

**Establishment of a Dialectical Logic Symbol System: Inspired by Hegel's Logic and
Buddhist Philosophy**

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Abstract

This paper presents an original dialectical logic symbol system designed to transcend the limitations of traditional logical symbols in capturing subjectivity, qualitative aspects, and contradictions inherent in the human mind. By introducing new symbols, such as “ \ddot{o} ” (being) and “ \emptyset ” (nothing), and arranging them based on principles of symmetry, the system’s operations capture complex dialectical relationships essential to both Hegelian philosophy and Buddhist thought. The operations of this system are primarily structured around the categories found in Hegel’s Logic, and it allows users to incorporate their own subjectivity into the logical processes, opening up new possibilities in the philosophy of subjectivity. This symbol system also has the potential to help us explore fundamental questions of language and to precisely describe processes of consciousness transformation through symbols. This form of logic offers a new, irreducible tool for qualitative methods, and it will spark a new philosophical reflection on its relationship with traditional logic and mathematical symbols.

Keywords: dialectical logic, Hegelian philosophy, symbolic logic, Buddhist thought, consciousness

In this paper, I introduce the dialectical logic symbol system I have created, which was initially published in a book in June 2020, “When Language Ceases: The Symbols of Hegel’s Logic and Buddhology” (Lin Chia Jen, 2020). The original system utilized Chinese characters, but in this paper, I have converted these into more succinct Greek letters. Dialectics, long employed in metaphysics, religion, and everyday psychological processes, has lacked a formal logical system of its own. I am deeply influenced by Descartes’ demand for clarity and distinctness in truth as mentioned in his *Meditations* (Descartes, 2008)¹, as well as by Hegel’s use of complex technical terms and philosophical language that profoundly articulate dialectical methods (Hegel, 1812). Similarly, the same dialectical approach is methodically presented in the Buddhist Agamas (Buddhist Agamas, n.d.), which also inspired me.

In the dialectical logic symbol system I have created, I introduced new logical symbols such as “ δ being,” “ \emptyset nothing,” “ $\dot{\epsilon}$ self,” and “ \cdot abstraction or concretization.” Additionally, I symbolized the “copula,” which is traditionally implicit in any propositional symbol of conventional propositional logic and yet lacks a symbol of its own, as “ \equiv affirmation.” This aspect is particularly significant because “ \equiv affirmation” in this new symbol

¹When Descartes explored the necessity of knowledge through methodological doubt, he also delved into how the existence of the object of knowledge affects this necessity. This inspired me to more firmly incorporate the concept of being into the logical symbols.

system functions as a logical operator, unlike in traditional propositional logic where it does not carry operational meaning and serves only as a presupposed semantic link. This indicates that the operations of dialectical logic symbols are not on “propositions,” but rather on what Hegel referred to as “the Absolute.” (Hegel, 1975) In dialectical logic, an input item p or x does not include a subject-predicate structure, thus detaching from the truthfulness and possibility of correspondence between subject and predicate, making the term “the Absolute” aptly significant.

However, since p or x do not represent propositions, they cannot convey the traditional knowledge significance concerning the truth or falsehood values related to “reality and property relationships.” A p does not encapsulate a copula and a subject, nor does it refer like a predicate does. Thus, the value of p or x clearly lies within themselves. Given the limited knowledge we have acquired from dialectical methods, I can only describe such self-valuable items p or x as having “spiritual value.”

I encourage readers to first abandon the habit of expecting a word to have a specific referent, and to consider the impact of language itself on the mind, such as the “spiritual value” brought by a poem. Unlike poetic and other literary languages that still possess qualities of “referentiality,” “functionality,” and “correspondence,” the symbol system I am about to introduce, using p or x , completely detaches from these aspects. The symbols p or x ,

endowed with spiritual value, do not refer to anything specific; thus, their function and significance lie solely in the symbols themselves, devoid of the ambiguities often associated with traditional literary expressions. This perspective, which views linguistic symbols not merely as tools for referring to reality but as having the power to shape reality itself, aligns with Heidegger's concept that "language is the house of Being" (Heidegger, 1971). From this viewpoint, language serves as an intermediary that constructs the world of meaning. It posits that the truth of beings must unfold through language, which further shapes our perception of reality. However, under this perspective, although the spiritual function of language is similar to what I intend to express, the form of language itself is not precise enough to clearly understand how it shapes the real world in a manner akin to the laws of the physical world.

Thus, to more precisely define this 'spiritual value', in this dialectical logic symbol system, the items p or x must be considered as having complete intrinsic spiritual value. This kind of spiritual effect, stemming from symbols, is commonly found in religious contexts, such as "mantras." Mantras are often hoped to have an immediate spiritual effect, enabling the mind to reach a certain state. However, a "mantra" still carries a certain "image" and thus includes a certain "correspondence relationship," which does not allow us to depart from the way of thinking in propositional logic. Since the logic system I am creating differs from the propositional logic approach, my design uses the operation of symbols to completely break

away from this correspondence relationship, one example being the transformation of the copula “is” into the symbol \equiv .

This spiritual value is further manifested in another crucial characteristic of this logical system: the dialectical logic symbol system is entirely “self-referential.” In the logical operations within this system, any item—be it x , y , or p —retains its inherent value regardless of changes in its logical position or its combinations with other items. This “self-referential” nature introduces a significant feature, allowing users of the symbols to “insert their self” into the system. This type of self-reference somewhat restores “intentionality,” but it differs from the functional external correspondence that constitutes propositional logic. Instead, it aligns with the type of “intentionality” associated with consciousness as described by Hegel in *Phenomenology of Spirit* (Hegel, 1977) and the “intentionality” in phenomenology by Husserl (Husserl, 2001). This phenomenological perspective explores how intentionality of consciousness is structured in the process of referring and how it interacts with the transformation process between consciousness and its object. Indeed, in my symbol system, the “self (ϵ)” is frequently utilized in the categories of “relationship of substantiality,” “relationship of causality,” “reciprocity,” and “The Notion” within this symbol system. I believe this feature of allowing the symbol’s users to “insert their self” offers a unique value

distinct from previous logical systems, providing a convenient tool for knowledge based fundamentally on “introspection.”

Method

In the creation of this dialectical logic symbol system, the primary focus was on the selection of symbols, their arrangement, and the method of operation. For symbol selection, I primarily referenced Hegel’s Logic (Hegel, 1975) and Buddhist scriptures, as both are fundamentally dialectical in nature, and their structural systems share significant similarities. This allowed me to cross-reference and select appropriate symbols for dialectical logic.

Hegel’s Logic integrates metaphysics and addresses the traditional philosophical question of “being,” while Buddhism similarly considers “subject” as a logical object. Therefore, I chose the Greek symbol δ (from $\delta\nu$) to represent “being,” and the symbol ϵ (from $\epsilon\gamma\acute{\omega}$) to represent “subject.” The introduction of these two symbols is key to distinguishing dialectical logic from other types of logic. As part of the logical system, I also chose \neg to represent “negation,” following traditional logic.

To make this symbol system capable of effectively describing Hegelian philosophy, Buddhist philosophy, and subjective, introspective idealist philosophy in general, I placed particular emphasis on geometric symmetry in the arrangement of symbols. For example, in the structure $\neg \cdot \delta \sim x \cdot \neg$, the \sim symbol separates two symmetric states: “negation abstracted”

and “negation concretized.” This symmetry captures important structures in idealism, such as “subject/object” or “consciousness/object.”

Finally, since dialectics is fundamentally a logic of transformation, I incorporated the principle of change into the operations. For example, I defined that \ddot{o} can be transformed into \emptyset (nothing), and the various structures of dialectical logic follow symmetrical patterns of calculation. In the calculation of the “category of substantiality,” for instance, the transformation of logical expressions shows a self-returning symmetry in the form of “ $\dot{\dot{e}} \rightarrow \dot{e} \cdot \rightarrow \dot{e} \cdot x \rightarrow \dot{e} \cdot \rightarrow \dot{e} \cdot$.”

As this is a completely original symbolic design, its effectiveness still awaits verification through practical application and experience by readers. This process, unlike applied sciences, cannot be validated by experimental data.

Basic Symbols of Dialectics

For ease of writing, I have selected some symbols to replace the original Chinese characters. These symbols are pre-existing, but their arrangement and operational methods are my creations for the purposes of dialectics. The following are the substitute symbols used in the dialectical logic symbol system, which comprises six basic symbols²:

² In my original Chinese work, the Chinese symbols for ‘ $\dot{\dot{e}}$, \equiv , \ddot{o} , \emptyset , \neg , \cdot ’ are respectively ‘我, 是, 有, 無, 不, 的’.

$\dot{\epsilon}$ represents “self”

\equiv represents “affirmation”

\ddot{o} represents “being”

\emptyset represents “nothing”

\neg represents “negation”

\cdot represents “abstraction or concretization”

Basic Concepts and Mastery

Firstly, within this dialectical logic symbol system, an item that can be inserted, such as an “x”, must be considered as “absolute,” “non-composite,” and “indivisible.”

Secondly, this x is not a proposition, meaning it does not inherently contain the copula “is” like the P in propositional logic. Nor is this x a predicate or an individual in predicate logic. In predicate logic, P(x) already implies “is,” and also encompasses the meaning of “having.” In dialectical logic, it is crucial to distinctly separate these aspects.

Thirdly, the dialectical logic symbol system I have developed differs from previous logical and mathematical symbol systems in that those systems could not incorporate the symbol user “I” into their operations; however, this revised symbol system fully allows the inclusion of the symbol user “I” in its calculations, meaning this system is “completely self-referential.”

Fourthly, regarding this x , we can “abstract” it to become “ $x\cdot$ ”; alternatively, it can be “concretized” to become “ $\cdot x$ ”. Ultimately, we can affirm it to become $\equiv x$.

The Introduction and Importance of the Symbol ‘ \cdot ’

The ‘ \cdot ’ symbol represents “abstraction” or “concretization.” For an item x , we can abstract it to become $x\cdot$ (placed to the left of “ \cdot ”); or it can be concretized to become $\cdot x$ (placed to the right of “ \cdot ”). Here, $x\cdot$ signifies that the function of X in thought is suppressed; whereas $\cdot x$ indicates that the function of x in thought is expressed.

The Introduction and Importance of the Symbols \checkmark and \emptyset

\checkmark signifies “being,” while \emptyset represents “nothing.” The introduction of these two symbols is for computing “becoming,” and their mutual transformation is closely related to the negation \neg . Additionally, \checkmark can represent “immediacy” or “mediacy,” whereas \emptyset can be used to eliminate items within a logical formula.

The Introduction of Double Negation

Following the introduction of the symbol “ \cdot ,” we can now discuss the critically important concept of double negation in dialectical logic. The commonly referred to “negation of negation” can be expressed as $\neg\cdot\neg$, which means that after negation \neg is abstracted to ‘ $\neg\cdot$ ’, it then concretizes back to \neg . This represents what Hegel referred to as “self-return” (Hegel, 1975). There is a familiar certainty, that the negation of negation is

indeed an affirmation. This principle is acknowledged in the dialectical logic symbol system, and its certainty is represented by the following self-evident symbol transformations:

$$\neg \cdot \neg$$

$$\equiv$$

or

$$\equiv$$

$$\neg \cdot \neg$$

The dialectical logic symbols use line breaks to represent transformations between logical formulas, showcasing the transformation from the double negation $\neg \cdot \neg$ to the affirmation symbol \equiv and the affirmation symbol \equiv back to double negation $\neg \cdot \neg$.

Affirmation \equiv and its Relationship with Being \checkmark and Nothing \emptyset :

$$\equiv \checkmark$$

$$\emptyset$$

and

$$\equiv \emptyset$$

$$\checkmark$$

The above sets of symbol transformations illustrate the relationship between affirmation \equiv and being \checkmark and nothing \emptyset . This relationship implies that if affirmation \equiv is removed from

$\equiv \check{\emptyset}$ or $\equiv \emptyset$, then being $\check{\emptyset}$ will transform into nothing \emptyset , and nothing \emptyset will transform into being $\check{\emptyset}$.

Further expansion on the relationship between being $\check{\emptyset}$ and nothing \emptyset :

$\check{\emptyset} \neg \cdot \neg$

$\equiv \check{\emptyset}$

\emptyset

Based on the transformation relationship between double negation $\neg \cdot \neg$ and affirmation \equiv , we have the above three formulas' transformation. With these transformations, we can formally enter the operational process of dialectical logic symbols.

Dialectical Logic Operations Require the Decomposition of Double Negation

We introduce a symbol \sim , which lacks specific logical meaning, to decompose double negation into the following formula:

$\neg \cdot \sim \cdot \neg$

The symbol \sim divides the abstraction of negation $\neg \cdot$ and the concretization of negation $\cdot \neg$ into left and right sides, thus generating two logical positions: the position between $\cdot \neg$ and \sim is called the "first logic position," and the position between $\neg \cdot$ and \sim is called the "second logic position." These two logical positions can insert items or determinateness, such as $\check{\emptyset}$, \emptyset , x , y , a , etc. Now, we will represent this in the form of

transformations between being \checkmark and nothing \emptyset .

Entering the Doctrine of Being

In this paper, I will display categories based on the structure of Hegel's Encyclopaedia of the Philosophical Sciences in Outline, specifically the section on logic (Hegel, 1975). However, the operations of the categories I have symbolized will not entirely correspond to those in Hegel's work due to the symbolization process. Some parts have developed their intrinsic nature and no longer fully align with the number of categories Hegel established using language. We now delve into the "The Doctrine of Being" category from Hegel's shorter logic, initiating a particularization process where negation \neg returns to itself, starting with "The Doctrine of Being."

The Category of Becoming

BC

$\neg \sim \checkmark \cdot \neg$

$\neg \cdot \checkmark \sim \neg$

$\neg \cdot \emptyset \sim \neg$

$\neg \sim \emptyset \cdot \neg$

$$\neg \cdot \sim \check{\circ} \cdot \neg$$

The above logical formulas constitute a cycle of formulas, which is referred to as a category. This category represents Hegel's category of becoming (Hegel, 1975). The cycle proceeds as follows: (a) $\check{\circ}$ is first abstracted by negation into $\check{\circ} \cdot \neg$, thereby suppressing $\check{\circ}$ in thought; (b) $\check{\circ}$ then abstracts the negation \neg , concretizing into $\neg \cdot \check{\circ}$; (c) As the first two formulas complete the self-return of $\check{\circ}$ and \neg , they implicitly contain the transformation between $\neg \cdot \neg$ and \equiv , and based on the transformation relationship between $\equiv \check{\circ}$ and \emptyset , reach the third logical formula $\neg \cdot \emptyset \sim \cdot \neg$. Thus, we arrive at the first category: becoming, abbreviated as the function BC.

The Category of Determinate Being

$$DB(x)$$

$$\neg \cdot \check{\circ} \sim x \cdot \neg$$

$$\neg \cdot x \sim \check{\circ} \cdot \neg$$

$$\neg \cdot x \sim \emptyset \cdot \neg$$

$$\neg \cdot \emptyset \sim \cdot \neg$$

$$\neg \cdot \check{\circ} \sim \cdot \neg$$

We now insert an item x , aside from $\check{\circ}$, into the first logic position, resulting in the category of determinate being, abbreviated as DB (). DB(x) represents the function of

inserting x into DB, encompassing five transformations of logical formulas. The part involving $\check{\text{o}}$ has already been explained in BC, and the process x undergoes is the same as that of $\check{\text{o}}$. In the first two logical formulas $\neg \cdot \check{\text{o}} \sim x \cdot \neg$ and $\neg \cdot x \sim \check{\text{o}} \cdot \neg$, x and \neg undergo mutual abstraction and concretization. Noteworthy are the transformations in the third and fourth formulas. When $\neg \cdot x \sim \emptyset \cdot \neg$ transforms into $\neg \cdot \emptyset \sim \cdot \neg$, x is eliminated by \emptyset , leaving only $\check{\text{o}}$ in the fifth formula.

In BC, I mentioned that the cycle of formulas constitutes a category. However, determinate being, unlike BC, is not a category that can return to itself, because the item x we insert is eliminated by \emptyset in the fourth formula. If we continue with the logical formulas of DB, it will return to BC. Thus, DB is a category that requires the insertion of other items such as x , y , a , etc., to return to itself.

Within $\neg \cdot \check{\text{o}} \sim x \cdot \neg$, we have filled the first and second logic positions, forming an entirety within the being-in-self stage. The meaning of the entirety lies in forming a concrete determinateness, unlike in BC, where the absence of items in the first and second logic positions allows the negation \neg to predominantly function in thought. In DB(x), we cannot merely negate in thought because the negation \neg is abstracted by $\check{\text{o}}$, nor can we simply treat $\check{\text{o}}$ as the absolute, because the necessary property x , united with $\check{\text{o}}$, is abstracted by \neg . The

meaning of the entirety is that it is erroneous to correctly analyze any items constituting the entirety, as any analysis or division would inevitably be incorrect.

From a linguistic perspective, $\neg \cdot \check{\text{ö}} \sim x \cdot \neg$ can be translated as 'x's being,' while $\neg \cdot x \sim \check{\text{ö}} \cdot \neg$ can be translated as 'the existing x.' Here, x oscillates between functioning as a noun and an adjective, but it is not yet a fully determined noun. The true noun will only emerge when we enter the category of matter. However, based on the principle that the entirety in DB(x) cannot be divided, although these translations may be used individually in practice, from the perspective of operations in dialectical logic symbols, such naturally isolated language use is always 'one-sided.'

The above two categories represent the categories of The Doctrine of Being, in which the symbol $\check{\text{ö}}$ primarily represents "immediacy." Next, we will enter the category of "The Doctrine of Essence" within the dialectical logic symbol system.

The Category of Essence: Reflective Categories

The Category of Identity

ID(x)

$x \neg \cdot \check{\text{ö}} \sim \cdot \neg$

$\check{\text{ö}} \neg \cdot x \sim \cdot \neg$

or

$$\checkmark \neg \cdot x \sim \cdot \neg$$

$$x \neg \cdot \checkmark \sim \cdot \neg$$

We now enter the category of essence, within which \checkmark can represent both “mediacy” and “immediacy.” The first point of focus is the third logic position, which is the position to the left of $\neg \cdot$. An item can only be inserted to the left of $\neg \cdot$ once an item has been placed in the second logic position. The significance of the third logic position is that it is not subject to the mutual abstraction and concretization relative to the negation \neg , since the negation \neg has already been abstracted by the item in the second logic position. The first category we encounter in the doctrine of essence is the category of identity, abbreviated as ID (). We can insert an x , formatted as ID(x). The function of ID(x) is to allow free interchange between items in the second and third logic positions across two lines of logical formulas, as seen in the interchange between $x \neg \cdot \checkmark \sim \cdot \neg$ and $\checkmark \neg \cdot x \sim \cdot \neg$. This interchange does not cause the being \checkmark to transform into nothing \emptyset . Finally, I define ID () as an operator, used as follows:

$$\checkmark \neg \cdot x \sim \cdot \neg$$

$$x \neg \cdot \checkmark \sim \cdot \neg \text{ID}$$

This indicates that the second logical expression $x \neg \cdot \checkmark \sim \cdot \neg$ is derived from the first

logical expression $\check{\emptyset} \neg \cdot x \sim \cdot \neg$ by applying the operator ID.

The Category of Opposition

OPP(x)

$x \neg \cdot \check{\emptyset} \sim \cdot \neg$

$\neg \cdot x \sim \check{\emptyset} \cdot \neg$

$\neg \cdot \check{\emptyset} \sim x \cdot \neg$

$\check{\emptyset} \neg \cdot x \sim \cdot \neg$

The category of opposition involves moving items located in the third and second logic positions backward in sequence to the second and first logic positions, respectively, and then performing the first two logical formulas of DB (), which interchange the second and first logic positions. However, the next step is not to proceed to the third formula of DB (), but rather to move the interchanged items back to the third and second logic positions. In the operation of the dialectical logic symbol system, a determinateness, such as x, as long as it does not continuously occupy the second and first logic position in relation to \emptyset within three consecutive logical formulas, $\check{\emptyset}$ will not be transformed into \emptyset .

The Category of The Thing Itself

TIF (x, y)

$\checkmark \neg \cdot x \sim y \cdot \neg$

$\checkmark \neg \cdot y \sim x \cdot \neg$

In the category of the thing itself, we have inserted a second item, y, thus filling the first to third logic positions. According to the previous interpretation of DB (), filling the first and second logic positions forms an indivisible whole, so \checkmark in the third logic position is not affected by the negation \neg and can be conceived as absolutely or independently itself. Thus, we have the following rule:

$\checkmark \neg \cdot x \sim y \cdot \neg$

\checkmark

and

$\checkmark \neg \cdot y \sim x \cdot \neg$

\checkmark

The two lines of logical formulas above represent that \checkmark in the third logic position can be independently transformed into a standalone formula. This rule is so definitive that I represent it succinctly in the following single-line transformations using an arrow to denote the logical conversion:

$$\checkmark \neg \cdot x \sim y \cdot \neg \rightarrow \checkmark$$

and

$$\checkmark \neg \cdot y \sim x \cdot \neg \rightarrow \checkmark$$

We can further interpret this: when our mind generates a thought, and the mind does not acknowledge that this thought has been affirmed or negated but appears purely naturally, the category used by the mind is the thing itself, TIF (). This category represents what we often refer to as “thing,” which is typically concrete and not represented as a proposition that can be judged true or false.

In $\checkmark \neg \cdot x \sim y \cdot \neg$ and $\checkmark \neg \cdot y \sim x \cdot \neg$, we can designate x and y as representing “immediacy” and “mediacy” respectively. For example, we might set x as immediacy and y as mediation, then in the entirety of $\checkmark \neg \cdot x \sim y \cdot \neg$, the thing itself or the thing is considered immediate, and in $\checkmark \neg \cdot y \sim x \cdot \neg$, it is considered mediated. However, since these two logical formulas can freely transform, \checkmark in the TIF () category is “simultaneously mediated and immediate.”

This capability for free transformation indicates “contradiction,” which is a frequent occurrence in our consciousness and thinking. The free interchange between $\checkmark \neg \cdot y \sim x \cdot \neg$ and $\checkmark \neg \cdot x \sim y \cdot \neg$ makes “the thing” seem as if it has been defined, yet because neither the first nor the second logic positions can be accurately analyzed, only \checkmark in the third logic position can be transformed, thus making “the thing” appear as purely undefined.

In fact, since only the third logic position can reasonably transform an “item,” and since any item in dialectical logic must be placed within the first to fourth logic positions, we must acknowledge that any item of dialectical logic “itself” comes from this transformation in the third logic position. Additionally, TIF () itself, like ID (), can be used as an operator in logical formulas and can be directly indicated on the far right of the logical formula. Next, I will explain that the logical transformations of this analyzable property are reversible, meaning:

ǒ

ǒ¬·x~y·¬

ǒ¬·y~x·¬TIF

x¬·ǒ~y·¬ ID → x

Since ǒ is a pure logical symbol, we can freely write it out to represent present thinking, and ǒ can be transformed into ǒ¬·x~y·¬, providing us with a rule for generating determinateness from nothing.

In this way, we can generate items such as x or y from nothing, and the items transformed through this arrow → x process, I call “free items,” meaning they can be substituted into the logic positions of either themselves or other categories.

The Category of Matter

MA(x)

$$x \cdot \check{\circ} \sim y \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot y \sim \check{\circ} \cdot \neg$$

$$x \cdot \neg \cdot y \sim \emptyset \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$x \cdot \neg \cdot \check{\circ} \sim \cdot \neg$$

(fi)*MA (): $x^* \text{ MA (1)} \vee x^* \text{ MA (3)}$

Based on the free interchangeability arising from the identity between the third and second logic positions, we exchange the positions of x and $\check{\circ}$ in TIF (), applying the principle of the third logic position's ability to independently analyze an item, and thus analyze x. This item x analyzed from the third logic position is an independent formula, which I refer to as a "free item." By exchanging x with $\check{\circ}$, we complete the process of generating the determinateness x from nothing, originating from TIF ().

This x, initially positioned in the third logic position, can transform into a free x and also enter into the second logical formula $x \cdot \neg \cdot y \sim \check{\circ} \cdot \neg$. The logical formulas displayed above are known as matter, abbreviated in function form as MA (). The process from $x \cdot \neg \cdot \check{\circ} \sim y \cdot \neg$ transforming into $x \cdot \neg \cdot y \sim \check{\circ} \cdot \neg$ is distinct from the categories in the Doctrine of Being where

negation \neg serves as the axis of self-return; here, it is x that can independently transform from the third logic position that serves as the axis of self-return. Thus, the transformation of the first three logical formulas can be simplified as “ $x \rightarrow x \cdot \rightarrow x$,” where x first abstracts into $x \cdot$ and then concretizes back into x , representing x 's self-return.

At the same time, the changes in the first and second logic positions are considered as resulting from changes in x . In the first logical formula, the possibility of x 's mediacy is based on the abstraction of y through the use of negation \neg in thought. If I define x as consciousness and y as sensation, in the first logical formula $x \neg \cdot \check{\text{o}} \sim y \cdot \neg$, the concreteness of consciousness through its identity relationship with $\check{\text{o}}$ in the third logical position, highlights the mediacy of $\check{\text{o}}$ within the mind. Meanwhile, by abstracting sensation, the $\check{\text{o}}$ that merges with y at the second logic position, and coincides with y at the first logic position, loses its immediacy in thought.

By the time we reach $x \cdot \neg \cdot y \sim \check{\text{o}} \cdot \neg$, x abstracts into $x \cdot$, such that x demonstrates its opposition to y , or rather, after y undergoes abstraction through negation \neg , it becomes mutually exclusive with the concrete x . Since the structure of $\neg \cdot y \sim \check{\text{o}} \cdot \neg$ is concrete, I like to say that $x \cdot \neg \cdot \check{\text{o}} \sim y \cdot \neg$ represents x “contacting” with $\neg \cdot \check{\text{o}} \sim y \cdot \neg$ ³, and $x \cdot \neg \cdot y \sim \check{\text{o}} \cdot \neg$ represents x “distancing” from $\neg \cdot y \sim \check{\text{o}} \cdot \neg$.

³In the twelve links of dependent origination in Buddhism, one link is ‘contact,’ which I

When we reach $x \neg \cdot y \sim \emptyset \cdot \neg$, x concretizes or returns to itself, transforming $\check{\emptyset}$ into \emptyset . This represents the identical relationship of x and y in the second and third logic positions, which cannot be framed by $\check{\emptyset}$, set as either “immediacy” or “mediacy,” but should be unrestricted. Originally, I defined consciousness as mediacy and sensation as immediacy, overall, a clear determinateness. However, when I unify consciousness and sensation, it is no longer appropriate to define them using the immediate or mediated $\check{\emptyset}$.

By the time we reach $x \cdot \neg \cdot \emptyset \sim \cdot \neg$, x , not being truly nothing, abstracts into $x \cdot$, indicating that x does not come into contact with \emptyset in the second and third logic positions.

Finally, we arrive at $x \neg \cdot \check{\emptyset} \sim \cdot \neg$, where $x \cdot$ once again returns to x , and \emptyset transforms back into $\check{\emptyset}$. However, at this time, the first logic position is empty, without any item inserted, which could either lead x into the pure category of identity or necessitate the insertion of a free item from elsewhere. It is important to note that in the $MA(x)$ category, the free item x initially transformed cannot be inserted into $x \neg \cdot \check{\emptyset} \sim \cdot \neg$, as this dialectical logic symbol system only allows negation \neg and self $\check{\emptyset}$ to appear twice within the same logical formula. Thus, $MA(x)$ is a category with a free item that cannot insert itself and requires insertion into other categories.

believe is very suitable to represent the concreteness of dialectical thinking in this context.

From a linguistic perspective, $x \cdot \ddot{\sim} y \cdot \neg$ can be translated as 'x possessing the quality of y (with $\ddot{\sim}$ functioning as 'possessing').' However, based on the explanation of the analyzability of the entirety mentioned above, what can truly be translated is only the single x. So, although this structure is a definition structure familiar to us in language, at this stage of the dialectical operation, since x is not yet a determinateness that is both being-in-self and being-for-self, the only thing that can truly be translated is the single noun x. The true definition structure will only emerge when we enter the category of actuality.

Summary of Operators and the Continuum of Thought

I define this type of logical formulas within the category that can generate free items using the arrow \rightarrow , each item can only be generated once. If there are different items, for example, formulas that can generate x, y, \equiv , $\equiv x$, $\equiv y$, $\equiv \dot{\epsilon}$, etc., each of them can only be generated once.

However, if we view the entire system of categories as a continuum—meaning that we can always connect the first and last logical formulas of one category to another in some way—the rule that a specific item can only be generated once within a category will apply to the entire continuum of the dialectical logic symbol system.

Once a continuum of categories has determined its direction based on transformation rules and has generated all possible items, that continuum becomes a definite set of thoughts or knowledge. To continue this continuum, we then enter the inference phase of the dialectical logic symbol system, which involves the process of moving free items between categories.

However, the \rightarrow used to generate free items interacts with the use of ID within a continuum, creating uncertainty before and after the application of ID. This causes a loss of completeness in one of the logical formulas that can be transformed. I will explain this with the second category of substantial relationship, SID1 ():

SID1(x, y)

$$\equiv \check{\sigma} \cdot y \sim x \cdot \neg \rightarrow \equiv \check{\sigma}$$

$$\dot{\epsilon} \equiv \check{\sigma} \cdot x \sim y \cdot \neg \rightarrow \dot{\epsilon} \equiv$$

$$\dot{\epsilon} \equiv x \cdot \check{\sigma} \sim y \cdot \neg \text{ID} \rightarrow \dot{\epsilon} \equiv x$$

The above are the first three logical formulas of SID1(x, y), where the third logical formula indicates that it is the result of using ID on the second logical formula. Without using

ID, the second logical formula would have been able to fully transform into $\dot{\epsilon} \equiv \check{\sigma}$. However, due to the use of ID, it can only transform into $\dot{\epsilon} \equiv$.

All the categories that can be used as operators have now been introduced. Now, I will introduce two ways to organize the free items generated within a continuum. The first type of free items can be found to form part of each other's structure in the collection of transformed free items, which I abbreviate as (fi). The second type cannot, and I call them non-composite free items, abbreviated as (ncfi). For example, if the set of items that can be transformed within the entire continuum are $\dot{\epsilon} \equiv$, $\dot{\epsilon} \equiv x$, $\equiv x$, $\equiv \check{\sigma}$, and $\dot{\epsilon} \cdot \equiv x$, it is clear that $\equiv \check{\sigma}$ and $\dot{\epsilon} \cdot \equiv x$ do not form part of the structure of the other members. After this comparison, we can organize their free items in the categories and continuum as follows (where the free item is placed to the left of *, and the number in parentheses next to the category name indicates the logical formula from which it is transformed):

(fi) * SID1(): $\dot{\epsilon} \equiv * \text{SID1}(2) \wedge \dot{\epsilon} \equiv x * \text{SID1}(3) \wedge \equiv x * \text{SID1}(8)$

(ncfi) * SID1(): $\equiv \check{\sigma} * \text{SID1}(1) \wedge \dot{\epsilon} \cdot \equiv x * \text{SID1}(4)$

The Category of Form

FM (x, y)

$$x \neg \cdot \check{\circ} \sim y \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot y \sim \check{\circ} \cdot \neg$$

$$x \neg \cdot y \sim \emptyset \cdot \neg$$

$$y \neg \cdot x \sim \emptyset \cdot \neg \text{ ID} \rightarrow y$$

$$y \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$y \neg \cdot \check{\circ} \sim \cdot \neg$$

$$(\text{ncfi})^* \text{FM} () : x^* \text{FM} (1) \wedge y^* \text{FM} (4)$$

We now consider other possibilities within MA (x, y). Originally, in the third logical formula of MA (x, y), $x \neg \cdot y \sim \emptyset \cdot \neg$, x and y are in a relationship of identity in the second and third logic positions, so we can use ID to exchange the positions of x and y to become $y \neg \cdot x \sim \emptyset \cdot \neg$. To demonstrate the transformation relationship with the line above $x \cdot \neg \cdot y \sim \check{\circ} \cdot \neg$, we should mark ID on $y \neg \cdot x \sim \emptyset \cdot \neg$, making it $y \neg \cdot x \sim \emptyset \cdot \neg \text{ ID}$.

$y \neg \cdot x \sim \emptyset \cdot \neg \text{ ID}$ indicates that in the unrestricted relationship of \emptyset , y swaps places with x according to ID, and manifests “itself” in the third logic position. According to my earlier setting, y represents sensation, and x represents consciousness. This swap of y and x reflects the mystical nature of dialectical logic. Sensation is inherently immediate, and the manifestation of consciousness is based on negating and abstracting the immediate sensation, yet sensation y is the necessary item that fills the logical positions, meaning that

consciousness x can only manifest itself because it denies the immediate nature of sensation.

Therefore, in $y \neg \cdot x \sim \emptyset \cdot \neg$ ID, the unrestricted nature of \emptyset makes it impossible to distinguish between the determinateness of consciousness and sensation.

Since $y \neg \cdot x \sim \emptyset \cdot \neg$ ID also fills all logic positions, y in the third logic position can also be freely transformed, hence marked as $y \neg \cdot x \sim \emptyset \cdot \neg$ ID $\rightarrow y$. Therefore, the category of form is a category that has two free items.

Note that because the last logical formula $y \neg \cdot \check{\emptyset} \sim \neg$ lacks the first logic position, and the item currently at the third logic position is different from that in the first logical formula, we can insert the x initially transformed into the last formula:

The Category of Force

$F(x, y)$

$FM(x, y)$

$x \neg \cdot \check{\emptyset} \sim y \cdot \neg \rightarrow x$

$x \cdot \neg \cdot y \sim \check{\emptyset} \cdot \neg$

$x \neg \cdot y \sim \emptyset \cdot \neg$

$y \neg \cdot x \sim \emptyset \cdot \neg$ ID $\rightarrow y$

$y \cdot \neg \cdot \emptyset \sim \neg$

$y \neg \cdot \check{\emptyset} \sim \neg \leftarrow x * FM(1)$

FM (y, x)

$y \cdot \neg \sim x \cdot \neg \rightarrow y$

$y \cdot \neg \cdot x \sim \neg$

$y \cdot \neg \cdot x \sim \emptyset \cdot \neg$

$x \cdot \neg \cdot y \sim \emptyset \cdot \neg \text{ID} \rightarrow x$

$x \cdot \neg \cdot \emptyset \sim \neg$

$x \cdot \neg \cdot \emptyset \sim \neg \leftarrow y * \text{FM} (1)$

(ncfi)*F (): y*F (4) \wedge x*F (10)

The category described above is formed when x is inserted into $y \cdot \neg \sim \neg$, establishing a new category called ‘force,’ abbreviated in function form as F (). The category of force is composed of two categories of form, FM (x, y) and FM (y, x), thus it is a ‘composite category.’ The structure of F() involves inserting the free item x from the category of form FM (x, y) into its last logical formula, thereby creating another category of form FM (y, x), where the free item y of this form can then be inserted into its own last logical formula, forming FM(x,y) in an endless cycle. This cyclic insertion of x or y into the last logical formula of the category of form represents the self-returning motion of x or y.

In summary, the category of force still contains two items that are not inserted into themselves; therefore, these two free items can still be inserted into other categories.

However, we must also consider the manner in which the category of matter, MA (x, y), extends and returns to itself. Since the free item of the category MA (x, y) cannot be inserted into its own last formula but only into other categories, and knowing that the category of the thing, TIF, can generate itself from nothing, we first create a TIF to attempt to insert x. Here is the process:

MA (x, y)

$x \cdot \check{\text{o}} \sim y \cdot \neg \text{ID} \rightarrow x$

TIF (x, y)

$\check{\text{o}} \cdot x \sim y \cdot \neg \rightarrow \check{\text{o}}$

$\check{\text{o}} \cdot y \sim x \cdot \neg \text{TIF}$

MA (y, x)

$y \cdot \check{\text{o}} \sim x \cdot \neg \text{ID} \rightarrow y$

$y \cdot \neg \cdot x \sim \check{\text{o}} \cdot \neg$

$y \cdot \neg \cdot x \sim \emptyset \cdot \neg$

$y \cdot \neg \cdot \emptyset \sim \cdot \neg$

$y \cdot \neg \cdot \check{\text{o}} \sim \cdot \neg \leftarrow x \cdot \text{FM} (1)$

Now, inserting the free item x of MA (x, y) into $y \cdot \neg \cdot \check{\text{o}} \sim \cdot \neg$ results in $y \cdot \neg \cdot \check{\text{o}} \sim x \cdot \neg \rightarrow y$, returning us to the first formula of MA (y, x) and, by swapping y with $\check{\text{o}}$, back to TIF (y, x):

$$y \cdot \check{\circ} \sim x \cdot \neg$$

$$\text{TIF } (y, x)$$

$$\check{\circ} \cdot y \sim x \cdot \neg \text{ ID} \rightarrow \check{\circ}$$

$$\check{\circ} \cdot x \sim y \cdot \neg \text{TIF}$$

Then, by swapping x and $\check{\circ}$ using ID, we return to MA (x, y):

$$\text{MA } (x, y)$$

$$x \cdot \check{\circ} \sim y \cdot \neg \text{ ID} \rightarrow x$$

$$x \cdot \neg \cdot y \sim \check{\circ} \cdot \neg$$

$$x \cdot \neg \cdot y \sim \emptyset \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$x \cdot \neg \cdot \check{\circ} \sim \cdot \neg$$

This process, returning to MA (x, y), is termed the category of appearance, abbreviated

as AP (x, y):

The Category of Appearance

$$\text{AP } (x, y)$$

$$\text{TIF } (x, y)$$

$$\check{\circ} \cdot x \sim y \cdot \neg \text{ ID}$$

$$\check{\circ} \cdot y \sim x \cdot \neg \text{TIF}$$

MA (y, x)

$y \cdot \check{\sim} x \cdot \neg \text{ID} \rightarrow y$

$y \cdot \neg \cdot x \sim \check{\sim} \cdot \neg$

$y \cdot \neg \cdot x \sim \emptyset \cdot \neg$

$y \cdot \neg \cdot \emptyset \sim \cdot \neg$

$y \cdot \neg \cdot \check{\sim} \cdot \neg \leftarrow x * \text{MA} (1)$

TIF (y, x)

$\check{\sim} \cdot y \sim x \cdot \neg \text{ID}$

$\check{\sim} \cdot x \sim y \cdot \neg \text{TIF}$

MA (x, y)

$x \cdot \neg \cdot \check{\sim} y \cdot \neg \text{ID} \rightarrow x$

$x \cdot \neg \cdot y \sim \check{\sim} \cdot \neg$

$x \cdot \neg \cdot y \sim \emptyset \cdot \neg$

$x \cdot \neg \cdot \emptyset \sim \cdot \neg$

$x \cdot \neg \cdot \check{\sim} \cdot \neg \leftarrow y * \text{MA} (1)$

(ncfi)*AP (): None

This process of x returning to itself in AP (x, y) differs from that in the category of force, as it needs to pass through the category of the thing. These two modes of self-return are

significant for thought: the category of TIF can transform \checkmark , and considering AP (x, y) where the free items transform as “ $x \rightarrow \checkmark \rightarrow y \rightarrow \checkmark \rightarrow x$,” involves uncertainty between “immediacy” and “mediacy” of \checkmark and the principle of “TIF () $\rightarrow \checkmark$.” Thus, when thought seeks to express ideas involving contradictions and uncertainties, we should employ the category of appearance, AP (); however, when thought intends to express a self-sufficient, automatic, and definite process of self-return, we should use the category of force, F ().

Before concluding the reflective categories of essence, I wish to introduce a special category, a variant of the category of matter, MA (x, y), which I call MA2().

MA2(x, y)

$y \cdot \neg x \sim \checkmark \cdot \neg \rightarrow y$

$x \cdot \neg y \sim \checkmark \cdot \neg$ ID

$x \cdot \neg \checkmark \sim y \cdot \neg$

$x \cdot \neg \emptyset \sim \cdot \neg$

$x \cdot \neg \sim \emptyset \cdot \neg$

$x \cdot \neg \sim \checkmark \cdot \neg \leftarrow y * \text{MA2} (1)$

(ncfi)*MA2 (): None

The foundation of this variant of the matter category is based on the following process:

1. \checkmark can appear arbitrarily.

2. \checkmark should be able to generate determinateness, becoming $\checkmark \neg \cdot x \sim y \cdot \neg$.
3. $\checkmark \neg \cdot x \sim y \cdot \neg \rightarrow x \neg \cdot \checkmark \sim y \cdot \neg$ ID.
4. $x \neg \cdot \checkmark \sim y \cdot \neg \rightarrow x$.
5. x is a symbol as simple as \checkmark .
6. x can also appear arbitrarily.
7. x can accompany a logical structure filled in the first and second logic positions, thus transforming into $x \neg \cdot y \sim \checkmark \cdot \neg$.

In the first two formulas of MA2's category, $y \neg \cdot x \sim \checkmark \cdot \neg$ and $x \neg \cdot y \sim \checkmark \cdot \neg$, x and y are positioned in the second and third logic positions in a relationship of identity, while \checkmark in the first logic position is abstracted by negation. This means that the \checkmark representing immediacy or mediacy is initially suppressed, and the free interchangeability between x and y due to their identity allows their determinateness to permeate each other. A notable aspect of this category is the formula $x \neg \cdot \emptyset \sim \cdot \neg$, where x and \emptyset are in a relationship of identity.

Another characteristic of the MA2 category is that since the first two formulas can respectively transform the free items x and y , as soon as we reach the last logical formula $x \neg \cdot \sim \checkmark \cdot \neg$, y can be inserted to become $x \neg \cdot y \sim \checkmark \cdot \neg$, and continue with the following category of

MA2 matter:

MA2(y, x)

$$x \cdot y \sim \check{o} \cdot \neg \rightarrow x$$

$$y \cdot \neg \cdot x \sim \check{o} \cdot \neg \text{ ID}$$

$$y \cdot \neg \cdot \check{o} \sim x \cdot \neg$$

$$y \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$y \cdot \neg \cdot \sim \emptyset \cdot \neg$$

$$y \cdot \neg \cdot \sim \check{o} \cdot \neg \leftarrow x \cdot \text{MA2 (1)}$$

(ncfi)*MA2 (y, x): None

Connecting reflective category —category of actuality: CRA ()

CRA (x, y)

$$x \cdot y \sim \check{o} \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot \check{o} \sim y \cdot \neg$$

$$x \cdot \neg \cdot \emptyset \sim y \cdot \neg$$

$$\equiv x \cdot \neg \cdot \check{o} \sim y \cdot \neg \rightarrow \equiv x$$

(fi)*CRA (): $x \cdot \text{CRA (1)} \wedge \equiv x \cdot \text{CRA (4)}$

The second category of matter also has a very important property; its first logical formula $x \cdot y \sim \check{o} \cdot \neg$ can be used as a bridge connecting the reflective categories of essence and the category of actuality. When the category of matter progresses to the third logical formula

$x \neg \emptyset \sim y \neg$, adding an affirmation symbol \equiv in front transforms the \emptyset in the second logic position into $\check{\emptyset}$, turning the entire formula into $\equiv x \neg \check{\emptyset} \sim y \neg$, which constitutes the complete structure of the category of actuality in the Doctrine of Essence.

Note: There are more linking formulas, but this introduction only covers this one, as it directly links to the category of actuality itself, rather than to the categories of possibility and contingency.

The Doctrine of Essence: The Category of Actuality

We have now explored how, within the reflective categories of the Doctrine of Essence, x and y returns to itself. Whether through the category of appearance or through the category of force, x and y become determinateness that simultaneously possesses both immediacy and mediacy in the cycle of logical formula. This totality of immediacy and mediacy brings thought into the crucial category of actuality within the Doctrine of Essence.

We know that double negation $\neg \neg$ equates to affirmation \equiv . In previous categories of being and reflective categories of essence, the insertion of free items and their movement with $\check{\emptyset}$ and \emptyset were just intermediary determinateness of double negation $\neg \neg$. This intermediary determinateness describes the particular process of thought where double negation $\neg \neg$ is transformed into affirmation \equiv . However, true affirmation \equiv is not yet

reached until the intermediary particular items x or y return to themselves through the categories of force or appearance.

In the DB () category of being determinate, $\neg \cdot x \sim \checkmark \cdot \neg$ can neither analyze the certainty of negation \neg nor the exact particularized nature of x , thus thought has not yet expressed affirmation \equiv . Hegel refers to the particularized nature of x in the Doctrine of Being as being-in-self, not being-for-self (Hegel, 1975). However, when the reflective categories in the Doctrine of Essence allow x or y to return to themselves “for themselves” through the categories of force F () or appearance AP (), the particularized nature of x or y becomes being-for-self.

In $x \neg \cdot \checkmark \sim y \cdot \neg$, thought can only transform an independent and free item x , represented as $x \neg \cdot \checkmark \sim y \cdot \neg \rightarrow x$. Since $\neg \cdot \checkmark \sim y \cdot \neg$ cannot be understood as affirmation \equiv , we cannot symbolically transform it into a truly being-for-self expression, which would be $\equiv x$.

Now, to truly express a being-for-self particularized nature, we finally enter the category of actuality within the Doctrine of Essence, introducing the symbol of affirmation \equiv . Here is the first logical formula of actuality itself, $AC(x, y)$:

$$\equiv x \neg \cdot \checkmark \sim y \cdot \neg \rightarrow \equiv x$$

With the addition of the affirmation sign \equiv , we gain an additional logic position to the left of the \equiv sign, termed “the fourth logic position”. The item in the fourth logic position is

“being-for-self”, meaning the entire formula $\equiv x \neg \cdot \check{\sim} y \cdot \neg$ can transform into $\equiv x$, not just x .

With the inclusion of the fourth logic position, x 's range of movement within the structure is expanded. Here is the complete operation formula of the AC (x, y) category of actuality:

The Category of Actuality Itself

AC (x, y)

$\equiv x \neg \cdot \check{\sim} y \cdot \neg \rightarrow \equiv x$

$x \equiv y \neg \cdot \check{\sim} \cdot \neg$

$\equiv x \neg \cdot y \sim \check{\sim} \cdot \neg \rightarrow \equiv x$

$\equiv x \cdot \neg \cdot y \sim \emptyset \cdot \neg \rightarrow \equiv x \cdot$

$\equiv x \neg \cdot \emptyset \sim \cdot \neg$

$x \equiv \neg \cdot \check{\sim} \cdot \neg$

$\equiv x \neg \cdot \check{\sim} a \cdot \neg \rightarrow \equiv x$

(fi)*AC (): $\equiv x^*AC (1) \vee \equiv x^*AC (3) \vee \equiv x^*AC (7)$

(ncfi)*AC (): $\equiv x^*AC (4)$

The affirmation \equiv enhances our dialectical logic's ability to capture our mental processes. The certainty expressed by affirmation \equiv manifests in the complete operational formula of the AC (x, y) category of actuality. Notice especially the last logical formula

$\equiv x \neg \cdot \check{\text{o}} \sim a \cdot \neg$, which structurally mirrors the first formula of the category, $\equiv x \neg \cdot \check{\text{o}} \sim y \cdot \neg$, unlike the reflective category where the last formula often lacks an item in the first logic position. This demonstrates that the AC (x, y) category of actuality can allow x to return to itself within its own operations.

Let us closely examine the process within these formulas. Initially, in $\equiv x \neg \cdot \check{\text{o}} \sim y \cdot \neg$, while we could use the ID to swap x and $\check{\text{o}}$ as we do in the categories of the thing and matter, the freely transformable item here is not x or $\check{\text{o}}$, but $\equiv x$ or $\equiv \check{\text{o}}$.

When x in $\equiv x \neg \cdot \check{\text{o}} \sim y \cdot \neg$ is treated as a specific immediacy or mediation, and transformed using ID into $\equiv \check{\text{o}} \neg \cdot x \sim y \cdot \neg$, x and y merge into one in the first and second logic positions making $\check{\text{o}}$ both immediate and mediated. This transformation between particularity and universality allows the transformable item to become uncertain; the only certain transformation is the totality represented by affirmation \equiv .

However, when we do not use ID to transform the first logical formula, when $\equiv x \neg \cdot \check{\text{o}} \sim y \cdot \neg$ is transformed into $x \equiv y \neg \cdot \check{\text{o}} \sim \cdot \neg$, x moves to the fourth logic position, while y and $\check{\text{o}}$ move to the second and third logic positions. This shows that x and “y identical with $\check{\text{o}}$ ” are in an affirmative relationship. This formula represents a “complete affirmation” because no item in the first logic position is abstracted by negation \neg .

As $x \equiv y \cdot \check{\sim} \cdot \neg$ transforms into $\equiv x \cdot y \sim \check{\sim} \cdot \neg$, x's being-for-self particularized nature returns to the third logic position, where x and y are in a relationship of identity, and the $\check{\sim}$ capable of particularizing immediacy or mediacy is abstracted by negation \neg . This differs from the third logical formula of the reflective categories of essence where the identity of x and y is based on \emptyset being abstracted by negation.

From $\equiv x \cdot y \sim \check{\sim} \cdot \neg$ to $\equiv x \cdot y \sim \emptyset \cdot \neg$, and finally to $\equiv x \cdot \emptyset \sim \cdot \neg$, these three transformations represent x's "being-in-self" return within the category of actuality. In $\equiv x \cdot \emptyset \sim \cdot \neg$, note that x and \emptyset are in an identity relationship, representing x's being-in-self particularized nature identical with the nothing \emptyset , unlike the reflective category where x is always identical with $\check{\sim}$.

However, in the category of actuality, x is not merely being-in-self; it is also being-for-self, thus $\equiv x \cdot \emptyset \sim \cdot \neg$ must progress to $x \equiv \cdot \check{\sim} \cdot \neg$.

In the final logical formula $\equiv x \cdot \check{\sim} \sim a \cdot \neg$, x returns to the third logic position, and due to the movement of x shown in the previous formulas representing x's being-in-self and being-for-self return, a latent affirmation arises producing a new particularized nature 'a'. The importance of the category of actuality lies here, because unlike the reflective category, which can only produce a new item through the TIF () category, the category of actuality can produce a new item within its own logical necessity.

Finally, it is important to note that $\equiv x$ cannot be substituted into the first and second logic positions.

The $\equiv x$ generated by the first logical formula should be marked as $\equiv x^*AC$ (1), and the third or fourth logical formula should allow the generation of $\equiv x^*AC$ (3) or $\equiv x^*AC$ (4), but only one place can be chosen to generate $\equiv x$.

In the second logical expression of the AC (x, y) category of actuality, $x \equiv y \neg \cdot \check{\sim} \sim \cdot \neg$, located at the third and fourth logical positions, x and y are in a relationship where they can freely interchange positions. We explore the logical effects produced by this free interchange, transforming $x \equiv \check{\sim} \neg \cdot y \sim \cdot \neg$ into $x \equiv y \neg \cdot \check{\sim} \sim \cdot \neg$ ID and $y \equiv x \neg \cdot \check{\sim} \sim \cdot \neg$ ID. The category formed by this free interchange of the third and fourth logical positions is called the first category of possibility, abbreviated as POS1().

Before entering the remaining categories of actualities there is a special category as follows:

Connect to Nothing

CON (x, y)

$\equiv x \neg \cdot \check{\sim} \sim y \cdot \neg \rightarrow \equiv$

$\equiv \check{\sim} \neg \cdot x \sim y \cdot \neg$ ID $\rightarrow \equiv \check{\sim}$

$\leftarrow \equiv \check{\sim}^*CON$ (2)

$$\equiv \check{\emptyset}$$

$$\emptyset$$

$$(fi)*CON(): \equiv *CON(1)$$

Category of Possibility 1

$$POS1(x, y)$$

$$\equiv x \neg \check{\sim} y \neg \rightarrow \equiv x$$

$$x \equiv \check{\sim} y \neg \rightarrow \equiv$$

$$x \equiv y \neg \check{\sim} \neg ID$$

$$y \equiv x \neg \check{\sim} \neg ID$$

$$\equiv y \neg x \sim \check{\sim} \rightarrow \equiv y$$

$$\equiv y \neg \check{\sim} x \neg \rightarrow \equiv y$$

$$\equiv y \neg \emptyset \sim x \neg \rightarrow \equiv y$$

$$y \equiv \neg \sim \emptyset \neg$$

$$\equiv y \neg \sim \check{\sim}$$

$$(fi)*POS1(): \equiv *POS1(2) \wedge [\equiv y*POS1(5) \vee \equiv y*POS1(7)]$$

$$(ncfi)*POS1(): \equiv x*POS1(1) \wedge y.*POS1(6)$$

Since at the first and second logical positions, \checkmark and determinateness need to be in a relationship of mutual concretization and abstraction for \checkmark to transform into \emptyset , when $\equiv y \cdot \neg \cdot x \sim \checkmark \cdot \neg$ progresses to $\equiv y \cdot \neg \cdot \checkmark \sim x \cdot \neg$, it does not transform \checkmark into \emptyset like in the AC(x,y) category of actuality. Following this, the last logical expression produced by the self-returning movement of y, being both being-in-itself and being-for-itself, $\equiv y \cdot \neg \cdot \sim \checkmark \cdot \neg$, creates a vacancy at the second logical position. This vacancy must be filled by a free item transformed from the reflective category.

Category of Possibility 2

POS2(x, y)

$$\equiv x \cdot \neg \cdot \checkmark \sim y \cdot \neg \rightarrow \equiv x$$

$$x \equiv \checkmark \cdot \neg \cdot y \sim \cdot \neg$$

$$x \equiv y \cdot \neg \cdot \checkmark \sim \cdot \neg \text{ ID}$$

$$y \equiv x \cdot \neg \cdot \checkmark \sim \cdot \neg \text{ ID}$$

$$y \equiv \checkmark \cdot \neg \cdot x \sim \cdot \neg \text{ ID}$$

$$\equiv y \cdot \neg \cdot \checkmark \sim x \cdot \neg \rightarrow \equiv y$$

$$\equiv y \cdot \neg \cdot x \sim \checkmark \cdot \neg \rightarrow \equiv y \cdot$$

$$\equiv y \cdot \neg \cdot x \sim \emptyset \cdot \neg \rightarrow \equiv y$$

$$y \equiv \neg \cdot \emptyset \sim \cdot \neg$$

$$\equiv y \neg \cdot \check{\emptyset} \sim \cdot \neg$$

$$(fi)*POS2(): \equiv y* POS2(6) \vee \equiv y* POS2(8)$$

$$(ncfi)*POS2(): \equiv x* POS2(1) \wedge \equiv y \cdot *POS2(7)$$

The second category of possibility, abbreviated as POS2(), differs from POS1(x, y) by utilizing ID three times consecutively. This results in a vacancy occurring in the first logical position in the last logical expression, $\equiv y \neg \cdot \check{\emptyset} \sim \cdot \neg$. Similarly, this vacancy must be filled by a free item transformed from the reflective category.

Now, let us look at the last category of actuality, which is the category of contingency.

Category of Contingency

$$CONT(x, y)$$

$$\equiv y \neg \cdot x \sim \check{\emptyset} \cdot \neg \rightarrow \equiv$$

$$\equiv x \neg \cdot y \sim \check{\emptyset} \cdot \neg \text{ ID } \rightarrow \equiv x$$

$$x \equiv y \neg \cdot \check{\emptyset} \sim \cdot \neg$$

$$\equiv x \neg \cdot \check{\emptyset} \sim y \cdot \neg \rightarrow \equiv x$$

$$\equiv x \cdot \neg \cdot \emptyset \sim y \cdot \neg \rightarrow \equiv x \cdot$$

$$\equiv x \neg \cdot \sim \emptyset \cdot \neg$$

$$x \equiv \neg \cdot \sim \check{\emptyset} \cdot \neg$$

$$\equiv x \neg \cdot b \sim \check{o} \cdot \neg \rightarrow \equiv$$

$$\equiv b \neg \cdot x \sim \check{o} \cdot \neg \text{ID} \rightarrow \equiv b$$

$$(\text{fi})^* \text{CONT}(): [\equiv^* \text{CONT}(1) \vee \equiv^* \text{CONT}(8)] \wedge [\equiv x^* \text{CONT}(2) \vee \equiv x^* \text{CONT}(4)]$$

$$(\text{ncfi})^* \text{CONT}() \equiv x^* \text{CONT}(5) \wedge \equiv b^* \text{CONT}(9)$$

The category of contingency, abbreviated as CONT (), begins with the second logical expression $\equiv x \neg \cdot y \sim \check{o} \cdot \neg$ which originates from the linkage formula in the category of matter:

$$\text{CRA2}(x, y)$$

$$x \neg \cdot \check{o} \sim y \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot x \sim \check{o} \cdot \neg$$

$$x \neg \cdot y \sim \check{o} \cdot \neg \text{ID} \rightarrow x$$

$$\equiv x \neg \cdot y \sim \check{o} \cdot \neg$$

We observe that the CONT (x, y) category, similar to the AC (x, y) category of actuality, can generate a new item b in its final logical expression. However, the self-sufficient and self-determined nature of x in the CONT (x, y) category is not as definitive as in the AC (x, y) category, which can only transform into $\equiv x$, whereas the category of contingency can transform into $\equiv x$, $\equiv y$, or $\equiv b$.

Now, let's review the entirety of the category of actuality. The category of actuality encompasses the first and second categories of possibility, POS1() and POS2(), and the

category of contingency, CONT (). In these categories, the self-sufficient and self-determined items or terms change, such as transforming from $\equiv x$ to $\equiv y$ or $\equiv b$, whereas the category of actuality itself, AC (), maintains the same item from start to finish, which is " $\equiv x \rightarrow \equiv x$ "—this represents "necessity." Let us continue to explore the category of actuality itself, AC (x, y):

AC (x, y)

$$\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg \rightarrow \equiv x$$

$$x \equiv y \neg \cdot \check{\circ} \sim \cdot \neg$$

$$\equiv x \neg \cdot y \sim \check{\circ} \cdot \neg \rightarrow \equiv x$$

$$\equiv x \cdot \neg \cdot y \sim \emptyset \cdot \neg \rightarrow \equiv x \cdot$$

$$\equiv x \neg \cdot \emptyset \sim \cdot \neg$$

$$x \equiv \neg \cdot \check{\circ} \sim \cdot \neg$$

$$\equiv x \neg \cdot \check{\circ} \sim a \cdot \neg \rightarrow \equiv x$$

$$(fi)^*AC (): \equiv x^*AC (1) \vee \equiv x^*AC (3) \vee \equiv x^*AC (7)$$

$$(ncfi)^*AC (): \equiv x \cdot^*AC (4)$$

AC (x, a)

$$\equiv x \neg \cdot \check{\circ} \sim a \cdot \neg \rightarrow \equiv x$$

$$x \equiv a \cdot \check{o} \sim \cdot \neg$$

$$\equiv x \neg \cdot a \sim \check{o} \cdot \neg \rightarrow \equiv x$$

$$\equiv x \cdot \neg \cdot a \sim \emptyset \cdot \neg \rightarrow \equiv x \cdot$$

$$\equiv x \neg \cdot \emptyset \sim \cdot \neg$$

$$x \equiv \neg \cdot \check{o} \sim \cdot \neg$$

$$\equiv x \neg \cdot \check{o} \sim c \cdot \neg \rightarrow \equiv x$$

$$(fi)*AC(): \equiv x*AC(1) \vee \equiv x*AC(3) \vee \equiv x*AC(7)$$

$$(ncfi)*AC(): \equiv x*AC(4)$$

We discover that if we extend the category of actuality to become AC (x, a), we continue to generate a new free item c, and the transformations still ultimately produce $\equiv x$.

The first, second categories of possibility, POS1() and POS2(), and the category of contingency, CONT (), relative to the category of actuality itself AC (), add flexibility in changing the both being-in-self and being-for-self free items. However, the first and second categories of possibility may have a vacancy in the last logical expression, while the category of contingency inherently carries an uncertainty in thought.

From a linguistic perspective, $\equiv x \neg \cdot \check{o} \sim y \cdot \neg$ can be translated as '...is x possessing the quality of y (with \check{o} functioning as "possessing").' At this stage of dialectical logic, since x has already acquired necessity within the category of actuality itself and is accompanied by \equiv ,

we can acknowledge that the structure of the category of actuality can formally be translated into a definitional structure.

Now, as the mind seeks to preserve the necessity of the category of actuality itself, AC (x, y), while also maintaining the flexibility of self-sufficient and self-determined free items, it will continue to delve deeper into the “category of substantial relationships.”

Category of Substantial Relationships

In my dialectical symbol system, readers will experience for the first time “what it means to treat the self as a purely symbolic entity.” In philosophical reflection, we understand the self as something pure, not as an experiential term used in everyday language to refer to this body or brain. For instance, Kant in his Critique of Pure Reason suggests that the “I think” accompanies all perceptions and thoughts, constituting the transcendental structure of the knowing subject (Kant, 1998). This not only provides the self with an innate, logical foundation but also positions the self transcendently as the center that orchestrates cognitive activities. For Kant, the transcendental self bestows unity on experience, ensures the coherence and possibility of individual experiences, and actively interprets sensory data according to the laws of causality and substance. These philosophical reflections underscore the self’s role as a purely functional and logical entity in human experience.

However, to date, no one has treated the self as a logical symbol. Considering the self's abstract qualities, such as invisibility and inaudibility, it seems plausible that it could be embraced as a logical symbol.

Category of Substantiality S ()

S (x, y)

$$\equiv x \neg \check{\text{o}} \sim y \neg \rightarrow \equiv x$$

$$\dot{\text{e}} \equiv x \neg y \sim \check{\text{o}} \neg \rightarrow \dot{\text{e}} \equiv x$$

$$\dot{\text{e}} \cdot \equiv x \neg y \sim \emptyset \neg \rightarrow \dot{\text{e}} \cdot \equiv x$$

$$\dot{\text{e}} \cdot x \equiv \neg \emptyset \sim \neg$$

$$\dot{\text{e}} \cdot \equiv x \neg \check{\text{o}} \sim \neg$$

$$\dot{\text{e}} \equiv x \neg \sim \check{\text{o}} \neg$$

$$\equiv x \neg a \sim \check{\text{o}} \neg \rightarrow \equiv x$$

$$(fi)*S () : [\equiv x *S (1) \vee \equiv x *S2 (7)] \wedge \dot{\text{e}} \equiv x *S (2)$$

$$(ncfi)*S () : \dot{\text{e}} \cdot \equiv x *S (3)$$

In this first category of substantiality, abbreviated as S (), we introduce the last symbol of the dialectical logic symbol system, the “self $\dot{\text{e}}$.” Like “being $\check{\text{o}}$,” “self $\dot{\text{e}}$ ” is a purely logical symbol, pre-assigned to the fourth logic position. The difference between $\dot{\text{e}}$ and $\check{\text{o}}$ is that (a) $\check{\text{o}}$

transforms into \emptyset , while \acute{e} does not; (b) \check{o} is regarded as “to be particularized as either immediacy or mediacy”; whereas \acute{e} is regarded as “both immediacy and mediacy, and even as indeterminate.”

$S(x, y)$ begins with the same first logic expression as the category of actuality itself, $\equiv x \neg \check{o} \sim y \neg$. Its second logic expression involves adding \acute{e} to the fourth logic position, leading to an exchange between y and \check{o} in the first and second logic positions. Previously, if we transformed $\equiv x \neg \check{o} \sim y \neg$ into $\equiv \check{o} \neg x \sim y \neg$ using the category of identity, the only symbol that could be definitely transformed was the affirmation \equiv , a pure connection, not “the absolute.” However, when we add \acute{e} to the fourth logic position and place \acute{e} and x in a true affirmation \equiv relationship, x becomes “substantiated” because of the “absoluteness” of the symbol \acute{e} . Yet, this does not mean x is in a “relationship” with another entity \acute{e} , hence not representing its own “absoluteness.” On the contrary, since \acute{e} is a purely logical symbol, connecting it with x through affirmation \equiv does not add any content to x . \acute{e} merely “substantiates” x through its logical significance of “both immediacy and mediacy, and even indeterminacy,” otherwise the only real transformation possible is through affirmation \equiv .

Therefore, the transformation of $\equiv x \neg \check{o} \sim y \neg$ into $\acute{e} \equiv x \neg y \sim \check{o} \neg$ can be interpreted as follows: In the previous logic expression, x was only superficially particularized as immediacy or mediacy. When \acute{e} is added to the fourth logic position, the superficial \check{o} is

abstracted by negation \neg . Now, as x becomes $\dot{\epsilon}$, the universal determination afforded by $\dot{\epsilon}$ to x , and since $\dot{\epsilon}$ is a purely logical symbol, x merely returns to its own substantiality and absoluteness within the affirmation with $\dot{\epsilon}$.

When $\dot{\epsilon} \equiv x \neg \cdot y \sim \check{\circ} \cdot \neg$ transforms into $\dot{\epsilon} \equiv x \neg \cdot y \sim \emptyset \cdot \neg$, we can interpret it as follows: The universality of $\dot{\epsilon}$ should not be understood as “indeterminacy,” hence “self $\dot{\epsilon}$ ” abstracts itself and transforms $\check{\circ}$ into \emptyset , indicating that the self does not come into “undifferentiated unity with x and y .”

When $\dot{\epsilon} \equiv x \neg \cdot y \sim \emptyset \cdot \neg$ transforms into $\dot{\epsilon} \cdot x \equiv \neg \cdot \emptyset \sim \cdot \neg$, we interpret it as: $\dot{\epsilon}$ is abstracted by x , or rather, $\dot{\epsilon}$ concretizes itself within x . “ $\dot{\epsilon} \cdot x$ ” often appears in our conscious perception, very colloquially as “my x .”

When $\dot{\epsilon} \cdot x \equiv \neg \cdot \emptyset \sim \cdot \neg$ transforms into $\dot{\epsilon} \equiv x \neg \cdot \check{\circ} \sim \cdot \neg$, we interpret it as: x leaves my representation and objectifies itself in the third logic position, transforming \emptyset into $\check{\circ}$.

When $\dot{\epsilon} \equiv x \neg \cdot \check{\circ} \sim \cdot \neg$ transforms into $\dot{\epsilon} \equiv x \neg \cdot \sim \check{\circ} \cdot \neg$, we interpret it as: “self $\dot{\epsilon}$ ” re-concretizes, moving $\check{\circ}$ to the first logic position, indicating that “self $\dot{\epsilon}$ ” now identifies with the concrete negation \neg .

Finally, when $\dot{\epsilon} \equiv x \neg \cdot \sim \check{\circ} \cdot \neg$ transforms into $\equiv x \neg \cdot a \sim \check{\circ} \cdot \neg$, we interpret it as: Since “self $\dot{\epsilon}$ ” should be concrete, but now there is a vacancy in the second logic position, the present objective actuality does not match the nature of $\dot{\epsilon}$, thus the self dissipates into a new

objectivity. In this logic expression, the self has experienced the movement of

“ $\dot{\epsilon} \rightarrow \dot{\epsilon} \rightarrow \dot{\epsilon} \cdot x \rightarrow \dot{\epsilon} \rightarrow \dot{\epsilon}$,” thus forming a potential affirmation, which transitions the entire movement of the self into a new determination or item ‘a’.

The free item $\equiv x$ can be generated by the first or last logical formula and should be marked as $\equiv x * S (1)$ or $\equiv x * S (7)$. However, if it is $\equiv x * S (7)$, then the first logical formula can only transform into $\equiv S^* (1)$, and the connection that follows can only be $\equiv \check{\circ} \neg \cdot x \sim y \cdot \neg$. The second logical formula can generate $\dot{\epsilon} \equiv x * S (2)$, and the third logical formula can generate $\dot{\epsilon} \equiv x * S (3)$, but only one can be chosen between $\dot{\epsilon} \equiv x * S (3)$ and $\dot{\epsilon} \equiv x * S (2)$. The same principle applies to the remaining three Substantial Relationships, so I won’t elaborate further.

Note: $S ()$ maintains that x acts as the both being-in-self and being-for-self determinateness of actuality.

Relationship of Substantiality SID1()

SID1(a, x)

$\equiv \check{\circ} \neg \cdot a \sim x \cdot \neg \rightarrow \equiv \check{\circ}$

$\dot{\epsilon} \equiv \check{\circ} \neg \cdot x \sim a \cdot \neg \rightarrow \dot{\epsilon} \equiv$

$\dot{\epsilon} \equiv x \neg \cdot \check{\circ} \sim a \cdot \neg ID \rightarrow \dot{\epsilon} \equiv x$

$\dot{\epsilon} \equiv x \neg \cdot a \sim \check{\circ} \cdot \neg \rightarrow \dot{\epsilon} \equiv x$

$$\dot{\epsilon} \cdot x \equiv \neg \cdot a \sim \emptyset \cdot \neg$$

$$\dot{\epsilon} \cdot \equiv x \neg \cdot \emptyset \sim \cdot \neg$$

$$\dot{\epsilon} \equiv x \neg \cdot \check{\emptyset} \sim \cdot \neg$$

$$\equiv x \neg \cdot \check{\emptyset} \sim b \cdot \neg \rightarrow \equiv x$$

$$(fi) * SID1(): \dot{\epsilon} \equiv * SID1(2) \wedge \dot{\epsilon} \equiv x * SID1(3) \wedge \equiv x * SID1(8)$$

$$(ncfi) * SID1(): \equiv \check{\emptyset} * SID1(1) \wedge \dot{\epsilon} \cdot \equiv x * SID1(4)$$

We now introduce another category within the relationship of substantiality, abbreviated with the symbol SID1() next to the S. More categories, such as SID2(), will be discussed subsequently. The initial formula of this category, $\equiv \check{\emptyset} \neg \cdot a \sim x \cdot \neg$, differs from S () as it positions $\check{\emptyset}$ at the third logic position instead of an item.

As we proceed, I'll clarify the transformations of this category, though many patterns have been previously discussed. Therefore, major differences will be highlighted, with minor ones only briefly mentioned.

When transforming from $\equiv \check{\emptyset} \neg \cdot a \sim x \cdot \neg$ to $\dot{\epsilon} \equiv \check{\emptyset} \neg \cdot x \sim a \cdot \neg$, consider that the truth of $\check{\emptyset}$ is not solely unified with a, as a represents either immediacy or mediacy, thus $\check{\emptyset}$'s truth must encompass the entirety, represented as $\dot{\epsilon}$. $\dot{\epsilon}$ allows for the exchange of positions between x and a, unifying $\check{\emptyset}$ with x.

When $\dot{\epsilon} \equiv \check{\emptyset} \neg \cdot x \sim a \cdot \neg$ transforms to $\dot{\epsilon} \equiv x \neg \cdot \check{\emptyset} \sim a \cdot \neg$ ID, it signifies that $\dot{\epsilon}$ not only maintains

universality but also particularizes itself, thus using ID to swap x with $\check{\circ}$, turning itself into x .

The last formula of $SID1(a, x)$, $\equiv x \neg \check{\circ} \sim b \neg$, structurally mirrors the first formula of $S()$ and the last formula of $SID1()$ matches the structure of the first formula of $S()$. This characteristic will be crucial in the subsequent discussions on dialectical logic symbol reasoning.

In summary, $SID1(a, x)$ transforms the originally first logic position item, here x , into an item that is both being-in-self and being-for-self within the actuality category. Therefore, $SID1()$ maintains ‘necessity’ while providing flexibility absent in the $AC()$ Actuality category itself. Lastly, it’s worth noting that $SID1()$ preserves sequential flexibility, meaning the original first logic position item shifts to the third logic position, becoming both being-in-self and being-for-self.

Relationship of Substantiality $SID2()$

$SID2(x, y)$

$$\equiv x \neg \check{\circ} \sim y \neg \rightarrow \equiv x$$

$$\dot{\equiv} x \neg y \sim \check{\circ} \neg \rightarrow \dot{\equiv} x$$

$$\dot{\equiv} x \neg y \sim \emptyset \neg \rightarrow \dot{\equiv} x$$

$$\dot{\equiv} x \equiv \neg \emptyset \sim \neg$$

$$\dot{\equiv} x \equiv x \neg \check{\circ} \sim \neg$$

$$\dot{\epsilon} \equiv x \neg \cdot \ddot{o} \sim a \cdot \neg \rightarrow \dot{\epsilon} \equiv$$

$$\dot{\epsilon} \equiv \ddot{o} \neg \cdot x \sim a \cdot \neg \text{ID} \rightarrow \dot{\epsilon} \equiv \ddot{o}$$

$$\equiv \ddot{o} \neg \cdot a \sim x \cdot \neg \rightarrow \equiv \ddot{o}$$

$$\text{(fi)*SID2 ()}: \equiv x^* \text{SID2}(1) \wedge \dot{\epsilon} \equiv x^* \text{SID2}(2) \wedge \dot{\epsilon} \equiv \text{*SID2}(6) \wedge \dot{\epsilon} \equiv \ddot{o}^* \text{SID1}(7) \wedge \equiv \ddot{o}^* \text{SID1}(8)$$

$$\text{(ncfi)*SID2 ()}: \dot{\epsilon} \cdot \equiv x^* \text{SID1}(3)$$

SID2(x, y) utilizes the ID in the penultimate formula, $\dot{\epsilon} \equiv \ddot{o} \neg \cdot x \sim a \cdot \neg \text{ID}$. Broadly, SID2(x, y) affects the structure of actuality by downgrading the initially being-in-self and being-for-self item, here x, to the first logic position, allowing a new item, here a, to become within the being-in-self and being-for-self category of actuality. Notably, SID2() offers the ‘necessity’ and ‘flexibility’ sought in actuality without the sequential flexibility of SID1(), as it positions a new item as independent and defined.

Relationship of Substantiality S2()

$$S2(x, a)$$

$$\equiv x \neg \cdot a \sim \ddot{o} \cdot \neg \rightarrow \equiv x$$

$$\dot{\epsilon} \equiv x \neg \cdot \ddot{o} \sim a \cdot \neg \rightarrow \dot{\epsilon} \equiv x$$

$$\dot{\epsilon} \cdot \equiv x \neg \cdot \emptyset \sim a \cdot \neg \rightarrow \dot{\epsilon} \cdot \equiv x$$

$$\dot{\epsilon} \cdot x \equiv \neg \cdot \sim \emptyset \cdot \neg$$

$$\dot{\epsilon} \cdot \equiv x \neg \cdot \sim \ddot{o} \cdot \neg$$

$$\dot{\epsilon} \equiv x \neg \check{\circ} \sim \neg$$

$$\equiv x \neg \check{\circ} \sim d \neg \rightarrow \equiv x$$

$$(fi)*S2 () : [\equiv x*S2 (1) \vee \equiv x*S2 (7)] \wedge \dot{\epsilon} \equiv x*S2 (2)$$

$$(ncfi) *S2 () : \dot{\epsilon} \equiv x*S2 (3)$$

S2() stands distinct from the previous three categories of substantiality: its first formula aligns with the initial formula of the CONT () category of contingency.

Category of Causality

CAUS (x, y)

$$\equiv x \neg \check{\circ} \sim y \neg \rightarrow \equiv x$$

$$\dot{\epsilon} \equiv x \neg y \sim \check{\circ} \neg$$

$$x \equiv \dot{\epsilon} \neg y \sim \check{\circ} \neg \text{ ID } \rightarrow x \equiv$$

$$x \equiv y \neg \dot{\epsilon} \sim \check{\circ} \neg \text{ ID } \rightarrow x \equiv y$$

$$x \cdot \equiv y \neg \check{\circ} \sim \dot{\epsilon} \neg$$

$$x \cdot y \equiv \neg \check{\circ} \sim \dot{\epsilon} \neg$$

$$x \cdot \equiv y \neg \sim \check{\circ} \neg$$

$$x \equiv y \neg \sim \check{\circ} \neg$$

$$x \equiv \neg y \sim \check{\circ} \neg \text{ ID}$$

$$\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg \rightarrow \equiv x$$

$$(fi)*CAUS(): x \equiv *CAUS(3) [\equiv x *CAUS(1) \vee \equiv x *CAUS(10)] \wedge x \equiv y *CAUS(4)$$

$$(ncfi)CAUS(): x \cdot \equiv y *CAUS(5)$$

Now, let's introduce the category of causality, abbreviated as CAUS (). For our thinking, causality generally represents a kind of "objective necessity." Therefore, the design of the logical operation formulas in the category of causality is intended to allow the subjective "self $\check{\epsilon}$ " to dissolve itself in the face of objective actuality. To achieve this, the thinking uses the technique of free interchangeability between the third and fourth logical positions.

When $\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg$ transforms into $\check{\epsilon} \equiv x \neg \cdot y \sim \check{\circ} \cdot \neg$, the self $\check{\epsilon}$ acts on the actuality structure similar to the first two formulas of S(), but at this time, thinking determines that x is the true "entity," hence it uses ID on $\check{\epsilon} \equiv x \neg \cdot y \sim \check{\circ} \cdot \neg$ to transform it into $x \equiv \check{\epsilon} \neg \cdot y \sim \check{\circ} \cdot \neg$. Then, thinking again applies ID to $x \equiv \check{\epsilon} \neg \cdot y \sim \check{\circ} \cdot \neg$, transforming it into $x \equiv y \neg \cdot \check{\epsilon} \sim \check{\circ} \cdot \neg$.

The operations of the subsequent logical formulas are all aimed at making the last logical formula return to the formula identical to the first logical formula $\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg$. We can see that the five logical transformations of $x \equiv y \neg \cdot \check{\epsilon} \sim \check{\circ} \cdot \neg$, $x \cdot \equiv y \neg \cdot \check{\circ} \sim \check{\epsilon} \cdot \neg$, $x \cdot y \equiv \neg \cdot \check{\circ} \sim \check{\epsilon} \cdot \neg$, $x \cdot \equiv y \neg \cdot \sim \check{\circ} \cdot \neg$, and $x \equiv y \neg \cdot \sim \check{\circ} \cdot \neg$ express the self-return movement of x at the level of being both being-in-itself and being-for-itself as " $x \rightarrow x \cdot \rightarrow x \cdot y \rightarrow x \cdot \rightarrow x$," and during " $x \cdot \rightarrow x \cdot y \rightarrow x$," the "self $\check{\epsilon}$," abstracted by negation, is eliminated. In common language, "x·y" can be

expressed as “the effect of x is y.”

In summary, in CAUS (x, y), we can consider x as the “cause” and y as the “effect.”

Conversely, CAUS (y, x) sets y as the “cause” and x as the “effect.”

When the following logical formulas of actualities occur between CAUS (x, y) and CAUS (y, x), x and y become “mutually causal”:

MUCAC (x, y)

$\equiv y \cdot \check{\text{O}} \sim x \cdot \neg \rightarrow \equiv$

$\equiv \check{\text{O}} \cdot y \sim x \cdot \neg \text{ID} \rightarrow \equiv \check{\text{O}}$

$\equiv \check{\text{O}} \cdot x \sim y \cdot \neg \text{TIF} \rightarrow \equiv \check{\text{O}}$

$\equiv x \cdot \check{\text{O}} \sim y \cdot \neg \text{ID} \rightarrow \equiv$

(fi)* MUCAC (x, y): $\equiv \text{*MUCAC (1)} \vee \equiv \text{*MUCAC (3)}$

(ncfi)* MUCAC (x, y): $\equiv \check{\text{O}} \cdot \text{*MUCAC (2)} \vee \equiv x \cdot \text{*MUCAC (4)}$

When x and y are thus set as “mutually causal” as above, thinking must enter the final category of the doctrine of essence, the category of reciprocity, abbreviated as RECI ().

The Category of Reciprocity

RECI (x, y)

$\dot{\varepsilon} \equiv x \neg \cdot \check{\circ} \sim y \cdot \neg$

$\dot{\varepsilon} \equiv \check{\circ} \neg \cdot x \sim y \cdot \neg$ ID

$\dot{\varepsilon} \equiv \check{\circ} \neg \cdot y \sim x \cdot \neg$

$\dot{\varepsilon} \equiv y \neg \cdot \check{\circ} \sim x \cdot \neg$ ID

$\dot{\varepsilon} \equiv y \neg \cdot x \sim \check{\circ} \cdot \neg$

$\dot{\varepsilon} \equiv x \neg \cdot y \sim \check{\circ} \cdot \neg$ ID

$\dot{\varepsilon} \equiv x \neg \cdot \check{\circ} \sim y \cdot \neg$

$\dot{\varepsilon} \equiv \check{\circ} \neg \cdot x \sim y \cdot \neg$ ID

$\dot{\varepsilon} \equiv \check{\circ} \neg \cdot y \sim x \cdot \neg$

$\dot{\varepsilon} \equiv y \neg \cdot \check{\circ} \sim x \cdot \neg$ ID

$\dot{\varepsilon} \equiv y \neg \cdot x \sim \check{\circ} \cdot \neg$

$\dot{\varepsilon} \equiv x \neg \cdot y \sim \check{\circ} \cdot \neg$ ID

$\dot{\varepsilon} \equiv x \neg \cdot \check{\circ} \sim y \cdot \neg$

(fi)* RECI (): $\dot{\varepsilon} \equiv *RECI () \wedge \dot{\varepsilon} \equiv *RECI ()$

In the category of reciprocity, RECI (), the items x and y represent the most purified particularizations within the dialectical logic symbol system. This purity is because the interchanges of x, y, and $\check{\circ}$ across the first three logic positions in RECI (x, y) do not

transform $\check{\circ}$ into \emptyset . Instead, these movements synchronize perfectly with the "self $\dot{\epsilon}$ " undergoing pure cyclic movements of " $\dot{\epsilon} \rightarrow \dot{\epsilon} \rightarrow \dot{\epsilon}$." The interactions within RECI () are entirely cyclic and do not require the insertion of any items, thus lacking a definitive start or end point.

Using an example from Buddhist doctrine, where the third and fourth stages of the twelve links of dependent origination involve "consciousness" and "name-and-form" reciprocally defining each other, we define x as "consciousness" and y as "name-and-form"⁴ for this illustration.

In this setup, $\dot{\epsilon} \equiv x \cdot \check{\circ} \sim y \cdot \neg$ can be interpreted as "the self is consciousness," with x maintaining a mediated particularization with $\check{\circ}$, and y combining the first and second logic positions as an immediate particularization. As the self becomes concrete in the form of $\dot{\epsilon}$, $y \cdot \neg$ asserts a "negated certainty." However, the self $\dot{\epsilon}$ should represent the totality of both immediate and mediate determinateness, thus driving the transformation to the next logic formula.

The second logic formula, $\dot{\epsilon} \equiv \check{\circ} \cdot x \sim y \cdot \neg$ ID, can be interpreted as "the self is pure being, returning to its totality," with x and y now unified in the first two logic positions, making the

⁴ In the twelve links of dependent origination in Buddhism, 'consciousness' and 'name-and-form' are key concepts of reciprocity.

allocation of immediacy or mediacy between x and y indeterminate. The transformation from $x \neg \check{\circ}$ to $\check{\circ} \neg x$ completes the shift from particularity to universality, resulting in a self that purely "is," rather than being merely a particular x or a universal $\check{\circ}$.

In the third logic formula, $\acute{\epsilon} \equiv \check{\circ} \neg y \sim x \neg$, this can be interpreted as the self, in its totality, should not be seen as "undetermined." Thus, $\acute{\epsilon}$, while inherently a pure logical setting, also allows us as users of these symbols to insert "the self" into systematic operations, representing "anything," potentially our sensory data. Therefore, the self abstracts itself to avoid contact with this reality and swaps positions between x and y , creating new possibilities. The notation $\acute{\epsilon} \cdot$ and $x \cdot$ indicates both the self and x are in an abstract state, implying they "exist mutually within each other."

In the fourth logic formula, $\acute{\epsilon} \equiv y \neg \check{\circ} \sim x \neg$ ID, thought explicitly reveals that the determinateness which $\acute{\epsilon}$ desires to distance itself from is y , thus abstracting itself into $\acute{\epsilon} \cdot$ and no longer functioning in thought. On the contrary, y undergoes the completion of the transformation from ' $\check{\circ} \neg y$ to $y \neg \check{\circ}$ ', thus finishing its process of concretization. Now, the thought is functioning through the pure affirmation \equiv , formed by the interaction between $y \neg \check{\circ}$ and $\check{\circ} \neg y$. However, true actuality is not such an abstract self, and pure affirmation \equiv must connect to a concrete item, progressing the thought to the next logic transformation.

By the fifth and sixth logic formulas, $\dot{\epsilon} \equiv y \neg \cdot x \sim \check{\epsilon} \cdot \neg$ and $\dot{\epsilon} \equiv x \neg \cdot y \sim \check{\epsilon} \cdot \neg$ ID, the self $\dot{\epsilon}$ realizes that what limited the particularization of self was the logical symbol $\check{\epsilon}$. Now, recognizing that only the self $\dot{\epsilon}$ as a pure logical symbol and an infinite particularization truly exists, it concretizes y , previously unified with x , into the concrete form of the self $\dot{\epsilon}$. The self becomes the complete particularity unified from x and y in the second and third logic positions. Notably, in the sixth formula, the self $\dot{\epsilon}$ again forms an affirmative relationship with x , but this time the relationship with x must manifest through y .

Upon reaching the fifth and sixth formulas, the self $\dot{\epsilon}$ completes a “ $\dot{\epsilon} \rightarrow \dot{\epsilon} \cdot \rightarrow \dot{\epsilon}$ ” movement. Before this self-returning movement, x and the self were in an “affirmative” and “mutually existing” relationship, whereas y and the self were in a “different” and “negated certainty” relationship. Now, as x and y unify in the fifth and sixth logic transformations, and since y , in its concretized form, previously underwent “negated certainty” and abstracted negation, now attains substantiality through an affirmative relationship with the self, thus becoming the primary particularity.

The seventh and eighth logic formulas, $\dot{\epsilon} \cdot \equiv x \neg \cdot \check{\epsilon} \sim y \cdot \neg$ and $\dot{\epsilon} \cdot \equiv \check{\epsilon} \neg \cdot x \sim y \cdot \neg$ ID, transition into a relationship where x and the self are “different,” while $\dot{\epsilon} \cdot$ and $y \cdot$ remain abstract, indicating they “exist mutually within each other.” This suggests that the self $\dot{\epsilon}$ is not any particular aspect of the unified consciousness and name-form but purely a symbol.

By the ninth and tenth formulas, $\dot{\epsilon} \equiv \ddot{o} \cdot y \sim x \cdot \neg$ and $\dot{\epsilon} \equiv y \cdot \neg \ddot{o} \sim x \cdot \neg$ ID, ‘the self $\dot{\epsilon}$ is name-form y,’ and x, abstracted by negation, embodies the self involving ‘negated certainty.’

The eleventh and twelfth formulas, $\dot{\epsilon} \cdot \equiv y \cdot \neg x \sim \ddot{o} \cdot \neg$ and $\dot{\epsilon} \cdot \equiv x \cdot \neg y \sim \ddot{o} \cdot \neg$ ID, echo the fifth and sixth in unifying x and y in the third and fourth logic positions. However, this time the self abstracts into $\dot{\epsilon} \cdot$, avoiding contact with this consciousness and name-form unity, and $\dot{\epsilon} \cdot$ and $\ddot{o} \cdot$ remain in the same abstract state, denoting the self $\dot{\epsilon}$ and being \ddot{o} ‘exist mutually within each other.’ This indicates the self $\dot{\epsilon}$ does not particularize itself as any aspect of this unified consciousness and name-form but remains a purely symbolic entity.

When the category of reciprocity develops to $\dot{\epsilon} \cdot \equiv x \cdot \neg y \sim \ddot{o} \cdot \neg$ ID, it reaches the final formula, because the next formula returns to the starting point of $\dot{\epsilon} \equiv x \cdot \neg \ddot{o} \sim y \cdot \neg$. This cycle of twelve formulas embodies the pure circularity of thought in the category of reciprocity. The cyclical category allows any mid-sequence logic formula to serve as the starting point, constructing RECI (y, x) reversely, achieving the same outcome.

In conclusion, the category of reciprocity, which Hegel describes as ‘the nullity of distinctions’ (Hegel, 1975, p. 218), embodies ‘the notion’—a concept representing truth and freedom. This leads us into ‘the category of the notion.’

The Category of The Notion

$N(\dot{\epsilon})$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \check{\circ} \neg \cdot x \sim y \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \neg \cdot x \sim y \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \neg \cdot y \sim x \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \neg \cdot x \sim y \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \neg \cdot \emptyset \sim y \cdot \neg \text{ ID}$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \dot{\epsilon} \neg \cdot \sim \emptyset \cdot \neg$$

$$\dot{\epsilon} \equiv \dot{\epsilon} \neg \cdot \sim \check{\circ} \cdot \neg$$

$$\dot{\epsilon} \equiv \dot{\epsilon} \cdot \neg \cdot \check{\circ} \sim \cdot \neg$$

$$\dot{\epsilon} \equiv \dot{\epsilon} \neg \cdot \emptyset \sim \cdot \neg$$

$$\dot{\epsilon} \equiv \emptyset \neg \cdot \sim \cdot \neg \text{ ID}$$

$$\dot{\epsilon} \neg \cdot \check{\circ} \sim \cdot \neg$$

In the category of the Notion $N(\dot{\epsilon})$, by the second logical formula, $\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \neg \cdot x \sim y \cdot \neg$, the mind introduces a second ‘self $\dot{\epsilon}$ ’, which represents ‘the Notion’. Adding this second $\dot{\epsilon}$ transforms the $\check{\circ}$ at the third logic position into \emptyset . $\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \neg \cdot x \sim y \cdot \neg$ can infinitely cycle with the third logical formula, $\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \neg \cdot y \sim x \cdot \neg$, because the fourth logic position is occupied by ‘ $\dot{\epsilon} \cdot \dot{\epsilon}$ ’. Thus, whether it is x or y in the second logic position, any unity with \emptyset in the third logic position cannot be distinctly recognized as either $\dot{\epsilon}$ or $\dot{\epsilon}$. Once any determinateness unites

with \emptyset and is identified as $\dot{\epsilon}$, it immediately becomes $\dot{\epsilon}$, leading to a swap of determinateness in the first and second logic positions.

Upon reaching the fifth logical formula, $\dot{\epsilon} \cdot \dot{\epsilon} \equiv \neg \cdot \emptyset \sim y \cdot \neg$ ID, the mind employs ID to eliminate x , indicating that any determinateness aiming to transcend the ‘negative certainty’ of the first and second logic positions is immediately annihilated. By the sixth formula, $\dot{\epsilon} \cdot \equiv \dot{\epsilon} \cdot \neg \sim \emptyset \cdot \neg$, all determinateness is erased. The next three formulas culminate in ‘self $\dot{\epsilon}$ ’ being absolutely unified with nothingness, and then the mind uses ID to remove one notion of self, transitioning from existence $\dot{\epsilon}$ to nothing \emptyset . The formula becomes $\dot{\epsilon} \cdot \neg \cdot \check{\emptyset} \sim \cdot \neg$, placing ‘self $\dot{\epsilon}$ ’ within the category of identity ID($\dot{\epsilon}$).”

Transition to Categories in the Doctrine of Being

DB($\dot{\epsilon}$)

$\neg \cdot \dot{\epsilon} \sim \check{\emptyset} \cdot \neg$

$\neg \cdot \check{\emptyset} \sim \dot{\epsilon} \cdot \neg$

$\neg \cdot \emptyset \sim \dot{\epsilon} \cdot \neg$

$\neg \cdot \sim \emptyset \cdot \neg$

BC

$\neg \cdot \sim \check{\emptyset} \cdot \neg$

$\neg \cdot \check{\emptyset} \sim \cdot \neg$

$$\neg \cdot \emptyset \sim \cdot \neg$$

$$\neg \cdot \sim \emptyset \cdot \neg$$

$$\neg \cdot \sim \check{\emptyset} \cdot \neg$$

Following the identity ID($\dot{\epsilon}$), the logic progresses to DB($\dot{\epsilon}$) and BC, thus returning the dialectic of The Notion back to The Doctrine of Being.

Buddhist Notion of The Category

NB($\dot{\epsilon}$)

$$\dot{\epsilon} \cdot \equiv \check{\emptyset} \cdot \neg \cdot x \sim y \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \cdot \neg \cdot x \sim y \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \cdot \neg \cdot y \sim x \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \emptyset \cdot \neg \cdot x \sim y \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \equiv \neg \cdot \emptyset \sim y \cdot \neg \text{ ID}$$

$$\dot{\epsilon} \cdot \equiv \dot{\epsilon} \cdot \neg \cdot \sim \emptyset \cdot \neg$$

$$\dot{\epsilon} \cdot \dot{\epsilon} \cdot \neg \cdot \sim \check{\emptyset} \cdot \neg$$

$$\dot{\epsilon} \cdot \neg \cdot \dot{\epsilon} \sim \check{\emptyset} \cdot \neg \text{ ID}$$

$$\dot{\epsilon} \cdot \neg \cdot \check{\emptyset} \sim \dot{\epsilon} \cdot \neg$$

$\check{\text{ö}}\neg\cdot\dot{\text{é}}\sim\dot{\text{é}}\cdot\neg$ ID

$\check{\text{ö}}\equiv$

\emptyset

In the Buddhist category of ‘The Notion’ NB($\dot{\text{é}}$), by the sixth logical formula $\dot{\text{é}}\cdot\equiv$
 $\dot{\text{é}}\neg\cdot\sim\emptyset\cdot\neg$, the affirmation \equiv is removed, transforming the formula into $\dot{\text{é}}\cdot\dot{\text{é}}\neg\cdot\sim\check{\text{ö}}\cdot\neg$.

Buddhism does not perceive ‘self’ as possessing an independent being-for-self existence.

Further computations lead to an important formula, $\check{\text{ö}}\neg\cdot\dot{\text{é}}\sim\dot{\text{é}}\cdot\neg$ ID, reflecting a frequently cited

Buddhist scripture: ‘All aggregates are impermanent and suffering, thus “not self, nor

belonging to self.’”, The structure $\neg\cdot\dot{\text{é}}\sim\dot{\text{é}}\cdot\neg$ represents the most solid affirmation in the

dialectical logic symbol system because it signifies ‘the mutual abstraction and concretization

of $\dot{\text{é}}$ and \neg ’. The final conversion of the formula from $\check{\text{ö}}\equiv$ to \emptyset indicates a cessation of

thinking as the mind transitions into nothing.

The Classical Continuum

Before entering the inference phase, we must first understand some continuums, among

which the continuums of the categories of substantial relationships can be infinite.

The Cyclical Continuum of S () and S2 ()

S (x, y)

$$\equiv x \neg \check{\circ} \sim y \neg \rightarrow \equiv x$$

$$\dot{\epsilon} \equiv x \neg y \sim \check{\circ} \neg \rightarrow \dot{\epsilon} \equiv x$$

$$\dot{\epsilon} \cdot \equiv x \neg y \sim \emptyset \neg \rightarrow \dot{\epsilon} \cdot \equiv x$$

$$\dot{\epsilon} \cdot x \equiv \neg \emptyset \sim \neg$$

$$\dot{\epsilon} \cdot \equiv x \neg \check{\circ} \sim \neg$$

$$\dot{\epsilon} \equiv x \neg \sim \check{\circ} \neg$$

$$\equiv x \neg a \sim \check{\circ} \neg \rightarrow \equiv x$$

$$(fi)*S () : [\equiv x*S (1) \vee \equiv x*S2 (7)] \wedge \dot{\epsilon} \equiv x*S (2)$$

$$(ncfi)*S () : \dot{\epsilon} \cdot \equiv x*S (3)$$

S2(x, c)

$$\equiv x \neg a \sim \check{\circ} \neg \rightarrow \equiv x$$

$$\dot{\epsilon} \equiv x \neg \check{\circ} \sim a \neg \rightarrow \dot{\epsilon} \equiv x$$

$$\dot{\varepsilon} \equiv x \neg \cdot \emptyset \sim a \cdot \neg \rightarrow \dot{\varepsilon} \equiv x$$

$$\dot{\varepsilon} \cdot x \equiv \neg \cdot \sim \emptyset \cdot \neg$$

$$\dot{\varepsilon} \equiv x \neg \cdot \sim \check{\emptyset} \cdot \neg$$

$$\dot{\varepsilon} \equiv x \neg \cdot \check{\emptyset} \sim \cdot \neg$$

$$\equiv x \neg \cdot \check{\emptyset} \sim d \cdot \neg \rightarrow \equiv x$$

$$(fi)*S2 () : [\equiv x *S2 (1) \vee \equiv x *S2 (7)] \wedge \dot{\varepsilon} \equiv x *S2 (2)$$

$$(ncfi) *S2 () : \dot{\varepsilon} \equiv x *S2 (3)$$

S (x, d)

$$\equiv x \neg \cdot \check{\emptyset} \sim d \cdot \neg \rightarrow \equiv x$$

$$\dot{\varepsilon} \equiv x \neg \cdot d \sim \check{\emptyset} \cdot \neg \rightarrow \dot{\varepsilon} \equiv x$$

$$\dot{\varepsilon} \equiv x \neg \cdot d \sim \emptyset \cdot \neg \rightarrow \dot{\varepsilon} \equiv x$$

$$\dot{\varepsilon} \cdot x \equiv \neg \cdot \emptyset \sim \cdot \neg$$

$$\dot{\varepsilon} \equiv x \neg \cdot \check{\emptyset} \sim \cdot \neg$$

$$\dot{\varepsilon} \equiv x \neg \cdot \sim \check{\emptyset} \cdot \neg$$

$$\equiv x \neg \cdot e \sim \check{o} \cdot \neg \rightarrow \equiv x$$

$$(fi)*S () : [\equiv x *S (1) \vee \equiv x *S2 (7)] \wedge \dot{\epsilon} \equiv x *S (2)$$

$$(ncfi) *S () : \dot{\epsilon} \cdot \equiv x *S (3)$$

S2 (x, e)

$$\equiv x \neg \cdot e \sim \check{o} \cdot \neg \rightarrow \equiv x$$

$$\dot{\epsilon} \equiv x \neg \cdot \check{o} \sim e \cdot \neg \rightarrow \dot{\epsilon} \equiv x$$

$$\dot{\epsilon} \cdot \equiv x \neg \cdot \emptyset \sim e \cdot \neg \rightarrow \dot{\epsilon} \cdot \equiv x$$

$$\dot{\epsilon} \cdot x \equiv \neg \cdot \sim \emptyset \cdot \neg$$

$$\dot{\epsilon} \cdot \equiv x \neg \cdot \sim \check{o} \cdot \neg$$

$$\dot{\epsilon} \equiv x \neg \cdot \check{o} \sim \cdot \neg$$

$$\equiv x \neg \cdot \check{o} \sim f \cdot \neg \rightarrow \equiv x$$

$$(fi)*S2 () : [\equiv x *S2 (1) \vee \equiv x *S2 (7)] \wedge \dot{\epsilon} \equiv x *S2 (2)$$

$$(ncfi) *S2 () : \dot{\epsilon} \cdot \equiv x *S2 (3)$$

The Cyclical Continuum of SID1 () and SID2 ()

SID1(x, y)

$$\equiv \check{\circ} \neg \cdot y \sim x \cdot \neg \rightarrow \equiv \check{\circ}$$

$$\dot{\epsilon} \equiv \check{\circ} \neg \cdot x \sim y \cdot \neg \rightarrow \dot{\epsilon} \equiv$$

$$\dot{\epsilon} \equiv x \neg \cdot \check{\circ} \sim y \cdot \neg \text{ID} \rightarrow \dot{\epsilon} \equiv x$$

$$\dot{\epsilon} \cdot \equiv x \neg \cdot y \sim \check{\circ} \cdot \neg \rightarrow \dot{\epsilon} \cdot \equiv x$$

$$\dot{\epsilon} \cdot x \equiv \neg \cdot y \sim \emptyset \cdot \neg$$

$$\dot{\epsilon} \cdot \equiv x \neg \cdot \emptyset \sim \cdot \neg$$

$$\dot{\epsilon} \equiv x \neg \cdot \check{\circ} \sim \cdot \neg$$

$$\equiv x \neg \cdot \check{\circ} \sim a \cdot \neg \rightarrow \equiv x$$

$$(fi) * \text{SID1}(): \dot{\epsilon} \equiv * \text{SID1}(2) \wedge \dot{\epsilon} \equiv x * \text{SID1}(3) \wedge \equiv x * \text{SID1}(8)$$

$$(ncfi) * \text{SID1}(): \equiv \check{\circ} * \text{SID1}(1) \wedge \dot{\epsilon} \cdot \equiv x * \text{SID1}(4)$$

SID2(x, a)

$$\equiv x \neg \cdot \check{\circ} \sim a \cdot \neg \rightarrow \equiv x$$

$$\dot{\varepsilon} \equiv x \neg \cdot y \sim \check{o} \neg \rightarrow \dot{\varepsilon} \equiv x$$

$$\dot{\varepsilon} \equiv x \neg \cdot y \sim \emptyset \neg \rightarrow \dot{\varepsilon} \equiv x$$

$$\dot{\varepsilon} \cdot x \equiv \neg \cdot \emptyset \sim \neg$$

$$\dot{\varepsilon} \equiv x \neg \cdot \check{o} \sim \neg$$

$$\dot{\varepsilon} \equiv x \neg \cdot \check{o} \sim a \neg \rightarrow \dot{\varepsilon} \equiv$$

$$\dot{\varepsilon} \equiv \check{o} \neg \cdot x \sim a \neg \text{ID} \rightarrow \dot{\varepsilon} \equiv \check{o}$$

$$\equiv \check{o} \neg \cdot a \sim x \neg \rightarrow \equiv \check{o}$$

$$(fi)*\text{SID2 } (): \equiv x * \text{SID2}(1) \wedge \dot{\varepsilon} \equiv x * \text{SID2}(2) \wedge \dot{\varepsilon} \equiv * \text{SID2}(6) \wedge \dot{\varepsilon} \equiv \check{o} * \text{SID1}(7) \wedge \equiv \check{o} * \text{SID1}(8)$$

$$(ncfi)*\text{SID2 } (): \dot{\varepsilon} \equiv x * \text{SID1}(3)$$

S (x, y) — SID1 (x, y): SSID1()

$$\equiv x \neg \cdot \check{o} \sim y \neg \rightarrow \equiv$$

$$\equiv \check{o} \neg \cdot x \sim y \neg \text{ID} \rightarrow \equiv \check{o}$$

$$\equiv \check{o} \neg \cdot x \sim y \neg \text{TIF}$$

$$\equiv \check{o} \neg \cdot y \sim x \neg \text{ID}$$

(fi)*SSID1 () : \equiv *SSID1(1) \wedge \equiv $\check{\circ}$ * SSID1(2)

The Influence of SSID1() on S (x, y)

Post S (x, y) by S (x, y) - SID1 (x, y)

$\equiv x \neg \check{\circ} \sim y \neg \rightarrow \equiv$

$\dot{\epsilon} \equiv x \neg y \sim \check{\circ} \neg \rightarrow \dot{\epsilon} \equiv x$

$\dot{\epsilon} \equiv x \neg y \sim \emptyset \neg \rightarrow \dot{\epsilon} \equiv x$

$\dot{\epsilon} \cdot x \equiv \neg \emptyset \sim \neg$

$\dot{\epsilon} \equiv x \neg \check{\circ} \sim \neg$

$\dot{\epsilon} \equiv x \neg \sim \check{\circ} \neg$

$\equiv x \neg a \sim \check{\circ} \neg \rightarrow \equiv x$

(fi)*S (x, y) : \equiv *S (1) \wedge $\dot{\epsilon} \equiv x$ *S (2) \wedge $\equiv x$ *S (7)

(ncfi)*S (x, y) : $\dot{\epsilon} \equiv x$ *S (3)

AC () — \emptyset

$$\equiv x \cdot \check{\circ} \sim y \cdot \neg \rightarrow \equiv$$

$$\equiv \check{\circ} \cdot x \sim y \cdot \neg \text{ID} \rightarrow \equiv \check{\circ}$$

$$\equiv \check{\circ} \leftarrow$$

\emptyset

Here, however, there is a possibility of leading to nothing \emptyset .

The Continuum of Causality

MUCAC ()

$$\equiv y \cdot \check{\circ} \sim x \cdot \neg \rightarrow \equiv$$

$$\equiv \check{\circ} \cdot y \sim x \cdot \neg \text{ID} \rightarrow \equiv \check{\circ}$$

$$\equiv \check{\circ} \cdot x \sim y \cdot \neg \text{TIF} \rightarrow \equiv \check{\circ}$$

$$\equiv x \cdot \check{\circ} \sim y \cdot \neg \text{ID} \rightarrow \equiv$$

or

MUCAC (x, y)

$$\equiv y \cdot \check{\circ} \sim x \cdot \neg \rightarrow \equiv$$

$$\equiv \check{\circ} \cdot y \sim x \cdot \neg \text{ID} \rightarrow \equiv \check{\circ}$$

$$\equiv \check{\circ} \cdot x \sim y \cdot \neg \text{TIF} \rightarrow \equiv$$

$$\equiv x \cdot \check{\circ} \sim y \cdot \neg \text{ID} \rightarrow \equiv x$$

or

MUCAC (y, x)

$$\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg$$

$$\equiv \check{\circ} \neg \cdot x \sim y \cdot \neg \text{ID} \rightarrow \equiv \check{\circ}$$

$$\equiv \check{\circ} \neg \cdot y \sim x \cdot \neg \text{TIF} \rightarrow \equiv$$

$$\equiv y \neg \cdot \check{\circ} \sim x \cdot \neg \text{ID} \rightarrow \equiv y$$

The three types of causal relationships and their connection modes to actualities are very important because the application of ID or TIF to either $\equiv \check{\circ} \neg \cdot y \sim x \cdot \neg$ or $\equiv \check{\circ} \neg \cdot x \sim y \cdot \neg$ affects which formula transforms into $\equiv \check{\circ}$. This impacts the direction of causality. If both are possible, then the two are mutually causal.

MUCAC (x, y)

$$\equiv y \neg \cdot \check{\circ} \sim x \cdot \neg \rightarrow \equiv$$

$$\equiv \check{\circ} \neg \cdot y \sim x \cdot \neg \text{ID} \rightarrow \equiv \check{\circ}$$

$$\equiv \check{\circ} \neg \cdot x \sim y \cdot \neg \text{TIF} \rightarrow \equiv$$

$$\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg \text{ID} \rightarrow \equiv x$$

(fi)* MUCAC (x, y): $\equiv^* \text{MUCAC (1)} \vee \equiv^* \text{MUCAC (3)}$

(ncfi)* MUCAC (x, y): $\equiv \check{\circ}^* \text{MUCAC (2)} \vee \equiv x^* \text{MUCAC (4)}$

CAUS (x, y)

$$\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg \rightarrow \equiv x$$

$$\dot{\equiv} x \neg \cdot y \sim \check{\circ} \cdot \neg$$

$$x \equiv \dot{\equiv} \neg \cdot y \sim \check{\circ} \cdot \neg \text{ID} \rightarrow x \equiv$$

$$x \equiv y \neg \cdot \dot{\equiv} \sim \check{\circ} \cdot \neg \text{ID} \rightarrow x \equiv y$$

$$x \cdot \equiv y \neg \cdot \check{\circ} \sim \dot{\equiv} \cdot \neg \rightarrow x \cdot \equiv y$$

$$x \cdot y \equiv \neg \cdot \emptyset \sim \dot{\epsilon} \cdot \neg$$

$$x \equiv y \neg \cdot \sim \emptyset \cdot \neg$$

$$x \equiv y \neg \cdot \sim \dot{\emptyset} \cdot \neg$$

$$x \equiv \neg \cdot y \sim \dot{\emptyset} \cdot \neg \text{ ID}$$

$$\equiv x \neg \cdot \dot{\emptyset} \sim y \cdot \neg \rightarrow \equiv x$$

$$(fi)*CAUS(): x \equiv *CAUS(3) [\equiv x *CAUS(1) \vee \equiv x *CAUS(10)] \wedge x \equiv y *CAUS(4)$$

$$(ncfi)CAUS(): x \cdot \equiv y *CAUS(5)$$

The above illustrates x as the cause.

MUCAC (y, x)

$$\equiv x \neg \cdot \dot{\emptyset} \sim y \cdot \neg$$

$$\equiv \dot{\emptyset} \neg \cdot x \sim y \cdot \neg \text{ ID} \rightarrow \equiv \dot{\emptyset}$$

$$\equiv \dot{\emptyset} \neg \cdot y \sim x \cdot \neg \text{ TIF} \rightarrow \equiv$$

$$\equiv y \neg \cdot \dot{\emptyset} \sim x \cdot \neg \text{ ID} \rightarrow \equiv y$$

$$(fi)* MUCAC(): \equiv *MUCAC(1) \vee \equiv *MUCAC(3)$$

$$(ncfi)* MUCAC(): \equiv \dot{\emptyset} * MUCAC(2) \vee \equiv y * MUCAC(4)$$

CAUS (y, x)

$$\equiv y \neg \cdot \dot{\emptyset} \sim x \cdot \neg \rightarrow \equiv y$$

$$\dot{\epsilon} \equiv y \neg \cdot x \sim \dot{\emptyset} \cdot \neg \rightarrow \equiv$$

$$y \equiv \dot{\epsilon} \neg \cdot x \sim \dot{\emptyset} \cdot \neg \text{ ID} \rightarrow y \equiv$$

$$y \equiv x \neg \cdot \dot{\epsilon} \sim \dot{\emptyset} \cdot \neg \text{ ID} \rightarrow y \equiv x$$

$$y \cdot \equiv x \neg \cdot \dot{\emptyset} \sim \dot{\epsilon} \cdot \neg \rightarrow y \cdot \equiv x$$

$$y \cdot x \equiv \neg \cdot \emptyset \sim \dot{\epsilon} \cdot \neg$$

$$y \cdot \equiv x \neg \sim \emptyset \neg$$

$$y \equiv x \neg \sim \check{\emptyset} \neg$$

$$y \equiv \neg x \sim \check{\emptyset} \neg \text{ ID}$$

$$\equiv y \neg \check{\emptyset} \sim x \neg \rightarrow \equiv y$$

$$\text{(fi)CAUS } (): y \equiv * \text{CAUS } (3) \wedge [\equiv y \text{CAUS } (1) \vee \equiv y \text{CAUS } (10)] \wedge y \equiv x \text{CAUS } (4)$$

$$\text{(ncfi)CAUS } (): y \cdot \equiv x * \text{CAUS } (5)$$

The above illustrates y as the cause.

MUCAC ()

$$\equiv y \neg \check{\emptyset} \sim x \neg \rightarrow \equiv$$

$$\equiv \check{\emptyset} \neg y \sim x \neg \text{ ID} \rightarrow \equiv \check{\emptyset}$$

$$\equiv \check{\emptyset} \neg x \sim y \neg \text{ TIF} \rightarrow \equiv \check{\emptyset}$$

$$\equiv x \neg \check{\emptyset} \sim y \neg \text{ ID} \rightarrow \equiv$$

CAUS (x, y)

$$\equiv x \neg \check{\emptyset} \sim y \neg \rightarrow \equiv$$

$$\dot{\equiv} x \neg y \sim \check{\emptyset} \neg \rightarrow \equiv$$

$$x \equiv \dot{\equiv} \neg y \sim \check{\emptyset} \neg \text{ ID} \rightarrow x \equiv$$

$$x \equiv y \neg \dot{\equiv} \sim \check{\emptyset} \neg \text{ ID} \rightarrow x \equiv y$$

$$x \cdot \equiv y \neg \check{\emptyset} \sim \dot{\equiv} \neg \rightarrow x \cdot \equiv y$$

$$x \cdot y \equiv \neg \emptyset \sim \dot{\equiv} \neg$$

$$x \cdot \equiv y \neg \sim \emptyset \neg$$

$$x \equiv y \neg \sim \check{\emptyset} \neg$$

$$x \equiv \neg y \sim \check{\emptyset} \neg \text{ ID}$$

$$\equiv x \neg \cdot \check{\circ} \sim y \cdot \neg \rightarrow \equiv$$

$$(fi)*CAUS(): x \equiv *CAUS(3) \wedge x \equiv y *CAUS(4) \wedge \equiv x *CAUS(10)$$

$$(ncfi)*CAUS(): x \cdot \equiv y *CAUS(5)$$

$$CAUS(y, x)$$

$$\equiv y \neg \cdot \check{\circ} \sim x \cdot \neg \rightarrow \equiv$$

$$\dot{\epsilon} \equiv y \neg \cdot x \sim \check{\circ} \cdot \neg \rightarrow \equiv$$

$$y \equiv \dot{\epsilon} \neg \cdot x \sim \check{\circ} \cdot \neg ID \rightarrow y \equiv$$

$$y \equiv x \neg \cdot \dot{\epsilon} \sim \check{\circ} \cdot \neg ID \rightarrow y \equiv x$$

$$y \cdot \equiv x \neg \cdot \check{\circ} \sim \dot{\epsilon} \cdot \neg \rightarrow y \cdot \equiv x$$

$$y \cdot x \equiv \neg \cdot \emptyset \sim \dot{\epsilon} \cdot \neg$$

$$y \cdot \equiv x \neg \cdot \sim \emptyset \cdot \neg$$

$$y \equiv x \neg \cdot \sim \check{\circ} \cdot \neg$$

$$y \equiv \neg \cdot x \sim \check{\circ} \cdot \neg ID$$

$$\equiv y \neg \cdot \check{\circ} \sim x \cdot \neg \rightarrow \equiv$$

$$(fi)CAUS(): y \equiv *CAUS(3) \wedge y \equiv x *CAUS(4)$$

$$(ncfi)CAUS(): y \cdot \equiv x CAUS(5)$$

The above illustrates mutual causality.

Preliminary Exploration of Inference in Dialectical Logic Symbolism

The inference within the system of dialectical logic symbolism is an extension of the self-returning thinking elaborated in the previous categories. Inferences in dialectical logic

aim to allow thinking to return to itself. The “conclusions” in dialectical reasoning are logical formulas that can return to their “premises.”

The so-called “calculation” in dialectical logic involves employing a self-circulating method of reasoning. Whenever a new “premise” is encountered, we must compute the most efficient “conclusion” that can return to that premise.

As the “calculation” part is still an area under my research, I will only introduce some important “inferences.” Note that I will present only a few of them.

Inference using F (x, y) and AP (a, b) as premises: Type One.

Premise 1 :

$F(x, y)$

$FM(x, y)$

$x \cdot \check{\sim} y \cdot \neg \rightarrow x$

$x \cdot \neg \cdot y \sim \check{\sim}$

$y \cdot \neg \cdot x \sim \emptyset \cdot \neg \text{ ID} \rightarrow y$

$y \cdot \neg \cdot \emptyset \sim \cdot \neg$

$y \cdot \check{\sim} \cdot \neg$

$FM(y, x)$

$y \cdot \check{\sim} x \cdot \neg \rightarrow y$

$$y \cdot \neg \cdot x \sim \check{\circ} \cdot \neg$$

$$x \cdot \neg \cdot y \sim \emptyset \cdot \neg \text{ID} \rightarrow x$$

$$x \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$x \cdot \neg \cdot \check{\circ} \sim \cdot \neg$$

Premise 2 :

AP (a, b)

TIF (a, b)

$$\check{\circ} \cdot \neg \cdot b \sim a \cdot \neg$$

$$\check{\circ} \cdot \neg \cdot a \sim b \cdot \neg$$

MA (b, a)

$$b \cdot \neg \cdot \check{\circ} \sim a \cdot \neg \rightarrow b$$

$$b \cdot \neg \cdot a \sim \check{\circ} \cdot \neg$$

$$b \cdot \neg \cdot a \sim \emptyset \cdot \neg$$

$$b \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$b \cdot \neg \cdot \check{\circ} \sim \cdot \neg$$

MA (a, b)

$$a \cdot \neg \cdot \check{\circ} \sim b \cdot \neg \rightarrow a$$

$$a \cdot \neg \cdot b \sim \check{\circ} \cdot \neg$$

$$a \cdot b \sim \emptyset \cdot \neg$$

$$a \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$a \cdot \check{\emptyset} \sim \cdot \neg$$

Conclusion:

FM (x, a)

$$a \cdot \check{\emptyset} \sim x \cdot \neg \rightarrow a$$

$$a \cdot \neg \cdot x \sim \check{\emptyset} \cdot \neg$$

$$x \cdot \neg \cdot a \sim \emptyset \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$x \cdot \neg \cdot \check{\emptyset} \sim \cdot \neg$$

In this inference, we work with two premises: the first is the category of force, F (x, y), and the second is the category of appearance, AP (a, b). In F (x, y), the free items x and y can be substituted into other categories. In AP (a, b), the items a and b, produced by the transformations within the category of appearance, must be substituted back into their respective last logical formula $b \cdot \neg \cdot \check{\emptyset} \sim \cdot \neg$ and $a \cdot \neg \cdot \check{\emptyset} \sim \cdot \neg$ to maintain the integrity of AP (a, b).

The category of AP (a, b) itself is composed of three categories: TIF (a, b), MA (b, a), and MA (a, b). We now substitute the free item x from FM (x, y) in F (x, y) into the final formula of MA (a, b), $a \cdot \neg \cdot \check{\emptyset} \sim \cdot \neg$, leading to the first conclusion: FM (x, a).

This substitution results in the free item b from $MA(b, a)$ being unable to be substituted into the final formula of $MA(a, b)$, $a \cdot \neg \sim \cdot \neg$, indicating a temporary disintegration of $AP(a, b)$. However, the free item a from $MA(a, b)$ remains unused and thus can be substituted into the final formula of $MA(b, a)$, $b \cdot \neg \sim \cdot \neg$, to reactivate $AP(a, b)$. Simultaneously, the free item x from $FM(x, a)$ can move freely back to become the free item x in $FM(x, y)$. This completes the cycle of inference, which we express simply with the following logical format:

Premises:

$F(x, y)$

$AP(a, b)$

Conclusion:

$FM(x, a)$

Inference using $F(x, y)$ and $AP(a, b)$ as premises: Type Two.

The premises one and two are the same as the first inference mentioned above, so they are not listed here again. However, based on different substitutions, we have the following different conclusions:

Conclusion 1 :

$FM(x, a)$

$$a \cdot \neg \check{\circ} \sim x \cdot \neg \rightarrow a$$

$$a \cdot \neg \cdot x \sim \check{\circ} \cdot \neg$$

$$x \cdot \neg \cdot a \sim \emptyset \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$x \cdot \neg \cdot \check{\circ} \sim \cdot \neg$$

Conclusion 2 :

FM (x, b)

$$b \cdot \neg \cdot \check{\circ} \sim y \cdot \neg \rightarrow b$$

$$b \cdot \neg \cdot y \sim \check{\circ} \cdot \neg$$

$$y \cdot \neg \cdot b \sim \emptyset \cdot \neg \rightarrow y$$

$$y \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$y \cdot \neg \cdot \check{\circ} \sim \cdot \neg$$

Conclusion 3 :

TIF (a, b)

$$\check{\circ} \cdot \neg \cdot b \sim a \cdot \neg$$

$$\checkmark \neg \cdot a \sim b \cdot \neg$$

$$\checkmark$$

$$BC$$

$$\neg \cdot \sim \checkmark \cdot \neg$$

$$\neg \cdot \checkmark \sim \cdot \neg$$

In this inference, we insert the free item x from $FM(x, y)$ in the category $F(x, y)$ into the last formula of $MA(a, b)$ in $AP(a, b)$, represented by $a \neg \cdot \checkmark \sim \cdot \neg$. This substitution results in the first conclusion: $FM(x, a)$. Similarly, the free item y from $FM(y, x)$ is substituted into the last formula of $MA(b, a)$, represented by $b \neg \cdot \checkmark \sim \cdot \neg$, leading to the second conclusion: $FM(y, b)$.

Now that the free items x and y from $FM(x, a)$ and $FM(y, b)$ can freely move back into the free items of $FM(x, y)$ in $F(x, y)$, the category of $F(x, y)$ is restored. However, since there are still two free items a and b from $FM(x, a)$ and $FM(y, b)$ that have not found places for substitution, the reasoning cannot yet cease.

At this point, the second premise, $AP(a, b)$, degenerates into $TIF(a, b)$ because no item is substituted into the final formulas of $MA(a, b)$ and $MA(b, a)$. $TIF(a, b)$ then further degenerates into \checkmark , entering into the BC . As the category of becoming progresses to its

second logical formula $\neg \cdot \check{\sim} \cdot \neg$, we can substitute the two free items a and b from FM (x, a) and FM (y, b) in any order, potentially forming either the first formulas of MA (a, b) or MA (b, a) categories. Thus, TIF (a, b) is restored through transformations stemming from using ID with MA (a, b) or MA (b, a), further restoring the category of AP (a, b). This completes the cycle of inference, which is succinctly expressed with the following logical format:

Premises:

F (x, y)

AP (a, b)

Conclusions:

FM (x, a)

FM (y, b)

BC

Inference using F (x, y) and AP (a, b) as premises: Type Three.

Conclusion 1 :

FM (x, a)

$a \neg \cdot \check{\sim} x \cdot \neg \rightarrow a$

$$a \cdot \neg \cdot x \sim \check{o} \cdot \neg$$

$$x \cdot \neg \cdot a \sim \emptyset \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$x \cdot \neg \cdot \check{o} \sim \cdot \neg$$

Conclusion 2 :

FM (y, b)

$$b \cdot \neg \cdot \check{o} \sim y \cdot \neg \rightarrow b$$

$$b \cdot \neg \cdot y \sim \check{o} \cdot \neg$$

$$y \cdot \neg \cdot b \sim \emptyset \cdot \neg \rightarrow y$$

$$y \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$y \cdot \neg \cdot \check{o} \sim \cdot \neg$$

Conclusion 3:

FM (a, y)

$$y \cdot \neg \cdot \check{o} \sim a \cdot \neg \rightarrow y$$

$$y \cdot \neg \cdot a \sim \check{o} \cdot \neg$$

$$a \cdot \neg \cdot y \sim \emptyset \cdot \neg \rightarrow a$$

$$a \cdot \neg \cdot \emptyset \sim \cdot \neg$$

$$a \cdot \check{\sim} \cdot \neg$$

Conclusion 4 :

FM (b, x)

$$x \cdot \check{\sim} \cdot b \cdot \neg \rightarrow x$$

$$x \cdot \neg \cdot b \cdot \check{\sim} \cdot \neg$$

$$b \cdot \neg \cdot x \cdot \check{\sim} \cdot \neg \rightarrow b$$

$$b \cdot \neg \cdot \check{\sim} \cdot \neg$$

$$b \cdot \neg \cdot \check{\sim} \cdot \neg$$

In this inference, we work with two premises: F (x, y) and AP (a, b). The process to reach conclusions one FM (x, a) and two FM (y, b) is similar to previous inferences. However, the free items x and y from FM (x, a) and FM (y, b) are not immediately moved back into the free items of F (x, y); thus, F (x, y) is temporarily not restored.

The reasoning now involves substituting the free items a and b from FM (x, a) and FM (y, b) oppositely into the last logical formulas of each other, $x \cdot \check{\sim} \cdot \neg$ and $y \cdot \check{\sim} \cdot \neg$. This leads to forming Conclusions Three and Four with the first logical formulas of FM (a, y) and FM (b, x) respectively: $y \cdot \check{\sim} \cdot a \cdot \neg \rightarrow y$ and $x \cdot \check{\sim} \cdot b \cdot \neg \rightarrow x$.

The free items x and y from $FM(a, y)$ and $FM(b, x)$ can now move back into the free items of $F(x, y)$, restoring $F(x, y)$. Additionally, the free items a and b from $FM(a, y)$ and $FM(b, x)$, can be substituted into the last logical formulas of each other, forming either $MA(a, b)$ or $MA(b, a)$, thus restoring $AP(a, b)$. This completes the cycle of inference, which we express with the following logical format:

Premises:

$F(x, y)$

$AP(a, b)$

Conclusions:

$FM(x, a)$

$FM(y, b)$

$FM(a, y)$

$FM(b, x)$

When seeking to expand knowledge, there is a very useful principle, which is:

$TIF(a, b)$

$\ddot{o}\neg\cdot b\sim a\cdot\neg$

$\checkmark \neg \cdot a \sim b \cdot \neg$

\checkmark

TIF (c, d)

$\checkmark \neg \cdot c \sim d \cdot \neg$

$\checkmark \neg \cdot d \sim c \cdot \neg$

This extends indefinitely.

Result

Through the establishment of various categories within the dialectical logic symbol system and the exploration of its potential inferential processes, we now have a tool that can accurately describe the different states of consciousness formed by the “subject/object” relationship. This allows us to overcome the complexity and obscurity for which idealism and continental philosophy have often been criticized.

The operation of this system also enables us to explore fundamental linguistic questions, such as “How is a noun generated?”, which may be related to the category of the thing itself, and “How is an adjective generated?”, connected to the first logical position in the category of determined being within the Doctrine of Being. Finally, we can examine “How is a definitional structure created and transformed?”, a question that pertains to the entire scope of the category of actuality.

Most importantly, this system provides us with a precise tool to incorporate the self into its operations. This has significant philosophical implications, as it suggests that our mind possesses a unique origin and mode of operation that cannot be reduced to mathematics or traditional logic.

Discussion

The creation of this dialectical logic symbol system was primarily aimed at capturing the contradictions, transformations, and truths of the human mind. If we try to consider the 'qualitative/quantitative' methodological divide, this system can be said to be designed to provide a precise logical tool for qualitative methods. In an era where quantitative methods are so meticulously developed, the lack of a logic system that can accommodate and utilize contradictions in an irreducible qualitative manner means that areas such as human emotions, society, ethics, and religion risk being interpreted through a predominantly materialist lens, because we tend to choose the model with the greatest explanatory and predictive power. From a humanistic perspective, this tendency towards reductive interpretation is not conducive to a healthy understanding of human experience.

There are always parts of the mind that are excluded from these quantitative frameworks, and as we become increasingly detached from the principles governing these mental forces, our understanding of them fades. In my experience with the Buddhist meditative system, I have

seen a rigorously structured mental process that is effective in guiding the human mind—one that cannot be quantified or reduced to numerical representation. This system, therefore, seeks to offer a framework capable of reflecting these qualitative aspects of human consciousness and spirituality.

Naturally, I am inclined to discuss how this novel system relates to traditional logic, particularly since I have adopted the negation symbol \neg , which aligns with previous logical frameworks. However, at this stage, I am not yet in a position to fully engage in this discussion, as the system is still in its early stages of development, lacking sufficient application and critique. I hope that as this system is adopted and refined in the future, these connections will become clearer. After all, we are entering a philosophical era, driven by the development of artificial intelligence, where serious reflection on the nature of human consciousness is more pressing than ever.

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