omega}$
- (9'')  $\{\langle \langle p'_b, e_b, p'_a, \mathbf{o}, w_a, p_a, e_a \rangle, \langle w_x, x \rangle \rangle |$   
 $\langle \langle p'_a, \mathbf{o}, w_a, p_a, e_a \rangle, \langle w_x, x \rangle \rangle \in (6) \ \& \ \uparrow e_b = \downarrow e_a \ \& \ \downarrow e_b = \uparrow e_a$   
 $\ \& \ p'_a \subseteq \{v \mid \langle e_b, \uparrow e_b, \downarrow e_b \rangle \in \llbracket spk.to \rrbracket(v)\}$   
 $\ \& \ w_a \in p'_b = \{v \in p'_a \mid \langle \mathbf{o}, x \rangle \in \llbracket congratulate \rrbracket(v)\}$

*Figure 1 Centering TAMP-universals*

- (T) TNS fills, or pushes down, the verb's time argument with a dref anchored to a top-ranked time and/or event ( $\uparrow \tau, \perp \tau, \uparrow \varepsilon, \perp \varepsilon$ ).
- (A) ASP fills, or pushes down, the verb's eventuality argument with a dref anchored to a top-ranked state and/or event ( $\uparrow \sigma, \perp \sigma, \uparrow \varepsilon, \perp \varepsilon$ ).
- (M) MOOD fills, or pushes down, the verb's world argument with a dref anchored to a top-ranked world and/or event ( $\uparrow \omega, \perp \omega, \uparrow \varepsilon, \perp \varepsilon$ ).
- (P) PRN fills the verb's subject or object argument with a dref anchored to a top-ranked individual and/or event ( $\uparrow \delta, \perp \delta, \uparrow \varepsilon, \perp \varepsilon$ ).

#### 4 WHY ONLY DECLARATIVES HAVE TRUTH CONDITIONS

BASIC IDEAS:

- All sentences (any MOOD, with or without RPT) express *updates* (a la Lewis 1972, Hamblin 1973, Bittner 2009, 2011, Murray 2010, ...)
- Only *declaratives* (DEC) express **at-issue** updates (D6, cf. AnderBois *et al.* '10) In reportative declaratives (RPT...DEC), *topic-setting* RPT-update is *not* at-issue,  $\uparrow \omega$ -comment in its scope is *at-issue* (cf. Murray 2010) *Illocutionary/context-setting* MOOD-updates (DEC, OPT, FCT, ...) are *not* at-issue
- Only *at-issue updates* are negotiable, e.g. raise the issue of **truth value** (see D7)

## APPENDIX A:

 $UC_{\delta^+}$  WITH ILLOCUTIONARY CENTERING ( $UC_{\omega^+}$ , Bittner 2012: Ch. 6)

D1 The set of  $UC_{\omega^+}$  types is the smallest set  $\Theta$  s.t.: (i)  $t, \delta, \varepsilon, \omega, s \in \Theta$ , (ii)  $(ab) \in \Theta$  if  $a, b \in \Theta$ . The subset  $DR(\Theta) = \{\delta, \varepsilon, \omega, \delta t, \varepsilon t, \omega t\}$  is the set of *dref types*.

D2 A  $UC_{\omega^+}$  frame is a set  $\mathcal{F} = \{\mathcal{D}_a \mid a \in \Theta\}$  such that:

- i.  $\mathcal{D}_t = \{0, 1\}$ ,  $\mathcal{D}_\delta$ ,  $\mathcal{D}_\varepsilon$ , and  $\mathcal{D}_\omega$  are non-empty pairwise disjoint sets
- ii.  $\mathcal{D}_s = \cup_{n \geq 0, m \geq 0} \{\langle \langle d_1, \dots, d_n \rangle, \langle d'_1, \dots, d'_m \rangle \rangle : d_i, d'_j \in \mathcal{D}_{dr}\}$ ,  
where  $\mathcal{D}_{dr} = \cup \{\mathcal{D}_a : a \in DR(\Theta)\}$
- iii.  $\mathcal{D}_{ab} = \{f \mid \emptyset \subset \text{Dom } f \subseteq \mathcal{D}_a \ \& \ \text{Ran } f \subseteq \mathcal{D}_b\}$

D3 A  $UC_{\omega^+}$  model is a tuple  $\mathcal{M} = \langle \mathcal{F}, \uparrow, \downarrow, \prec, \succ, \llbracket \cdot \rrbracket \rangle$  where  $\mathcal{F}$  is a  $UC_{\omega^+}$  frame and  $\llbracket \cdot \rrbracket$  is an interpretation function that maps  $A \in \text{Con}_a$  to  $\llbracket A \rrbracket \in \mathcal{D}_a$ . Moreover:

- i.  $x = \uparrow e \Rightarrow e \in \mathcal{D}_\varepsilon \ \& \ x \in \mathcal{D}_\delta$  (central individual)  
 $y = \downarrow e \Rightarrow e \in \mathcal{D}_\varepsilon \ \& \ y \in \mathcal{D}_\delta \ \& \ \exists x: x = \uparrow u$  (background individual)
- ii.  $e' = \prec e \Rightarrow e, e' \in \mathcal{D}_\varepsilon \ \& \ \uparrow e' \neq \uparrow e$  (source)  
 $e' = \succ e \Rightarrow e \in \mathcal{D}_{\varepsilon'} \ \& \ e' \in \mathcal{D}_{\omega\varepsilon} \ \& \ \forall w \in \text{Dom } e': \uparrow(e'_w) = \downarrow e$  (directive)
- iii.  $\forall w \in \mathcal{D}_\omega, e \in \mathcal{D}_{\varepsilon'}, d, d' \in \mathcal{D}_\delta$ :  
 $\langle e, d, \dots \rangle \in \llbracket A \rrbracket(w) \Leftrightarrow \langle e, \uparrow e, \dots \rangle \in \llbracket A \rrbracket(w) \ \& \ \uparrow e = d$   
 $\langle e, d, d' \dots \rangle \in \llbracket A \rrbracket(w) \Leftrightarrow \langle e, d, \downarrow e \dots \rangle \in \llbracket A \rrbracket(w) \ \& \ \downarrow e = d'$
- iv.  $\forall a \in DR(\Theta), i \in \mathcal{D}_s$ :  
 $\llbracket \top a \rrbracket(i) \doteq ((\textcircled{1}i)_a)_1$      $\llbracket \top' a \rrbracket(i) \doteq ((\textcircled{1}i)_a)_2$      $\llbracket \top \Rightarrow a \rrbracket(i) \doteq \times \{((\textcircled{1}i)_a)_n : n \geq 1\}$   
 $\llbracket \perp a \rrbracket(i) \doteq ((\textcircled{2}i)_a)_1$      $\llbracket \perp' a \rrbracket(i) \doteq ((\textcircled{2}i)_a)_2$      $\llbracket \perp \Rightarrow a \rrbracket(i) \doteq \times \{((\textcircled{2}i)_a)_n : n \geq 1\}$

D4 ( $UC_{\omega^+}$  syntax) For any type  $a \in \Theta$ , define the set of *a-terms*,  $Term_a$ :

- i.  $\text{Con}_a \cup \text{Var}_a \subseteq Term_a$
- ii.  $(A_a = B_a) \in Term_t$  , if  $A_a, B_a \in Term_a$   
 $(A_s \lesssim B_s) \in Term_t$  , if  $A_s, B_s \in Term_s$
- iii.  $\neg \varphi, (\varphi \wedge \psi) \in Term_t$  , if  $\varphi, \psi \in Term_t$
- iv.  $\exists u_a \varphi \in Term_t$  , if  $u_a \in \text{Var}_a$  and  $\varphi \in Term_t$
- v.  $\lambda u_a(B) \in Term_{ab}$  , if  $u_a \in \text{Var}_a$  and  $B \in Term_b$
- vi.  $BA \in Term_b$  , if  $B \in Term_{ab}$  and  $A \in Term_a$
- vii.  $\uparrow A, \downarrow A \in Term_\delta$  , if  $A \in Term_\varepsilon$   
 $\prec A \in Term_\varepsilon \ \& \ \succ A \in Term_{\omega\varepsilon}$  , if  $A \in Term_\varepsilon$
- viii.  $(A_a \top \cdot B), (A_a \perp \cdot B) \in Term_s$  , if  $a \in DR(\Theta), A_a \in Trm_a$  and  $B \in Term_s$
- ix.  $(A \top; B), (A \perp; B) \in Term_{(st)st}$  , if  $A, B \in Trm_{(st)st}$

D5 ( $UC_{\omega+}$  semantics) For any  $\mathcal{M} = \langle \mathcal{F}, \uparrow, \downarrow, <, >, \llbracket \cdot \rrbracket \rangle$  and  $\mathcal{M}$ -assignment  $g$ :

- i.  $\llbracket A \rrbracket^g = \llbracket A \rrbracket$  , if  $A \in Con_a$   
 $\llbracket A \rrbracket^g = g(A)$  , if  $A \in Var_a$
- ii.  $\llbracket (A_a = B_a) \rrbracket^g = 1$  , if  $\llbracket A_a \rrbracket^g = \llbracket B_a \rrbracket^g$ ; else, 0  
 $\llbracket (A_s \lesssim B_s) \rrbracket^g = 1$  , if  $\llbracket A_s \rrbracket^g \lesssim_s \llbracket B_s \rrbracket^g$ ; else, 0
- iii.  $\llbracket \neg \varphi \rrbracket^g = 1$  , if  $\llbracket \varphi \rrbracket^g = 0$ ; else, 0  
 $\llbracket (\varphi \wedge \psi) \rrbracket^g = 1$  , if  $\llbracket \varphi \rrbracket^g = 1$  and  $\llbracket \psi \rrbracket^g = 1$ ; else, 0
- iv.  $\llbracket \exists u_a \varphi \rrbracket^g = 1$  , if  $\{d \in \mathcal{D}_a \mid \llbracket \varphi \rrbracket^{g[u/d]} = 1\} \neq \emptyset$ ; else, 0
- v.  $\llbracket \lambda u_a (B) \rrbracket^g(d) \doteq \llbracket B \rrbracket^{g[u/d]}$  , if  $d \in \mathcal{D}_a$
- vi.  $\llbracket BA \rrbracket^g \doteq \llbracket B \rrbracket^g(\llbracket A \rrbracket^g)$
- vii.  $\llbracket fA \rrbracket^g \doteq f(\llbracket A \rrbracket^g)$  , if  $f \in \{\uparrow, \downarrow, <, >\}$
- viii.  $\llbracket (A^\top \bullet B) \rrbracket^g \doteq \langle (\llbracket A \rrbracket^g \cdot \textcircled{1}\llbracket B \rrbracket^g), \textcircled{2}\llbracket B \rrbracket^g \rangle$   
 $\llbracket (A^\perp \bullet B) \rrbracket^g \doteq \langle \textcircled{1}\llbracket B \rrbracket^g, (\llbracket A \rrbracket^g \cdot \textcircled{2}\llbracket B \rrbracket^g) \rangle$
- ix.  $c\llbracket (A^\top; B) \rrbracket^g = \{k \in c\llbracket A \rrbracket^g\llbracket B \rrbracket^g \mid \exists i \in c \exists j \in c\llbracket A \rrbracket^g \exists a \in DR(\Theta): (\textcircled{1}j)_1 \in \mathcal{D}_a$   
 $\& \textcircled{1}i < \textcircled{1}j \& (\textcircled{1}j)_a = (\textcircled{1}k)_a \& \llbracket B \rrbracket^g \neq \llbracket B[\top a / \perp a] \rrbracket^g\}$   
 $c\llbracket (A^\perp; B) \rrbracket^g = \{k \in c\llbracket A \rrbracket^g\llbracket B \rrbracket^g \mid \exists i \in c \exists j \in c\llbracket A \rrbracket^g \exists a \in DR(\Theta): (\textcircled{2}j)_1 \in \mathcal{D}_a$   
 $\& \textcircled{2}i < \textcircled{2}j \& (\textcircled{2}j)_a = (\textcircled{2}k)_a \& \llbracket B \rrbracket^g \neq \llbracket B[\perp a / \top a] \rrbracket^g\}$

D6 For  $(st)st$ -term  $K$ , model  $\mathcal{M}$ , assignment  $g$ , info-state  $c$ , and  $i \in c$ :

- i.  $I_{\mathcal{M},g}(i, c) = \llbracket \lambda w (\exists k (Ik \wedge \top \Omega i = \top \Omega k \wedge w = \top \omega k)) \rrbracket^{g[i^i][c]}$
- ii.  $K$  is *at-issue* iff  $\exists \mathcal{M} \exists c \exists i \in c \exists j: i \lesssim_s j \& \forall g: I_{\mathcal{M},g}(j, c\llbracket K \rrbracket^g) \subset I_{\mathcal{M},g}(i, c)$

D7 (truth, falsity) For  $(st)st$ -term  $K$ , model  $\mathcal{M}$ , info-state  $c$ , and world  $w$ :

- i.  $\mathcal{M}, c, w \models K$  iff  $K$  is at-issue &  $\exists i \in c \forall g: w \in I_{\mathcal{M},g}(i, c\llbracket K \rrbracket^g)$
- ii.  $\mathcal{M}, c, w \not\models K$  iff  $K$  is at-issue &  $\neg \exists i \in c \forall g: w \in I_{\mathcal{M},g}(i, c\llbracket K \rrbracket^g)$

D8 For any speech act  $e \in \mathcal{D}_e$ :

- i.  $e$  is a *real speech act* iff  $\exists w \in \mathcal{D}_\omega: \langle e, \uparrow e, \downarrow e \rangle \in \llbracket \underline{spk.to} \rrbracket(w)$ .
- ii.  ${}^{st}e = \lambda \langle \langle w, p, e \rangle, \langle \rangle \rangle \mid w \in p \subseteq \{v \mid \langle e, \uparrow e, \downarrow e \rangle \in \llbracket \underline{spk.to} \rrbracket(v)\}$   
is the *e-startup info-state*.

APPENDIX B: DRT-ABBREVIATIONS FOR  $UC_{\omega^+}$ 

Figure 5. Drt-abbreviations for  $UC_{\omega^+}$ -terms

- static conditions (type  $t$ )

$$A_{at} \subseteq B_{at} \quad := \quad \forall u_a (Au \rightarrow Bu)$$

$$B_w(A_1, \dots, A_n) \quad := \quad BwA_1 \dots A_n$$

$$B_p(A_1, \dots, A_n) \quad := \quad p \subseteq \lambda w (BwA_1 \dots A_n)$$

- local condition (type  $st$ )

$$B_w \langle A_1, \dots, A_n \rangle \quad := \quad \lambda i (B_{w \circ i} (A_1 \circ i, \dots, A_n \circ i))$$

- global condition (type  $s(st)t$ )

$$B_w \{A_1, \dots, A_n\} \quad := \quad \lambda i \lambda I (B_{w * i I} (A_1 * i I, \dots, A_n * i I))$$

- global update (type  $(st)st$ )

$$\partial K \quad := \quad \lambda I \lambda j (Ij \leftrightarrow (KI = I))$$

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