

Deductively Sound Formal Proofs

Using the sound deductive inference model as the basis of the a notion of a formal system defines [Deductively Sound Formal Proofs]. Within (DSFP) closed Well-formed formula that were undecidable in other formal systems are excluded on the basis that they do not belong to deductively sound inference.

The sound deductive inference model specifies:

[a connected sequence of valid deductions from true premises to a true conclusion]

A wf C is said to be a consequence in S of a set Γ of wfs if and only if there is a sequence B_1, \dots, B_k of wfs such that C is B_k and, for each i , either B_i is an axiom or B_i is in Γ , or B_i is a direct consequence by some rule of inference of some of the preceding wfs in the sequence. Such a sequence is called a proof (or deduction) of C from Γ . The members of Γ are called the hypotheses or premisses of the proof. We use $\Gamma \vdash C$ as an abbreviation for "C is a consequence of Γ "... (Mendelson 2015:28)

Within the sound deductive inference model True(x) is defined as:

- (1) An expression of language having the semantic value of Boolean True.
- (2) The consequence of any formal mathematical proof where every element of its set of premises is True.

Let \mathcal{T} be such a theory. Then the elementary statements which belong to \mathcal{T} we shall call the elementary theorems of \mathcal{T} ; we also say that these elementary statements are true for \mathcal{T} . Thus, given \mathcal{T} , **an elementary theorem is an elementary statement which is true**. A theory is thus a way of picking out from the statements of F a certain subclass of true statements. (Curry 1977:45)

To define formal systems having true premises we adopt the (Curry 1977:45) convention stipulating that axioms (Curry elementary theorems) are true. This stipulation derives a corresponding pair of True(x) and False(x) predicates:

- (1) True(x) := ($\vdash x$) " := " is defined below **
- (2) False(x) := ($\vdash \neg x$)

With True and False formalized we specify a semantic criterion of Well-formedness:

- (3) Deductively_Sound_Consequent(x) := (True(x) \vee False(x))

Sound deduction excludes conclusions/consequences not based on sound deduction, as deductively unsound conclusions. Thus a semantic criteria of well-formedness is provided in addition to the conventional syntactic notion of WFF.

Mendelson, Elliott 2015. Introduction to Mathematical Logic (sixth edition). Boca Raton: Taylor & Francis Group LLC.

Curry, Harkell B. 1977. Foundations of Mathematical Logic. New York: Dover Publications, Inc.

** LHS := RHS (The LHS is defined as an alias for the RHS)

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