

## Plan

- Review:
  - scope prediction
  - Kalaallisut data
- Analysis of Kalaallisut data
- Questions & discussion

# Scope in Kalaallisut: Analysis in CCG+UC<sub>2</sub>

Maria Bittner

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## Scope prediction

- The scope of **BA** (morphologically bound argument) and any **iv-** or **cn-**modifiers is unambiguous

Kalaallisut bound pronouns: Wide scope only

- (Last month Ole<sup>T</sup> ordered three books<sup>L</sup>.)

- Transitive S<sup>+</sup>...-pn<sub>L</sub>: wide only

Suli   atuagaq   ataasiq   tigu-nngi(t)-la-a-Ø.  
still   book<sub>1</sub>   one<sub>2</sub>   receive-not-DEC-3S<sub>(T)</sub>-3S<sub>(L)</sub>

Ǝ→. one book still missing

- Passive S<sup>+</sup>...-pn<sub>L</sub>: wide only

Suli   atuagaq   ataasiq   tigu-niqa(r)-nngi(t)-la-q.  
still   book<sub>T</sub>   one<sub>r</sub>   receive-pssv-not-DEC-3S<sub>(T)</sub>

Ǝ→. one book still missing

Kalaallisut argument *base- | -suffix*: Narrow scope only

- (Last month Ole<sup>T</sup> ordered three books<sup>⊥</sup>.)

- Antipassive  $s^+ \dots -antip$ : narrow only

*Suli* *atuakkamik* *ataatsimik* *tigusinngilaq.*  
*Suli* *atuagaq-mik* *ataasiq-mik* *tigu-si-nngi(t)-la-q.*  
 still book-MOD<sub>δ</sub> one-<sub>δ</sub>MOD receive-antip-not-DEC-3S<sub>(T)</sub>  
 - $\exists$ . hasn't received any

- Noun incorporation  $+s \dots cn-i\backslash cn-$ : narrow only

*Suli* *ataatsimik* *atuagarsinngilaq*  
*Suli* *ataasiq-mik* *atuagaq-si-nngi(t)-la-q.*  
 still one-<sub>δ</sub>MOD book-rcv-not-DEC-3S<sub>(T)</sub>  
 - $\exists$ . hasn't received any

Multiple *cn*-modifiers:  $[ [+s \dots cn-\dots] +s ]$  only

- (Yesterday I saw a bear near the village. And today...)

Ole *alla-mik* *nanu-si-pu-q* *angisuu-mik.*  
 Ole other-<sub>δ</sub>MOD bear-see-DEC<sub>IV</sub>-3S<sub>(T)</sub> big-<sub>δ</sub>MOD  
 Ole saw another bear, a big one.

- (Yesterday I saw a big bear near the village. And today...)

Ole *angisuu-mik* *nanu-si-pu-q* *alla-mik.*  
 Ole big-<sub>δ</sub>MOD bear-see-DEC<sub>IV</sub>-3S<sub>(T)</sub> other-<sub>δ</sub>MOD  
 Ole (too) saw a big bear, another one.

## Kalaallisut Lexicon (part 1)

LEXICAL CATEGORIES (iv = s\pn, tv = iv\pn)

- roots

book-	cn: $\lambda \underline{x} [bk(\underline{x})]$
one-	cn: $\lambda \underline{x} [? \delta \in \underline{x}]; [\underline{x} \in ? \delta]$
other-	cn: $\lambda \underline{x} [? \delta \in \underline{x}]; [\underline{x} \neq_i ? \delta]$
big-	cn: $\lambda \underline{x} [big\{\underline{x}, \underline{x}\}]$
receive-	tv: $\lambda \underline{y} \lambda \underline{x} ([e]^\perp; [rcv(\perp \varepsilon, \underline{x}, \underline{y})])$

- derivational suffixes

-rcv	iv\cn: $\lambda \underline{P} \lambda \underline{x} (\underline{P} \perp \delta^\perp; ([e]^\perp; [rcv(\perp \varepsilon, \underline{x}, \perp \delta)]))$	-si
-see	iv\cn: $\lambda \underline{P} \lambda \underline{x} (\underline{P} \perp \delta^\perp; ([e]^\perp; [see(\perp \varepsilon, \underline{x}, \perp \delta)]))$	-si
-antip	iv\tv: $\lambda \underline{R} \lambda \underline{x}. \underline{R} \underline{x}$ BCK $\langle \perp \varepsilon \rangle \underline{x}$	-si   -(ss)i l...
-pssv	iv\tv: $\lambda \underline{R} \lambda \underline{x}. \underline{R} \underline{x}$ CTR $\langle \perp \varepsilon \rangle$	-niqar l -taa
-not	iv\iv: $\lambda \underline{P} \lambda \underline{x} [\sim (\underline{P} \underline{x})]$	-nngit

## Kalaallisut Lexicon (part 2)

GRAMMATICAL CATEGORIES ( $s^+ = s/s$ ,  $+s = s/s$ )

-DEC	(s\pn)iv: $\lambda \underline{P} \lambda \underline{x}. \underline{P} \underline{x}$	-pu l -pa l -la
- (ERG) <sub>T</sub>	x\cn: $\lambda \underline{P} \lambda K (\underline{P} \top \delta^\top; K)$	(x $\in \{s^+, +s\}$ )
- (ERG) <sub>⊥</sub>	$s^+ \backslash cn: \lambda \underline{P} \lambda K (\underline{P} \perp \delta^\perp; K)$	(x $\in \{s^+, +s\}$ )
-MOD <sub>δ</sub>	$s^+ \backslash cn: \lambda \underline{P} \lambda K (\underline{P} \perp \delta^\perp; (K^\perp; [BCK \perp \varepsilon =_i \perp \delta]))$	-mik
- <sub>δ</sub> MOD	$+s \backslash cn: \lambda \underline{P} \lambda K (K^\perp; \underline{P} \perp \delta)$	-mik
-3S <sub>(T)</sub>	s\pn: $\lambda \underline{P}. \underline{P} \top \delta$	-q   -a...
-3S <sub>(⊥)</sub>	s\pn: $\lambda \underline{P}. \underline{P} \perp \delta$	-Ø   ...

LEXICAL OPERATORS

$\top(\cdot)$ -	cn/cn: $\lambda \underline{P} \lambda \underline{x} ([\underline{x}]^\top; \underline{P} \underline{x})$	$\top \delta$ -accom.
$\perp(\cdot)$ -	cn/cn: $\lambda \underline{P} \lambda \underline{x} ([\underline{y}]^\perp; \underline{P} \underline{x})$	$\perp \delta$ -accom.
$\neg(\cdot)$	(cn\+s)\cn: $\lambda \underline{P} \lambda J \lambda \underline{x}. J(\underline{P} \underline{x})$ (iv\+s)\iv: $\lambda \underline{P} \lambda J \lambda \underline{x}. J(\underline{P} \underline{x})$	cn-lift iv-lift
$\neg\neg(\cdot)$	(s\pn)\pn: $\lambda J \lambda H. H J$	postposed x-lift

TRANSITIVE: Wide scope  $s^+ \dots -pn_{\perp}$  (part 1)

Kalaallisut bound pronouns: Wide scope  $s^+ \dots -pn_{\perp} \text{-}pn_{(T)}$

- (Last month Ole<sup>T</sup> ordered three books<sup>⊥</sup>.)

TRANSITIVE

(suli) atuaqaq ataasiq tigu-nngi(t)-la-a-Ø  
 (still)  $\perp$ book<sub>⊥</sub> one<sub>⊥</sub> receive-not-DEC-3S<sub>(T)</sub>-3S<sub>(⊥)</sub>  
 $\exists \neg$ . one is still missing

'PASSIVE'

(suli) atuaqaq ataasiq tigu-niqar-nngi(t)-la-q  
 (still) Tbook<sub>↑</sub> one<sub>↑</sub> receive-passv-not-DEC-3S<sub>(T)</sub>  
 $\exists \neg$ . one book is still missing

$$\begin{array}{c}
 \frac{\perp \text{book}_{\perp} \text{ one}_{\perp} \dots}{\perp(\cdot)\text{-} \quad \text{book} \quad \perp} \\
 \frac{\text{cn/cn: } \lambda P \lambda x ([y] \perp; P_x) \quad \lambda x [bk(x)] \quad \lambda P \lambda K (P \perp \delta^{\perp}; K)}{\text{cn: } \lambda x ([y] \perp; [bk(x)])} \\
 \frac{\text{cn: } \lambda x ([y] \perp; [bk(x)])}{\exists s: \lambda K ([y] bk(y) \perp; K)} \\
 \frac{\text{one-} \quad \perp(\cdot) \quad \perp}{\text{cn: } \lambda x ([\perp \delta_2 \in x]; [x \in \perp \delta_2]) \quad (\text{cn} \setminus s) \text{ cn: } s^{\perp} \text{ cn: } \lambda P \lambda J x. J(P_x) \quad \lambda P \lambda K (P \perp \delta^{\perp}; K)} \\
 \frac{\text{cn} \setminus s: \lambda J \lambda x. J([\perp \delta_2 \in x]; [\perp \delta \in \perp \delta_2])}{\text{s}^{\perp} (= s/s): \lambda J \lambda K (J([\perp \delta_2 \in \perp \delta]); [\perp \delta \in \perp \delta_2]) \perp; K)} \\
 \frac{\text{s}^{\perp} (= s/s): \lambda K (([y] bk(y)); [\perp \delta_2 \in \perp \delta]); [\perp \delta \in \perp \delta_2]) \perp; K)}{\text{s}^{\perp} (= s/s): \lambda K (([y] bk(y)); [\perp \delta_2 \in \perp \delta]); [\perp \delta \in \perp \delta_2]) \perp; K)}
 \end{array}$$

TRANSITIVE: Wide scope  $s^+ \dots -pn_{\perp}$  (conclusion)

$$\begin{array}{c}
 \frac{\perp \text{book}_{\perp} \text{ one}_{\perp} \dots}{\perp(\cdot)\text{-} \quad \text{book} \quad \perp} \\
 \frac{s^{\perp} (= s/s): \lambda K (([y] bk(y)); [\perp \delta_2 \in \perp \delta]); [\perp \delta \in \perp \delta_2]) \perp; K)}{\text{tv} (= iv \setminus pn): \lambda \underline{y} \lambda x ([e] \perp; [rcv(\perp \varepsilon, x, y)]) \quad \lambda P \lambda x. \underline{P}_x \quad \lambda P. P \perp \delta \quad \lambda P. P \perp \delta} \\
 \frac{\text{tv} (= iv \setminus pn): \lambda \underline{y} \lambda x ([e] \perp; [rcv(\perp \varepsilon, x, y)])}{\text{receive-} \quad \text{not} \quad \text{-DEC} \quad \text{-3S}_{(T)} \quad \text{-3S}_{(⊥)}} \\
 \frac{\text{tv} (= iv \setminus pn): \lambda \underline{y} \lambda x ([e] \perp; [rcv(\perp \varepsilon, x, y)])}{\text{tv} (= iv \setminus pn): \lambda \underline{y} \lambda x \neg ([e] \perp; [rcv(\perp \varepsilon, x, y)])} \\
 \frac{\text{(s/pn)pn: } \lambda \underline{y} \lambda x \neg ([e] \perp; [rcv(\perp \varepsilon, x, y)])}{\text{s/pn: } \lambda \underline{y} \neg ([e] \perp; [rcv(\perp \varepsilon, \tau \delta, y)])} \\
 \frac{\text{s/pn: } \lambda \underline{y} \neg ([e] \perp; [rcv(\perp \varepsilon, \tau \delta, y)])}{\text{s: } [\neg [el rcv(e, \tau \delta, \perp \delta)]]} \\
 \frac{\text{s: } [\neg [el rcv(e, \tau \delta, \perp \delta)]]}{\text{s: } ([y] bk(y)); [\perp \delta_2 \in \perp \delta]; [\perp \delta \in \perp \delta_2]; [\neg [el rcv(e, \tau \delta, \perp \delta)]])
 }
 \end{array}$$

'PASSIVE': Wide scope  $s^+ \dots -pn_T$  (part 1)

$$\begin{array}{c}
 \frac{\perp \text{book}_{\perp} \text{ one}_{\perp} \dots}{\perp(\cdot)\text{-} \quad \text{book} \quad \perp} \\
 \frac{\text{cn/cn: } \lambda P \lambda x ([x] T; P_x) \quad \lambda x [bk(x)] \quad \lambda P \lambda K (P \perp \delta^T; K)}{\text{cn/cn: } \lambda x ([x] T; [bk(x)])} \\
 \frac{\text{cn/cn: } \lambda x ([x] T; [bk(x)])}{\exists s: \lambda K ([x] bk(x) T; K)} \\
 \frac{\text{one-} \quad \perp(\cdot) \quad \perp}{\text{cn: } \lambda x ([\perp \delta_2 \in x]; [x \in \perp \delta_2]) \quad (\text{cn} \setminus s) \text{ cn: } s^{\perp} \text{ cn: } \lambda P \lambda J x. J(P_x) \quad \lambda P \lambda K (P \perp \delta^T; K)} \\
 \frac{\text{cn} \setminus s: \lambda J \lambda x. J([\perp \delta_2 \in x]; [\perp \delta \in \perp \delta_2])}{\text{s}^{\perp} (= s/s): \lambda J \lambda K (J([\perp \delta_2 \in \perp \delta]); [\perp \delta \in \perp \delta_2]) T; K)} \\
 \frac{\text{s}^{\perp} (= s/s): \lambda K (([x] bk(x)); [\perp \delta_2 \in \perp \delta]); [\perp \delta \in \perp \delta_2]) T; K)}{\text{s}^{\perp} (= s/s): \lambda K (([x] bk(x)); [\perp \delta_2 \in \perp \delta]); [\perp \delta \in \perp \delta_2]) T; K)}
 \end{array}$$

'PASSIVE': Wide scope  $s^+ \dots -\text{pn}_T$  (conclusion)

- $\text{book}_T \text{ one}_T \dots$ 
  - $s^+ (= s/s): \lambda K(([\mathbf{x}] bk(\mathbf{x})); [\perp \delta_2 \in T \delta ||]; [T \delta \in \perp \delta_2 ||])^T; K)$
- ... T hasn't been received.
 

receive-	-passv	-not	-DEC	$\text{-3s}_{(T)}$
tv:	iv\tv:	iv\iv:	(s\pn)\iv:	$s(s\pn):$
$\lambda y \lambda x([e]^\perp; [rcv(\perp \varepsilon, x, y)])$	$\lambda R \lambda y. R y \text{ CTR}(\perp \varepsilon)$	$\lambda P \lambda x. \sim(P x)$	$\lambda P \lambda x. P x$	$\lambda P. P \top \delta$
- $s: [\sim[e] rcv(e, \text{CTR } e, \top \delta)]$

Kalaallisut 'object' reduction:  
Narrow scope  $s^+ \dots -\text{antip} | +s \dots cn-$

- (Last month Ole<sup>T</sup> ordered three books<sup>⊥</sup>.)
- (suli) *atuakkamik* *ataatsimik* *tigusinngilaq*  
(suli) *atuaga(q)-mik* *ataasi(q)-mik* *tigu-si-nngi(t)-la-q*  
(still) *book-MOD<sub>δ</sub>* *one-<sub>δ</sub>MOD* *(receive-**antip**)-not-DEC-3s<sub>(T)</sub>*  
 $\neg \exists$ . hasn't received even one book
- (suli) *ataatsimik* *atuagarsinngilaq*  
(suli) *ataasi(q)-mik* *atuaga(q)-si-nngi(t)-la-q*  
(still) *one-<sub>δ</sub>MOD* *(book-)-rcv-not-DEC-3s<sub>(T)</sub>*  
 $\neg \exists$ . hasn't received even one book

ANTIPASSIVE: Narrow scope  $s^+ \dots -\text{antip}$  (part 1)

- $\text{book-MOD}_\delta \text{ one-}_\delta \text{MOD} \dots$ 

$\perp(\cdot)$ -	book	$-\text{MOD}_\delta$
cn/cn	en:	$s^+ \backslash cn$
$\lambda P \lambda x([y]^\perp; P x)$	$\lambda y [bk(y)]$	$\lambda P \lambda K(P \perp \delta^\perp; (K^\perp; [\text{BCK}(\perp \varepsilon) =_i \perp \delta]))$
- $s^+ (= s/s): \lambda K([\mathbf{y}] bk(\mathbf{y})^\perp; (K^\perp; [\text{BCK}(\perp \varepsilon) =_i \perp \delta]))$
- $\text{one-} \quad -\delta \text{MOD}$ 

cn:	$+s \backslash cn:$
$\lambda x([\perp \delta_2 \in x   ]; [x \in \perp \delta_2   ])$	$\lambda P \lambda K(K^\perp; P \perp \delta)$
- $s^+ (= s/s): \lambda K(K^\perp; ([\perp \delta_2 \in \perp \delta ||]; [\perp \delta \in \perp \delta_2 ||]))$

ANTIPASSIVE: Narrow scope  $s^+ \dots -\text{antip}$  (conclusion)

- $\text{book-MOD}_\delta \text{ one-}_\delta \text{MOD} \dots$ 

$\perp(\cdot)$ -	$-\text{antip}$	$\perp(\cdot)$	$-\text{not-DEC-3s}_{(T)}$
tv:	iv\tv	$(iv \backslash s^+) \backslash iv:$	$s \backslash iv:$
$\lambda y \lambda x([e]^\perp; [rcv(\perp \varepsilon, x, y)])$	$\lambda R \lambda x. R \text{ BCK}(\perp \varepsilon) x$	$\lambda P \lambda x. \perp(P x)$	$\lambda P \lambda x. \sim(P \top \delta)$
- $iv: \lambda x([e]^\perp; [rcv(\perp \varepsilon, x, \text{BCK } \perp \delta)])$
- $iv \backslash s^+: \lambda \perp x. \perp([e]^\perp; [rcv(\perp \varepsilon, x, \text{BCK } \perp \delta)])$
- $s \backslash s^+: \lambda \perp x. [\sim el rcv(e, \top \delta, \text{BCK } e)]$
- $s: [\sim([\mathbf{y}] bk(\mathbf{y})]; [el rcv(e, \top \delta, \text{BCK } e)]; [\perp \delta_2 \in \perp \delta ||]; [\perp \delta \in \perp \delta_2 ||]; [\text{BCK}(\perp \varepsilon) =_i \perp \delta])]$
- $s: [\sim([\mathbf{y}] bk(\mathbf{y})]; [\perp \delta_2 \in \perp \delta ||]; [\perp \delta \in \perp \delta_2 ||]; [el rcv(e, \top \delta, \text{BCK } e)]; [\text{BCK}(\perp \varepsilon) =_i \perp \delta])]$
- $s: [\sim([\mathbf{y}] bk(\mathbf{y})]; [\perp \delta_2 \in \perp \delta ||]; [\perp \delta \in \perp \delta_2 ||]; [el rcv(e, \top \delta, \perp \delta)])]$

‘NOUN INCORPORATION’: Narrow scope  ${}^+s$  ...  $cn\text{-}$  (part 1)

- one-  $\neg_{\delta}^{\text{MOD}}$
- $$\frac{\text{cn: } \lambda x([ \perp \delta_2 \in \underline{x} ]); [ \underline{x} \in \perp \delta_2 ]]}{{}^+s (\text{cn: } \lambda P \underline{K}(K^\perp; P \perp \delta))} <$$
- $${}^+s (= s \backslash s): \lambda K(K^\perp; ([ \perp \delta_2 \in \perp \delta ]; [ \perp \delta \in \perp \delta_2 ]))$$
- ... T hasn’t [received any  $\underline{J}$ -book].
- |  |                                |   |  |                                   |
|--|--------------------------------|---|--|-----------------------------------|
| $\perp(\text{-})$  | $\text{book-}$                 | $\neg(\text{-})$  | $\text{-rcv}$  | $\text{-not-DEC-3S}_{(T)}$        |
| —  | —                              | —   | —  | —                                 |
| $\text{cn/cn: }$   | $\text{cn: }$                  | $(\text{cn} \backslash s) \text{cn: }$                              | $\text{iv} \backslash \text{cn: }$   | $s \backslash \text{iv: }$        |
| $\lambda P \lambda x([y]^\perp; P \underline{x})$  | $\lambda x[\underline{bk}(x)]$ | $\lambda P \underline{J} \lambda x. \underline{J}(P \underline{x})$ | $\lambda P \lambda x(P \perp \delta^\perp; [e]^\perp; [rcv(\perp \epsilon, x, \perp \delta)])$ | $\lambda P [\neg(P \tau \delta)]$ |
| —  | —                              | —   | —  | —                                 |
| $\text{cn: } \lambda x([y]^\perp; [bk(x)])$  | —                              | —   | —  | —                                 |
| —  | —                              | —   | —  | —                                 |
| $\text{cn} \backslash {}^+s: \underline{J} \lambda x. \underline{J}([y]^\perp; [bk(x)])$   | —                              | —   | —  | —                                 |
| —  | —                              | —   | —  | $\text{B}$                        |
| $\text{iv} \backslash {}^+s: \underline{J} \lambda x. \underline{J}([y]^\perp; ([e]^\perp; [rcv(\perp \epsilon, x, \perp \delta)]))$ | —                              | —   | —  | —                                 |
| —  | —                              | —   | —  | $\text{B}$                        |
| $s \backslash {}^+s: \underline{J} \lambda x. \underline{J}([y]^\perp; [el rcv(e, \tau \delta, \perp \delta)])$                      | —                              | —   | —  | —                                 |

‘NOUN INCORPORATION’ vs. Antipassive: Comparison

- ‘NOUN INCORPORATION’
- |  |   |
|--|---|
| $\text{one-} \neg_{\delta}^{\text{MOD}}$   | $\text{V}(\underline{book}-)\text{-rcv}\text{-not-DEC-3S}_{(T)}$  |
| —  | —   |
| $\text{+s } (= s \backslash s):$   | $s \backslash {}^+s:$   |
| $\lambda K(K^\perp; ([ \perp \delta_2 \in \perp \delta ]; [ \perp \delta \in \perp \delta_2 ]))$   | $\lambda J[\sim \underline{J}([y] \underline{bk}(y))^\perp; [el rcv(e, \tau \delta, \perp \delta)])$  |
| —  | —   |
| $s: [\sim ([y] \underline{bk}(y))^\perp; ([ \perp \delta_2 \in \perp \delta ]; [ \perp \delta \in \perp \delta_2 ])]^\perp; [el rcv(e, \tau \delta, \perp \delta)])$ | $s: [\sim ([y] \underline{bk}(y)); [ \perp \delta_2 \in \perp \delta ]; [ \perp \delta \in \perp \delta_2 ]; [el rcv(e, \tau \delta, \perp \delta)])$ |
- ANTIPASSIVE
- |   |  |
|---|--|
| $\text{book-MOD}_{\delta} \text{ one-} \neg_{\delta}^{\text{MOD}}$  | $\text{V}(\text{receive-antip-})\text{-not-DEC-3S}_{(T)}$  |
| —   | —  |
| $\text{s}^* \text{ } (= s \backslash s):$   | $s \backslash {}^*s:$  |
| $\lambda K([\underline{y} \underline{bk}(y)]^\perp; (K^\perp; ([ \perp \delta_2 \in \perp \delta ]; [ \perp \delta \in \perp \delta_2 ]; [BCK(\perp \epsilon) =_i \perp \delta])))$           | $\lambda J[\sim \underline{J}[el rcv(e, \tau \delta, BCK \epsilon)]]$  |
| —   | —  |
| $s: [\sim ([y] \underline{bk}(y)); [el rcv(e, \tau \delta, BCK \epsilon)]; [ \perp \delta_2 \in \perp \delta]; [ \perp \delta \in \perp \delta_2 ]; [BCK(\perp \epsilon) =_i \perp \delta])]$ | $s: [\sim ([y] \underline{bk}(y)); [ \perp \delta_2 \in \perp \delta]; [ \perp \delta \in \perp \delta_2 ]; [el rcv(e, \tau \delta, \perp \delta)])$ |

Multiple  $cn$ -modifiers:  $[{}^+s \dots cn\text{-} \dots] {}^+s$  only

- (Yesterday I saw a bear near the village. And today...)
- Ole alla-mik nanu-si-pu-q angisuu-mik.  
 $\text{T Ole}_T \text{ other-} \delta^{\text{MOD}} \text{ V}(\text{bear-})\text{-see-DEC}_{iv}\text{-3S}_{(T)} \text{ V}(\text{big-} \delta^{\text{MOD}})$   
Ole saw another bear, a big one.
- (Yesterday I saw a big bear near the village. And today...)
- Ole angisuu-mik nanu-si-pu-q alla-mik.  
 $\text{T Ole}_T \text{ big-} \delta^{\text{MOD}} \text{ V}(\text{bear-})\text{-see-DEC}_{iv}\text{-3S}_{(T)} \text{ V}(\text{other-} \delta^{\text{MOD}})$   
Ole (too) saw a big bear, another one.

Multiple  $cn$ -modifiers:  $[{}^+s \dots cn\text{-} \dots] {}^+s$  only (part 1)

- T saw a  $\underline{J}(\underline{J}'\text{bear})$  ...
- |   |                             |   |   |
|---|-----------------------------|---|---|
| $\perp(\text{-})$   | $\text{bear-}$              | $\neg(\text{-})$  | $\neg(\text{-})$                                      |
| —   | —                           | —   | —   |
| $\text{cn/cn: }$  | $\text{cn: }$               | $(\text{cn} \backslash s) \text{cn: }$                              | $(\text{cn} \backslash s) \text{cn: }$                |
| $\lambda P \lambda x([y]^\perp; P \underline{x})$   | $\lambda x[\text{bear}(x)]$ | $\lambda P \underline{J} \lambda x. \underline{J}(P \underline{x})$ | $\lambda P \lambda x. \underline{J}(P \underline{x})$ |
| —   | —                           | —   | —   |
| $\text{cn: } \lambda x([y]^\perp; [\text{bear}(x)])$  | —                           | —   | —   |
| —   | —                           | —   | —   |
| $\text{cn} \backslash {}^+s: \underline{J} \lambda x. \underline{J}([y]^\perp; [\text{bear}(x)])$   | —                           | —   | —   |
| —   | —                           | —   | $\text{B}$  |
| $(\text{cn} \backslash {}^+s) \backslash {}^+s: \underline{J} \lambda x. \underline{J}(\underline{J}([y]^\perp; [\text{bear}(x)]))$   | $\text{-see}$               | $\text{-DEC}_{iv}\text{-3S}_{(T)}$                                  | —   |
| —   | —                           | —   | —   |
| $\text{iv} \backslash \text{cn: }$  | $s \backslash \text{iv: }$  | —   | —   |
| $\lambda P \lambda x(P \perp \delta^\perp; ([e]^\perp; [\text{see}(\perp \epsilon, x, \perp \delta)]))$   | $\lambda P. P \tau \delta$  | —   | —   |
| —   | —                           | —   | $\text{B}$  |
| $s \backslash \text{cn: } \lambda P(P \perp \delta^\perp; [el \text{see}(e, \tau \delta, \perp \delta)])$   | —                           | —   | —   |
| —   | —                           | —   | $\text{B}$  |
| $(s \backslash {}^+s) \backslash {}^+s: \underline{J} \lambda x. \underline{J}(\underline{J}([y] \underline{\text{bear}}(y))^\perp; [el \text{see}(e, \tau \delta, \perp \delta)])$ | —                           | —   | —   |

$[[^{\text{t}}\text{s} \dots \text{cn}\dots] \text{+s}]$ : Ole saw another bear, a big one.

- T saw a  $\underline{J}$ (other bear) ...
 
$$\begin{array}{c} \text{other-} \quad \neg_{\delta}^{\text{MOD}} \\ \hline \text{cn: } \lambda \underline{x} ([\perp \delta_2 \in \underline{x}[]]; [\underline{x} \neq_i \perp \delta_2]) \quad +s \text{cn: } \lambda \underline{P} \lambda \underline{K} (\underline{K}^{\perp}; \underline{P} \perp \delta) \\ \hline +s (= s \backslash s): \lambda K (K^{\perp}; (\perp \delta_2 \in \underline{K}[]); [\perp \delta \neq_i \perp \delta_2]) \\ \quad \backslash(\underline{(\text{bear})}) \text{-see-DEC}_{\text{IV}} \text{-3s}_{(\text{T})} \\ \hline (s \backslash^{\text{t}}\text{s}) \text{+s: } \lambda \underline{J} \lambda \underline{J} (\underline{J} (\underline{J} \text{y} \mid \text{bear} \langle \text{y} \rangle) \mid^{\perp}; [\text{el see}(e, \top \delta, \perp \delta)]) \\ \hline +s \text{+s: } \lambda \underline{J} (\underline{J} (\text{y} \mid \text{bear} \langle \text{y} \rangle); [\perp \delta_2 \in \perp \delta[]]; [\perp \delta \neq_i \perp \delta_2]) \perp; [\text{el see}(e, \top \delta, \perp \delta)] \\ \hline \dots \text{a big one.} \end{array}$$
- big-  $\neg_{\delta}^{\text{MOD}}$   $\text{-!}(\cdot)$ 

$$\begin{array}{c} \text{big-} \quad \neg_{\delta}^{\text{MOD}} \quad \text{-!}(\cdot) \\ \hline \text{cn: } \lambda \underline{x} [big\{\underline{x}, \underline{x}\}] \quad +s \text{cn: } \lambda \underline{P} \lambda \underline{K} (K^{\perp}; \underline{P} \perp \delta) \quad (s \backslash (s \backslash^{\text{t}}\text{s})) \text{+s: } \lambda \underline{J} \lambda H \underline{H} \underline{J} \\ \hline +s (= s \backslash s): \lambda H \underline{H} (\lambda K (K^{\perp}; [big\{\perp \delta, \perp \delta\}])) \\ \hline s: (\text{y} \mid \text{bear} \langle \text{y} \rangle); [\perp \delta_2 \in \perp \delta[]]; [\perp \delta \neq_i \perp \delta_2]; [big\{\perp \delta, \perp \delta\}]; [\text{el see}(e, \top \delta, \perp \delta)] \end{array}$$

$[[^{\text{t}}\text{s} \dots \text{cn}\dots] \text{+s}]$ : Ole saw a big bear, another one

- T saw a  $\underline{J}$ (big bear) ...
 
$$\begin{array}{c} \text{big-} \quad \neg_{\delta}^{\text{MOD}} \\ \hline \text{cn: } \lambda \underline{x} [big\{\underline{x}, \underline{x}\}] \quad +s \text{cn: } \lambda \underline{P} \lambda \underline{K} (\underline{K}^{\perp}; \underline{P} \perp \delta) \\ \hline +s (= s \backslash s): \lambda K (K^{\perp}; [big\{\perp \delta, \perp \delta\}]) \\ \quad \backslash(\underline{(\text{bear})}) \text{-see-DEC}_{\text{IV}} \text{-3s}_{(\text{T})} \\ \hline (s \backslash^{\text{t}}\text{s}) \text{+s: } \lambda \underline{J} \lambda \underline{J} (\underline{J} (\underline{J} \text{y} \mid \text{bear} \langle \text{y} \rangle) \mid^{\perp}; [\text{el see}(e, \top \delta, \perp \delta)]) \\ \hline s \text{+s: } \lambda \underline{J} (\underline{J} (\text{y} \mid \text{bear} \langle \text{y} \rangle); [big\{\perp \delta, \perp \delta\}]) \perp; [\text{el see}(e, \top \delta, \perp \delta)] \\ \hline \dots \text{another one.} \end{array}$$
- other-  $\neg_{\delta}^{\text{MOD}}$   $\text{-!}(\cdot)$ 

$$\begin{array}{c} \text{other-} \quad \neg_{\delta}^{\text{MOD}} \quad \text{-!}(\cdot) \\ \hline \text{cn: } \lambda \underline{x} ([\perp \delta_2 \in \underline{x}[]]; [\underline{x} \neq_i \perp \delta_2]) \quad +s \text{cn: } \lambda \underline{P} \lambda \underline{K} (K^{\perp}; \underline{P} \perp \delta) \quad (s \backslash (s \backslash^{\text{t}}\text{s})) \text{+s: } \lambda \underline{J} \lambda H \underline{H} \underline{J} \\ \hline +s (= s \backslash s): \lambda H \underline{H} (\lambda K (K^{\perp}; ([\perp \delta_2 \in \perp \delta[]]; [\perp \delta \neq_i \perp \delta_2]))) \\ \hline s: (\text{y} \mid \text{bear} \langle \text{y} \rangle); [big\{\perp \delta, \perp \delta\}]; [\perp \delta_2 \in \perp \delta[]]; [\perp \delta \neq_i \perp \delta_2]; [\text{el see}(e, \top \delta, \perp \delta)] \end{array}$$