

Halting problem proofs refuted on the basis of software engineering

This is an explanation of a key new insight into the halting problem provided in the language of software engineering. Technical computer science terms are explained using software engineering terms. No knowledge of the halting problem is required.

It is based on fully operational software executed in the x86utm operating system. The x86utm operating system (based on an excellent open source x86 emulator) was created to study the details of the halting problem proof counter-examples at the much higher level of abstraction of C/x86.

```
typedef void (*ptr)();
int H(ptr p, ptr i); // simulating halt decider

void P(ptr x)
{
    int Halt_Status = H(x, x);
    if (Halt_Status)
        HERE: goto HERE;
    return;
}

int main()
{
    output("Input_Halts = ", H(P, P));
}
```

When simulating halt decider H(P,P) simulates its input we can see that:

- (1) Function H() is called from P().
- (2) With the same arguments to H().
- (3) With no instructions in P preceding its invocation of H(P,P).

The above shows that the simulated P cannot possibly terminate normally. Because H can see the same (1)(2)(3) that we see H aborts its simulation of P and rejects P as non-halting.

In computability theory, the halting problem is the problem of determining, from a description of an arbitrary computer program and an input, whether the program will finish running, or continue to run forever. Alan Turing proved in 1936 that a general algorithm to solve the halting problem for all possible program-input pairs cannot exist.

For any program H that might determine if programs halt, a "pathological" program P, called with some input, can pass its own source and its input to H and then specifically do the opposite of what H predicts P will do. **No H can exist that handles this case.**
https://en.wikipedia.org/wiki/Halting_problem

H and P implement the exact pathological relationship to each other as described above. Because H(P,P) does handle this case the above halting problem undecidable input template has been refuted.

When this halt deciding principle understood to be correct:

A halt decider must compute the mapping from its inputs to an accept or reject state on the basis of the actual behavior that is actually specified by these inputs.

Then (by logical necessity) this implements that principle:

Every simulating halt decider that correctly simulates its input until it correctly predicts that this simulated input would never terminate normally, correctly rejects this input as non-halting.

H is a [Pure function](#) thus implements a [Computable function](#) Thus H is Turing computable.

Example 01: H0 correctly determines that Infinite_Loop() never halts

```
void Infinite_Loop()
{
    HERE: goto HERE;
}

int main()
{
    Output("Input_Halts = ", H0((u32)Infinite_Loop));
}
```

```
_Infinite_Loop()
[00001102](01) 55      push ebp
[00001103](02) 8bec     mov ebp,esp
[00001105](02) ebfe     jmp 00001105
[00001107](01) 5d      pop ebp
[00001108](01) c3      ret
Size in bytes:(0007) [00001108]
```

```
_main()
[00001192](01) 55      push ebp
[00001193](02) 8bec     mov ebp,esp
[00001195](05) 6802110000 push 00001102
[0000119a](05) e8d3fbffff call 00000d72
[0000119f](03) 83c404   add esp,+04
[000011a2](01) 50      push eax
[000011a3](05) 68a3040000 push 000004a3
[000011a8](05) e845f3ffff call 000004f2
[000011ad](03) 83c408   add esp,+08
[000011b0](02) 33c0     xor eax,eax
[000011b2](01) 5d      pop ebp
[000011b3](01) c3      ret
Size in bytes:(0034) [000011b3]
```

machine address	stack address	stack data	machine code	assembly language
[00001192]	[00101ef8]	[00000000]	55	push ebp
[00001193]	[00101ef8]	[00000000]	8bec	mov ebp,esp
[00001195]	[00101ef4]	[00001102]	6802110000	push 00001102
[0000119a]	[00101ef0]	[0000119f]	e8d3fbffff	call 00000d72

```
H0: Begin Simulation      Execution Trace Stored at:211fac
[00001102][00211f9c][00211fa0] 55      push ebp
[00001103][00211f9c][00211fa0] 8bec     mov ebp,esp
[00001105][00211f9c][00211fa0] ebfe     jmp 00001105
[00001105][00211f9c][00211fa0] ebfe     jmp 00001105
H0: Infinite Loop Detected Simulation Stopped
```

```
if (current->Simplified_Opcode == JMP) // JMP
    if (current->Decode_Target <= current->Address) // upward
        if (traced->Address == current->Decode_Target) // to this address
            if (Conditional_Branch_Count == 0) // no escape
                return 1;
```

```
[0000119f][00101ef8][00000000] 83c404   add esp,+04
[000011a2][00101ef4][00000000] 50      push eax
[000011a3][00101ef0][000004a3] 68a3040000 push 000004a3
[000011a8][00101ef0][000004a3] e845f3ffff call 000004f2
Input_Halts = 0
[000011ad][00101ef8][00000000] 83c408   add esp,+08
[000011b0][00101ef8][00000000] 33c0     xor eax,eax
[000011b2][00101efc][00100000] 5d      pop ebp
[000011b3][00101f00][00000004] c3      ret
Number of Instructions Executed(554) == 8 Pages
```

Example 02: H correctly determines that Infinite_Recursion() never halts

```
void Infinite_Recursion(int N)
{
    Infinite_Recursion(N);
}

int main()
{
    output("Input_Halts = ", H((u32)Infinite_Recursion, 0x777));
}
```

```
_Infinite_Recursion()
[000010f2](01) 55      push ebp
[000010f3](02) 8bec     mov ebp,esp
[000010f5](03) 8b4508   mov eax,[ebp+08]
[000010f8](01) 50      push eax
[000010f9](05) e8f4ffff call 000010f2
[000010fe](03) 83c404   add esp,+04
[00001101](01) 5d      pop ebp
[00001102](01) c3      ret
Size in bytes:(0017) [00001102]
```

```
_main()
[000011b2](01) 55      push ebp
[000011b3](02) 8bec     mov ebp,esp
[000011b5](05) 6877070000 push 00000777
[000011ba](05) 68f2100000 push 000010f2
[000011bf](05) e8aefdffff call 00000f72
[000011c4](03) 83c408   add esp,+08
[000011c7](01) 50      push eax
[000011c8](05) 68a3040000 push 000004a3
[000011cd](05) e820f3ffff call 000004f2
[000011d2](03) 83c408   add esp,+08
[000011d5](02) 33c0     xor eax,eax
[000011d7](01) 5d      pop ebp
[000011d8](01) c3      ret
Size in bytes:(0039) [000011d8]
```

machine address	stack address	stack data	machine code	assembly language
[000011b2]	[00101f39]	[00000000]	55	push ebp
[000011b3]	[00101f39]	[00000000]	8bec	mov ebp,esp
[000011b5]	[00101f35]	[00000777]	6877070000	push 00000777
[000011ba]	[00101f31]	[000010f2]	68f2100000	push 000010f2
[000011bf]	[00101f2d]	[000011c4]	e8aefdffff	call 00000f72

```
H: Begin Simulation Execution Trace Stored at:111fe5
[000010f2][00111fd1][00111fd5] 55      push ebp
[000010f3][00111fd1][00111fd5] 8bec     mov ebp,esp
[000010f5][00111fd1][00111fd5] 8b4508   mov eax,[ebp+08]
[000010f8][00111fd1][00000777] 50      push eax // push 0x777
[000010f9][00111fd1][000010fe] e8f4ffff call 000010f2 // call Infinite_Recursion
[000010f2][00111fd1][00111fd1] 55      push ebp
[000010f3][00111fd1][00111fd1] 8bec     mov ebp,esp
[000010f5][00111fd1][00111fd1] 8b4508   mov eax,[ebp+08]
[000010f8][00111fd1][00000777] 50      push eax // push 0x777
[000010f9][00111fd1][000010fe] e8f4ffff call 000010f2 // call Infinite_Recursion
H: Infinite Recursion Detected Simulation Stopped
```

```
if (current->Simplified_Opcode == CALL)
    if (current->Simplified_Opcode == traced->Simplified_Opcode) // CALL
        if (current->Address == traced->Address) // from same address
            if (current->Decode_Target == traced->Decode_Target) // to Same Function
                if (Conditional_Branch_Count == 0) // no escape
                    return 2;
```

```
[000011c4][00101f39][00000000] 83c408   add esp,+08
[000011c7][00101f35][00000000] 50      push eax
[000011c8][00101f31][000004a3] 68a3040000 push 000004a3
[000011cd][00101f31][000004a3] e820f3ffff call 000004f2
Input_Halts = 0
[000011d2][00101f39][00000000] 83c408   add esp,+08
[000011d5][00101f39][00000000] 33c0     xor eax,eax
[000011d7][00101f3d][00000018] 5d      pop ebp
[000011d8][00101f41][00000000] c3      ret
Number of Instructions Executed(1118) == 17 Pages
```

Example 03: H(P,P) correctly determines that its input never halts

```
void P(ptr x)
{
    int Halt_Status = H(x, x);
    if (Halt_Status)
        HERE: goto HERE;
    return;
}

int main()
{
    output("Input_Halts = ", H(P, P));
}
```

From a purely software engineering perspective (anchored in the semantics of the x86 language) it is proven that H(P,P) correctly predicts that its correct and complete x86 emulation of its input would never reach the "ret" instruction (final state) of this input. **Copyright 2022 PL Olcott**

```
_P()
[000013c6] (01) 55      push ebp          // Save Base Pointer register onto the stack
[000013c7] (02) 8bec     mov ebp,esp      // Load Base Pointer with Stack Pointer
[000013c9] (01) 51      push ecx         // Save the value of ecx on the stack
[000013ca] (03) 8b4508   mov eax,[ebp+08] // Load eax with argument to P
[000013cd] (01) 50      push eax         // push 2nd argument to H onto the stack
[000013ce] (03) 8b4d08   mov ecx,[ebp+08] // Load ecx with with argument to P
[000013d1] (01) 51      push ecx         // push 1st argument to H onto the stack
[000013d2] (05) e2ffdf    call 00001106    // push return address on the stack; call simulated H
[000013d7] (03) 83c408   add esp,+08     // remove call arguments from stack
[000013da] (03) 8945fc   mov [ebp-04],eax // load Halt_Status with return value from H
[000013dd] (04) 837dfc00  cmp dword [ebp-04],+00 // compare Halt_Status to 0
[000013e1] (02) 7402     jz 000013e5     // if Halt_Status == 0 goto 000013e5
[000013e3] (02) ebfe     jmp 000013e3    // goto 13e3
[000013e5] (02) 8be5     mov esp,ebp     // Load Stack Pointer with Base Pointer
[000013e7] (01) 5d      pop ebp         // Restore Base Pointer value from stack
[000013e8] (01) c3      ret             // return to caller
Size in bytes:(0035) [000013e8]
```

```
_main()
[000013f6] (01) 55      push ebp          // Save Base Pointer register onto the stack
[000013f7] (02) 8bec     mov ebp,esp      // Load Base Pointer with Stack Pointer
[000013f9] (05) 68c6130000  push 000013c6   // Push P (2nd argument to H) onto the stack
[000013fe] (05) 68c6130000  push 000013c6   // Push P (1nd argument to H) onto the stack
[00001403] (05) e8fefcffff    call 00001106    // push return address onto the stack and call executed H
[00001408] (03) 83c408   add esp,+08     // remove call arguments from stack frame
[0000140b] (01) 50      push eax         // Push return value from H onto the stack
[0000140c] (05) 6837050000  push 00000537   // Push address of "Input_Halts = " onto the stack
[00001411] (05) e870f1ffff    call 00000586    // call Output with its pushed arguments.
[00001416] (03) 83c408   add esp,+08     // remove call arguments from stack frame
[00001419] (02) 33c0     xor eax,eax     // set eax to 0
[0000141b] (01) 5d      pop ebp         // Restore Base Pointer register from stack
[0000141c] (01) c3      ret             // return to 0 operating system
Size in bytes:(0039) [0000141c]
```

machine address	stack address	stack data	machine code	assembly language
[000013f6]	[0010235f]	[00000000]	55	push ebp // Save Base Pointer register onto the stack
[000013f7]	[0010235f]	[00000000]	8bec	mov ebp,esp // Load Base Pointer with Stack Pointer
[000013f9]	[0010235b]	[000013c6]	68c6130000	push 000013c6 // Push P (2nd argument to H) onto the stack
[000013fe]	[00102357]	[000013c6]	68c6130000	push 000013c6 // Push P (1nd argument to H) onto the stack
[00001403]	[00102353]	[00001408]	e8fefcffff	call 00001106 // push return address; call executed H

H: Begin Simulation Execution Trace Stored at:11240b

```
Address_of_H:1106
[000013c6] [001123f7] [001123fb] 55      push ebp          // Save Base Pointer register onto the stack
[000013c7] [001123f7] [001123fb] 8bec     mov ebp,esp      // Load Base Pointer with Stack Pointer
[000013c9] [001123f3] [001023c7] 51      push ecx         // Save the value of ecx on the stack
[000013ca] [001123f3] [001023c7] 8b4508   mov eax,[ebp+08] // Load eax with argument to P
[000013cd] [001123ef] [000013c6] 50      push eax         // push 2nd argument to H onto the stack
[000013ce] [001123ef] [000013c6] 8b4d08   mov ecx,[ebp+08] // Load ecx with with argument to P
[000013d1] [001123eb] [000013c6] 51      push ecx         // push 1st argument to H onto the stack
[000013d2] [001123e7] [000013d7] e2ffdf    call 00001106    // push return address; call simulated H
H: Infinitely Recursive Simulation Detected Simulation Stopped
```

```
[00001408] [0010235f] [00000000] 83c408   add esp,+08     // remove call arguments from stack frame
[0000140b] [0010235b] [00000000] 50      push eax       // Push return value from H onto the stack
[0000140c] [00102357] [00000537] 6837050000  push 00000537 // Push address of "Input_Halts = " onto stack
[00001411] [00102357] [00000537] e870f1ffff    call 00000586 // call Output with its pushed arguments
Input_Halts = 0
[00001416] [0010235f] [00000000] 83c408   add esp,+08     // remove call arguments from stack frame
[00001419] [0010235f] [00000000] 33c0     xor eax,eax     // set eax to 0
[0000141b] [00102363] [00000018] 5d      pop ebp         // Restore Base Pointer register from stack
[0000141c] [00102367] [00000000] c3      ret             // return to 0 operating system
Number of Instructions Executed(987) == 15 Pages
```

Example 04: An impossible program: Strachey(1965)

The Computer Journal, Volume 7, Issue 4, January 1965, Page 313,

<https://doi.org/10.1093/comjnl/7.4.313>

```
typedef void (*ptr)();
// rec routine P
// $L :if T[P] go to L
// Return $
void Strachey_P()
{
  L: if (T(Strachey_P)) goto L;
  return;
}

int main()
{
  output("Input_Halts = ", T(Strachey_P));
}
```

```
_Strachey_P()
[000012a6] (01) 55      push ebp
[000012a7] (02) 8bec     mov ebp,esp
[000012a9] (05) 68a6120000 push 000012a6
[000012ae] (05) e833fcffff call 00000ee6
[000012b3] (03) 83c404   add esp,+04
[000012b6] (02) 85c0     test eax,eax
[000012b8] (02) 7402     jz 000012bc
[000012ba] (02) ebed     jmp 000012a9
[000012bc] (01) 5d      pop ebp
[000012bd] (01) c3      ret
Size in bytes:(0024) [000012bd]
```

```
_main()
[00001346] (01) 55      push ebp
[00001347] (02) 8bec     mov ebp,esp
[00001349] (05) 68a6120000 push 000012a6
[0000134e] (05) e893fbffff call 00000ee6
[00001353] (03) 83c404   add esp,+04
[00001356] (01) 50      push eax
[00001357] (05) 6817050000 push 00000517
[0000135c] (05) e805f2ffff call 00000566
[00001361] (03) 83c408   add esp,+08
[00001364] (02) 33c0     xor eax,eax
[00001366] (01) 5d      pop ebp
[00001367] (01) c3      ret
Size in bytes:(0034) [00001367]
```

machine address	stack address	stack data	machine code	assembly language
[00001346]	[0010221b]	[00000000]	55	push ebp
[00001347]	[0010221b]	[00000000]	8bec	mov ebp,esp
[00001349]	[00102217]	[000012a6]	68a6120000	push 000012a6
[0000134e]	[00102213]	[00001353]	e893fbffff	call 00000ee6

T: Begin Simulation Execution Trace Stored at:1122c7
Address_of_T:ee6

```
[000012a6] [001122b7] [001122bb] 55      push ebp
[000012a7] [001122b7] [001122bb] 8bec     mov ebp,esp
[000012a9] [001122b3] [000012a6] 68a6120000 push 000012a6
[000012ae] [001122af] [000012b3] e833fcffff call 00000ee6
```

T: Infinitely Recursive Simulation Detected Simulation Stopped

T knows its own machine address and on this basis it can easily examine its stored execution_trace of Strachey_P (see above) to determine:

- Strachey_P is calling T with the same arguments that T was called with.
- No instructions in Strachey_P could possibly escape this otherwise infinitely recursive emulation.
- T aborts its emulation of Strachey_P before its call to T is emulated.

```
[00001353] [0010221b] [00000000] 83c404   add esp,+04
[00001356] [00102217] [00000000] 50      push eax
[00001357] [00102213] [00000517] 6817050000 push 00000517
[0000135c] [00102213] [00000517] e805f2ffff call 00000566
```

Input_Halts = 0

```
[00001361] [0010221b] [00000000] 83c408   add esp,+08
[00001364] [0010221b] [00000000] 33c0     xor eax,eax
[00001366] [0010221f] [00000018] 5d      pop ebp
[00001367] [00102223] [00000000] c3      ret
```

Number of Instructions Executed(538) == 8 Pages

Appendix (Simulating halt decider applied to Peter Linz proof)

The following is the same idea as shown above this time it is applied to the Peter Linz Halting Problem proof. It can only be understood within the context of this proof.

A simulating halt decider (SHD) computes the mapping from its inputs to its own final states on the basis of the behavior of its correctly simulated input.

All of the conventional halting problem counter-example inputs are simply rejected by a simulating halt decider as non-halting because they fail to meet the Linz definition of halting:

computation that halts ... the Turing machine will halt whenever it enters a final state.
(Linz:1990:234)

USENET comp.theory: On 4/11/2022 3:19 PM, Malcolm McLean wrote:

- > PO's idea is to have a simulator with an infinite cycle detector.
- > You would achieve this by modifying a UTM, so describing it as
- > a "modified UTM", or "acts like a UTM until it detects an infinite
- > cycle", is reasonable. And such a machine is a fairly powerful
- > halt decider. Even if the infinite cycle detector isn't very
- > sophisticated, it will still catch a large subset of non-halting
- > machines.

The following simplifies the syntax for the definition of the Linz Turing machine \hat{H} . There is no need for the infinite loop after $H.qy$ because it is never reached. The halting criteria has been adapted so that it applies to a simulating halt decider (SHD).

$\hat{H}.q_0 \langle \hat{H} \rangle \vdash^* H \langle \hat{H} \rangle \langle \hat{H} \rangle \vdash^* \hat{H}.qy$

If the correctly simulated input $\langle \hat{H} \rangle \langle \hat{H} \rangle$ to H would reach its own final state of $\langle \hat{H}.qy \rangle$ or $\langle \hat{H}.qn \rangle$.

$\hat{H}.q_0 \langle \hat{H} \rangle \vdash^* H \langle \hat{H} \rangle \langle \hat{H} \rangle \vdash^* \hat{H}.qn$

If the correctly simulated input $\langle \hat{H} \rangle \langle \hat{H} \rangle$ to H would never reach its own final state of $\langle \hat{H}.qy \rangle$ or $\langle \hat{H}.qn \rangle$.

When \hat{H} is applied to $\langle \hat{H} \rangle$ // subscripts indicate unique finite strings

\hat{H} copies its input $\langle \hat{H}_0 \rangle$ to $\langle \hat{H}_1 \rangle$ then H simulates $\langle \hat{H}_0 \rangle \langle \hat{H}_1 \rangle$

Then these steps would keep repeating: (unless their simulation is aborted)

\hat{H}_0 copies its input $\langle \hat{H}_1 \rangle$ to $\langle \hat{H}_2 \rangle$ then H_0 simulates $\langle \hat{H}_1 \rangle \langle \hat{H}_2 \rangle$

\hat{H}_1 copies its input $\langle \hat{H}_2 \rangle$ to $\langle \hat{H}_3 \rangle$ then H_1 simulates $\langle \hat{H}_2 \rangle \langle \hat{H}_3 \rangle$

\hat{H}_2 copies its input $\langle \hat{H}_3 \rangle$ to $\langle \hat{H}_4 \rangle$ then H_2 simulates $\langle \hat{H}_3 \rangle \langle \hat{H}_4 \rangle$...

Since we can see that the simulated input: $\langle \hat{H}_0 \rangle$ to H would never reach its own final state of $\langle \hat{H}_0.qy \rangle$ or $\langle \hat{H}_0.qn \rangle$ we know that it is non-halting.

Linz, Peter 1990. An Introduction to Formal Languages and Automata. Lexington/Toronto: D. C. Heath and Company. (317-320) **this paper copyright 2022 by PL Olcott**