### Simulating (partial) Halt Deciders Defeat the Halting Problem Proofs

A simulating halt decider correctly predicts whether or not its correctly simulated input can possibly reach its own final state and halt. It does this by correctly recognizing several non-halting behavior patterns in a finite number of steps of correct simulation. Inputs that do terminate are simply simulated until they complete.

When a simulating halt decider correctly simulates N steps of its input it derives the exact same N steps that a pure UTM would derive because it is itself a UTM with extra features.

My reviewers cannot show that any of the extra features added to the UTM change the behavior of the simulated input for the first N steps of simulation:

- Watching the behavior doesn't change it.
- Matching non-halting behavior patterns doesn't change it
- Even aborting the simulation after N steps doesn't change the first N steps.

Because of all this we can know that the first N steps of input D simulated by simulating halt decider H are the actual behavior that D presents to H for these same N steps.

**computation that halts...** "the Turing machine will halt whenever it enters a final state" (Linz:1990:234)

When we see (after N steps) that D correctly simulated by H cannot possibly reach its simulated final state in any finite number of steps of correct simulation then we have conclusive proof that D presents non-halting behavior to H.

# **Summary of Linz Halting Problem Proof**

The Linz halting problem proof constructs its counter-example input  $\langle \hat{H} \rangle$  on the basis of prepending and appending states to the original Linz H, (assumed halt decider) thus is named embedded H.

#### Original Linz Turing Machine H

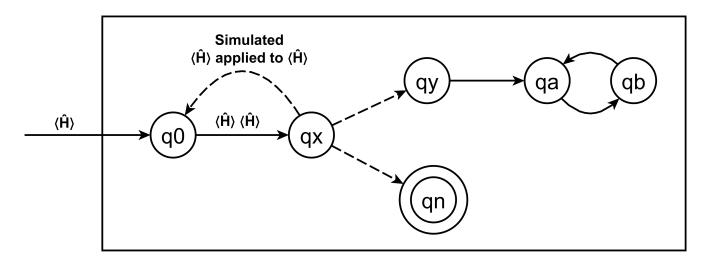
H.q0 ⟨M⟩ w ⊢\* H.qy // Turing Machine description M and finite string w, reject state H.q0 ⟨M⟩ w ⊢\* Hqn // Turing Machine description M and finite string w, accept state

(q0) is prepended to H to copy the  $\langle \hat{H} \rangle$  input of  $\hat{H}$ . The transition from (qa) to (qb) is the conventional infinite loop appended to the (qy) accept state of embedded\_H.  $\vdash$ \* indicates an arbitrary number of moves.

The Linz term "move" means a state transition and its corresponding tape head action {move left, move right, read, write}.

$$\hat{H}$$
.q0 (M)  $\vdash$ \* embedded\_H (M) (M)  $\vdash$ \*  $\hat{H}$ .qy  $\infty$   $\hat{H}$ .q0 (M)  $\vdash$ \* embedded\_H (M) (M)  $\vdash$ \*  $\hat{H}$ .qn

# **Analysis of Linz Halting Problem Proof**



#### When Ĥ is applied to 〈Ĥ〉

 $\hat{H}$ .q0  $\langle \hat{H} \rangle \vdash^* \text{embedded}_H \langle \hat{H} \rangle \langle \hat{H} \rangle \vdash^* \hat{H}$ .qy  $\infty$   $\hat{H}$ .q0  $\langle \hat{H} \rangle \vdash^* \text{embedded}_H \langle \hat{H} \rangle \langle \hat{H} \rangle \vdash^* \hat{H}$ .qn

(q0) The input  $\langle \hat{H} \rangle$  is copied then transitions to (qx)

(qx) embedded\_H is applied to  $\langle \hat{H} \rangle \langle \hat{H} \rangle$  (input and copy)

which simulates  $\langle \hat{H} \rangle$  applied to  $\langle \hat{H} \rangle$  which begins at its own (q0) to repeat the process.

This process continues to repeat until embedded\_H recognizes the repeating pattern and aborts its simulation of  $\langle \hat{H} \rangle \langle \hat{H} \rangle$ . embedded\_H can see the same repeating pattern that we see.

**computation that halts...** "the Turing machine will halt whenever it enters a final state" (Linz:1990:234)

### Every "rebuttal" simply ignores this key fact

 $\langle \hat{H} \rangle \langle \hat{H} \rangle$  correctly simulated by embedded\_H cannot possibly reach its own simulated final state of  $\langle \hat{H}.qn \rangle$  and halt in any finite number of steps of correct simulation.

**Linz, Peter 1990.** An Introduction to Formal Languages and Automata. Lexington/Toronto: D. C. Heath and Company. (317-320)

#### Concrete Rebuttal of Halting Theorem with fully operational software

Simulating halt decider H correctly predicts whether or not D correctly simulated by H can possibly reach its own final state in any finite number of correctly simulated steps. It does this by correctly recognizing several non-halting behavior patterns in a finite number of steps of correct simulation. Inputs that do terminate are simply simulated until they complete.

```
01 int D(int (*x)())
02 {
03    int Halt_Status = H(x, x);
04    if (Halt_Status)
05        HERE: goto HERE;
06    return Halt_Status;
07 }
08
09 void main()
10 {
11    H(D,D);
12 }
```

We form an isomorphism to the Linz Turing Machine analysis in the C programming language. It has the key required element that D attempts to do the opposite of whatever value H returns.

For any program H that might determine whether programs halt, a "pathological" program D, called with some input, can pass its own source and its input to H and then specifically do the opposite of what H predicts D will do. No H can exist that handles this case. <a href="https://en.wikipedia.org/wiki/Halting\_problem">https://en.wikipedia.org/wiki/Halting\_problem</a>

As was proved above in the Linz example the first N steps of input D correctly simulated by simulating halt decider H are the actual behavior that D presents to H for these same N steps.

When H correctly determines (after N simulated steps) that D correctly simulated by H cannot possibly reach its own final state on line 6 and halt then H has conclusive proof that D presents non-halting behavior to H.

# Here is the sequence if H would never abort it simulation main() calls H(D,D) that simulates D(D) at line 11

**keeps repeating:** simulated D(D) calls simulated H(D,D) that simulates D(D) at line 03 ...

This proves that D correctly simulated by H cannot possibly reach its own simulated final state at line 6 and halt in any finite number of steps of correct simulation.

When H aborts its simulation and returns 0 it is only affirming this verified fact.

H(D,D) fully operational in x86utm operating system: <a href="https://github.com/plolcott/x86utm">https://github.com/plolcott/x86utm</a>

Source-code of several different partial halt deciders and their sample input. <a href="https://github.com/plolcott/x86utm/blob/master/Halt7.c">https://github.com/plolcott/x86utm/blob/master/Halt7.c</a>