

# Halting problem proofs refuted on the basis of software engineering

This is an explanation of a key new insight into the halting problem provided in the language of software engineering. Technical computer science terms are explained using software engineering terms. No knowledge of the halting problem is required.

It is based on fully operational software executed in the x86utm operating system. The x86utm operating system (based on an excellent open source x86 emulator) was created to study the details of the halting problem proof counter-examples at the much higher level of abstraction of C/x86.

```
typedef void (*ptr)();
int H(ptr p, ptr i); // simulating halt decider

void P(ptr x)
{
    int Halt_Status = H(x, x);
    if (Halt_Status)
        HERE: goto HERE;
    return;
}

int main()
{
    output("Input_Halts = ", H(P, P));
}
```

**When simulating halt decider H(P,P) simulates its input it can see that:**

- (1) Function H() is called from P().
- (2) With the same arguments to H().
- (3) With no instructions in P preceding its invocation of H(P,P) that could escape repeated simulations.

The above shows that the simulated P cannot possibly (reaches its "return" instruction and) terminate normally. H(P,P) simulates its input then P calls H(P,P) to simulate itself again. When H sees that this otherwise infinitely nested simulation would never end it aborts its simulation of P and rejects P as non-halting.

In computability theory, the halting problem is the problem of determining, from a description of an arbitrary computer program and an input, whether the program will finish running, or continue to run forever. Alan Turing proved in 1936 that a general algorithm to solve the halting problem for all possible program-input pairs cannot exist.

For any program H that might determine if programs halt, a "pathological" program P, called with some input, can pass its own source and its input to H and then specifically do the opposite of what H predicts P will do. **No H can exist that handles this case.**  
[https://en.wikipedia.org/wiki/Halting\\_problem](https://en.wikipedia.org/wiki/Halting_problem)

H and P implement the exact pathological relationship to each other as described above. Because H(P,P) does handle this case the above halting problem undecidable input template has been refuted.

**When this halt deciding principle understood to be correct:**

A halt decider must compute the mapping from its inputs to an accept or reject state on the basis of the actual behavior that is actually specified by these inputs.

Within the common knowledge that the correct simulation of a program (or TM description) accurately measures the actual behavior of this program:

**Then (by logical necessity) this correctly implements the halting deciding principle:**

Every simulating halt decider that correctly simulates its input until it correctly predicts that this simulated input would never terminate normally, correctly rejects this input as non-halting.

H is a [Pure function](#) thus implements a [Computable function](#) Thus H is Turing computable.

A halt decider must compute the mapping from its inputs to an accept or reject state on the basis of the actual behavior that is actually specified by these inputs.

It is common knowledge that a correct simulation of a program is a correct measure of the behavior of this program. The concept of a Universal Turing Machine (UTM) is invalidated unless it is accepted that the correct simulation of a machine description is computationally equivalent to the underlying computation.

**Example 05** proves that both the simulation of the input to H(P,P) and the direct execution of P(P) are correct: The execution trace of behavior of the correctly simulated input to H(P,P) and the execution trace of behavior of the directly executed P(P) exactly matches the line-by-line x86 source code of P. Because these behaviors diverge this proves that the direct execution of P(P) is not the behavior that H(P,P) must report on.

### Example 01: H0 correctly determines that Infinite\_Loop() never halts

```
void Infinite_Loop()
{
    HERE: goto HERE;
}

int main()
{
    Output("Input_Halts = ", H0((u32)Infinite_Loop));
}
```

```
_Infinite_Loop()
[00001102](01) 55          push ebp
[00001103](02) 8bec         mov ebp,esp
[00001105](02) ebfe         jmp 00001105
[00001107](01) 5d          pop ebp
[00001108](01) c3          ret
Size in bytes:(0007) [00001108]
```

```
_main()
[00001192](01) 55          push ebp
[00001193](02) 8bec         mov ebp,esp
[00001195](05) 6802110000 push 00001102
[0000119a](05) e8d3fbffff call 00000d72
[0000119f](03) 83c404      add esp,+04
[000011a2](01) 50          push eax
[000011a3](05) 68a3040000 push 000004a3
[000011a8](05) e845f3ffff call 000004f2
[000011ad](03) 83c408      add esp,+08
[000011b0](02) 33c0        xor eax,eax
[000011b2](01) 5d          pop ebp
[000011b3](01) c3          ret
Size in bytes:(0034) [000011b3]
```

machine address	stack address	stack data	machine code	assembly language
[00001192]	[00101ef8]	[00000000]	55	push ebp
[00001193]	[00101ef8]	[00000000]	8bec	mov ebp,esp
[00001195]	[00101ef4]	[00001102]	6802110000	push 00001102
[0000119a]	[00101ef0]	[0000119f]	e8d3fbffff	call 00000d72

```
H0: Begin Simulation Execution Trace Stored at:211fac
[00001102][00211f9c][00211fa0] 55          push ebp
[00001103][00211f9c][00211fa0] 8bec         mov ebp,esp
[00001105][00211f9c][00211fa0] ebfe         jmp 00001105
[00001105][00211f9c][00211fa0] ebfe         jmp 00001105
H0: Infinite Loop Detected Simulation Stopped
```

```
if (current->Simplified_Opcode == JMP) // JMP
    if (current->Decode_Target <= current->Address) // upward
        if (traced->Address == current->Decode_Target) // to this address
            if (Conditional_Branch_Count == 0) // no escape
                return 1;
```

```
[0000119f][00101ef8][00000000] 83c404      add esp,+04
[000011a2][00101ef4][00000000] 50          push eax
[000011a3][00101ef0][000004a3] 68a3040000 push 000004a3
[000011a8][00101ef0][000004a3] e845f3ffff call 000004f2
Input_Halts = 0
[000011ad][00101ef8][00000000] 83c408      add esp,+08
[000011b0][00101ef8][00000000] 33c0        xor eax,eax
[000011b2][00101efc][00100000] 5d          pop ebp
[000011b3][00101f00][00000004] c3          ret
Number of Instructions Executed(554) == 8 Pages
```

## Example 02: H correctly determines that Infinite\_Recursion() never halts

```
void Infinite_Recursion(int N)
{
    Infinite_Recursion(N);
}

int main()
{
    output("Input_Halts = ", H((u32)Infinite_Recursion, 0x777));
}
```

```
_Infinite_Recursion()
[000010f2](01) 55      push ebp
[000010f3](02) 8bec     mov ebp,esp
[000010f5](03) 8b4508   mov eax,[ebp+08]
[000010f8](01) 50      push eax
[000010f9](05) e8f4ffff call 000010f2
[000010fe](03) 83c404   add esp,+04
[00001101](01) 5d      pop ebp
[00001102](01) c3      ret
Size in bytes:(0017) [00001102]
```

```
_main()
[000011b2](01) 55      push ebp
[000011b3](02) 8bec     mov ebp,esp
[000011b5](05) 6877070000 push 00000777
[000011ba](05) 68f2100000 push 000010f2
[000011bf](05) e8aefdffff call 00000f72
[000011c4](03) 83c408   add esp,+08
[000011c7](01) 50      push eax
[000011c8](05) 68a3040000 push 000004a3
[000011cd](05) e820f3ffff call 000004f2
[000011d2](03) 83c408   add esp,+08
[000011d5](02) 33c0     xor eax,eax
[000011d7](01) 5d      pop ebp
[000011d8](01) c3      ret
Size in bytes:(0039) [000011d8]
```

machine address	stack address	stack data	machine code	assembly language
[000011b2]	[00101f39]	[00000000]	55	push ebp
[000011b3]	[00101f39]	[00000000]	8bec	mov ebp,esp
[000011b5]	[00101f35]	[00000777]	6877070000	push 00000777
[000011ba]	[00101f31]	[000010f2]	68f2100000	push 000010f2
[000011bf]	[00101f2d]	[000011c4]	e8aefdffff	call 00000f72

```
H: Begin Simulation Execution Trace Stored at:111fe5
[000010f2][00111fd1][00111fd5] 55      push ebp
[000010f3][00111fd1][00111fd5] 8bec     mov ebp,esp
[000010f5][00111fd1][00111fd5] 8b4508   mov eax,[ebp+08]
[000010f8][00111fd1][00000777] 50      push eax // push 0x777
[000010f9][00111fd1][000010fe] e8f4ffff call 000010f2 // call Infinite_Recursion
[000010f2][00111fd1][00111fd1] 55      push ebp
[000010f3][00111fd1][00111fd1] 8bec     mov ebp,esp
[000010f5][00111fd1][00111fd1] 8b4508   mov eax,[ebp+08]
[000010f8][00111fd1][00000777] 50      push eax // push 0x777
[000010f9][00111fd1][000010fe] e8f4ffff call 000010f2 // call Infinite_Recursion
H: Infinite Recursion Detected Simulation Stopped
```

```
if (current->Simplified_Opcode == CALL)
    if (current->Simplified_Opcode == traced->Simplified_Opcode) // CALL
        if (current->Address == traced->Address) // from same address
            if (current->Decode_Target == traced->Decode_Target) // to Same Function
                if (Conditional_Branch_Count == 0) // no escape
                    return 2;
```

[000011c4]	[00101f39]	[00000000]	83c408	add esp,+08
[000011c7]	[00101f35]	[00000000]	50	push eax
[000011c8]	[00101f31]	[000004a3]	68a3040000	push 000004a3
[000011cd]	[00101f31]	[000004a3]	e820f3ffff	call 000004f2

```
Input_Halts = 0
[000011d2][00101f39][00000000] 83c408   add esp,+08
[000011d5][00101f39][00000000] 33c0     xor eax,eax
[000011d7][00101f3d][00000018] 5d      pop ebp
[000011d8][00101f41][00000000] c3      ret
Number of Instructions Executed(1118) == 17 Pages
```

### Example 03: H(P,P) correctly determines that its input never halts

```
void P(ptr x)
{
    int Halt_Status = H(x, x);
    if (Halt_Status)
        HERE: goto HERE;
    return;
}

int main()
{
    output("Input_Halts = ", H(P, P));
}
```

From a purely software engineering perspective (anchored in the semantics of the x86 language) it is proven that H(P,P) correctly predicts that its correct and complete x86 emulation of its input would never reach the "ret" instruction (final state) of this input. **Copyright 2022 PL Olcott**

```
_P()
[000013c6] (01) 55          push ebp          // Save Base Pointer register onto the stack
[000013c7] (02) 8bec        mov ebp,esp      // Load Base Pointer with Stack Pointer
[000013c9] (01) 51          push ecx         // Save the value of ecx on the stack
[000013ca] (03) 8b4508      mov eax,[ebp+08] // Load eax with argument to P
[000013cd] (01) 50          push eax         // push 2nd argument to H onto the stack
[000013ce] (03) 8b4d08      mov ecx,[ebp+08] // Load ecx with with argument to P
[000013d1] (01) 51          push ecx         // push 1st argument to H onto the stack
[000013d2] (05) e82ffdffff   call 00001106    // push return address on the stack; call simulated H
[000013d7] (03) 83c408      add esp,+08     // remove call arguments from stack
[000013da] (03) 8945fc      mov [ebp-04],eax // load Halt_Status with return value from H
[000013dd] (04) 837dfc00    cmp dword [ebp-04],+00 // compare Halt_Status to 0
[000013e1] (02) 7402       jz 000013e5     // if Halt_Status == 0 goto 000013e5
[000013e3] (02) ebfe       jmp 000013e3    // goto 13e3
[000013e5] (02) 8be5       mov esp,ebp     // Load Stack Pointer with Base Pointer
[000013e7] (01) 5d          pop ebp         // Restore Base Pointer value from stack
[000013e8] (01) c3          ret             // return to caller
Size in bytes:(0035) [000013e8]
```

```
_main()
[000013f6] (01) 55          push ebp          // Save Base Pointer register onto the stack
[000013f7] (02) 8bec        mov ebp,esp      // Load Base Pointer with Stack Pointer
[000013f9] (05) 68c6130000  push 000013c6    // Push P (2nd argument to H) onto the stack
[000013fe] (05) 68c6130000  push 000013c6    // Push P (1nd argument to H) onto the stack
[00001403] (05) e8fefcffff   call 00001106    // push return address onto the stack and call executed H
[00001408] (03) 83c408      add esp,+08     // remove call arguments from stack frame
[0000140b] (01) 50          push eax         // Push return value from H onto the stack
[0000140c] (05) 6837050000  push 00000537    // Push address of "Input_Halts = " onto the stack
[00001411] (05) e870f1ffff   call 00000586    // call Output with its pushed arguments.
[00001416] (03) 83c408      add esp,+08     // remove call arguments from stack frame
[00001419] (02) 33c0       xor eax,eax     // set eax to 0
[0000141b] (01) 5d          pop ebp         // Restore Base Pointer register from stack
[0000141c] (01) c3          ret             // return to 0 operating system
Size in bytes:(0039) [0000141c]
```

machine address	stack address	stack data	machine code	assembly language
[000013f6]	[0010235f]	[00000000]	55	push ebp
[000013f7]	[0010235f]	[00000000]	8bec	mov ebp,esp
[000013f9]	[0010235b]	[000013c6]	68c6130000	push 000013c6 // Push P (2nd argument to H) onto the stack
[000013fe]	[00102357]	[000013c6]	68c6130000	push 000013c6 // Push P (1nd argument to H) onto the stack
[00001403]	[00102353]	[00001408]	e8fefcffff	call 00001106 // push return address; call executed H

H: Begin Simulation Execution Trace Stored at:11240b

```
Address_of_H:1106
[000013c6] [001123f7] [001123fb] 55          push ebp
[000013c7] [001123f7] [001123fb] 8bec        mov ebp,esp
[000013c9] [001123f3] [001023c7] 51          push ecx      // Save the value of ecx on the stack
[000013ca] [001123f3] [001023c7] 8b4508      mov eax,[ebp+08] // Load eax with argument to P
[000013cd] [001123ef] [000013c6] 50          push eax     // push 2nd argument to H onto the stack
[000013ce] [001123ef] [000013c6] 8b4d08      mov ecx,[ebp+08] // Load ecx with with argument to P
[000013d1] [001123eb] [000013c6] 51          push ecx     // push 1st argument to H onto the stack
[000013d2] [001123e7] [000013d7] e82ffdffff  call 00001106 // push return address; call simulated H
```

H: Infinitely Recursive Simulation Detected Simulation Stopped

```
[00001408] [0010235f] [00000000] 83c408      add esp,+08
[0000140b] [0010235b] [00000000] 50          push eax     // Push return value from H onto the stack
[0000140c] [00102357] [00000537] 6837050000  push 00000537 // Push address of "Input_Halts = " onto stack
[00001411] [00102357] [00000537] e870f1ffff  call 00000586 // call Output with its pushed arguments
Input_Halts = 0
[00001416] [0010235f] [00000000] 83c408      add esp,+08
[00001419] [0010235f] [00000000] 33c0       xor eax,eax  // set eax to 0
[0000141b] [00102363] [00000018] 5d          pop ebp
[0000141c] [00102367] [00000000] c3          ret         // return to 0 operating system
Number of Instructions Executed(987) == 15 Pages
```

## Example 04: An impossible program: Strachey(1965)

The Computer Journal, Volume 7, Issue 4, January 1965, Page 313,

<https://doi.org/10.1093/comjnl/7.4.313>

```
typedef void (*ptr)();
// rec routine P
// $L :if T[P] go to L
// Return $
void Strachey_P()
{
  L: if (T(Strachey_P)) goto L;
  return;
}

int main()
{
  output("Input_Halts = ", T(Strachey_P));
}
```

```
_Strachey_P()
[000012a6] (01) 55      push ebp
[000012a7] (02) 8bec     mov ebp,esp
[000012a9] (05) 68a6120000 push 000012a6
[000012ae] (05) e833fcffff call 00000ee6
[000012b3] (03) 83c404   add esp,+04
[000012b6] (02) 85c0     test eax,eax
[000012b8] (02) 7402     jz 000012bc
[000012ba] (02) ebcd     jmp 000012a9
[000012bc] (01) 5d       pop ebp
[000012bd] (01) c3       ret
Size in bytes:(0024) [000012bd]
```

```
_main()
[00001346] (01) 55      push ebp
[00001347] (02) 8bec     mov ebp,esp
[00001349] (05) 68a6120000 push 000012a6
[0000134e] (05) e893fbffff call 00000ee6
[00001353] (03) 83c404   add esp,+04
[00001356] (01) 50       push eax
[00001357] (05) 6817050000 push 00000517
[0000135c] (05) e805f2ffff call 00000566
[00001361] (03) 83c408   add esp,+08
[00001364] (02) 33c0     xor eax,eax
[00001366] (01) 5d       pop ebp
[00001367] (01) c3       ret
Size in bytes:(0034) [00001367]
```

machine address	stack address	stack data	machine code	assembly language
[00001346]	[0010221b]	[00000000]	55	push ebp
[00001347]	[0010221b]	[00000000]	8bec	mov ebp,esp
[00001349]	[00102217]	[000012a6]	68a6120000	push 000012a6
[0000134e]	[00102213]	[00001353]	e893fbffff	call 00000ee6

T: Begin Simulation Execution Trace Stored at:1122c7  
Address\_of\_T:ee6

```
[000012a6] [001122b7] [001122bb] 55      push ebp
[000012a7] [001122b7] [001122bb] 8bec     mov ebp,esp
[000012a9] [001122b3] [000012a6] 68a6120000 push 000012a6
[000012ae] [001122af] [000012b3] e833fcffff call 00000ee6
```

T: Infinitely Recursive Simulation Detected Simulation Stopped

**T knows its own machine address and on this basis it can easily examine its stored execution\_trace of Strachey\_P (see above) to determine:**

- Strachey\_P is calling T with the same arguments that T was called with.
- No instructions in Strachey\_P could possibly escape this otherwise infinitely recursive emulation.
- T aborts its emulation of Strachey\_P before its call to T is emulated.

```
[00001353] [0010221b] [00000000] 83c404   add esp,+04
[00001356] [00102217] [00000000] 50       push eax
[00001357] [00102213] [00000517] 6817050000 push 00000517
[0000135c] [00102213] [00000517] e805f2ffff call 00000566
```

Input\_Halts = 0

```
[00001361] [0010221b] [00000000] 83c408   add esp,+08
[00001364] [0010221b] [00000000] 33c0     xor eax,eax
[00001366] [0010221f] [00000018] 5d       pop ebp
[00001367] [00102223] [00000000] c3       ret
```

Number of Instructions Executed(538) == 8 Pages

## Example 05: P(P) halts because H(P,P) correctly determines that its input never halts

This conclusively proves that H(P,P) correctly simulates its input and that the behavior of the correctly simulated P is very different than the directly executed P(P).

The correctly simulated P cannot possibly terminate normally by reaching its own "return" instruction. The executed P does terminate normally and reaches its own "return" instruction.

If you are not an expert in the x86 language then you lack the basis to determine that the input to H(P,P) is not simulated correctly. The strongest claim that you can make is that on the basis that you do not understand the x86 language you do not understand the proof.

```
typedef void (*ptr)();
int H(ptr p, ptr i); // simulating halt decider
```

```
void P(ptr x)
{
    int Halt_Status = H(x, x);
    if (Halt_Status)
        HERE: goto HERE;
    return;
}
```

```
int main()
{
    P(P);
}
```

```
_P()
[0000143b](01) 55          push ebp
[0000143c](02) 8bec         mov ebp,esp
[0000143e](01) 51          push ecx
[0000143f](03) 8b4508      mov eax,[ebp+08]
[00001442](01) 50          push eax
[00001443](03) 8b4d08      mov ecx,[ebp+08]
[00001446](01) 51          push ecx
[00001447](05) e8affcffff  call 000010fb
[0000144c](03) 83c408      add esp,+08
[0000144f](03) 8945fc      mov [ebp-04],eax
[00001452](04) 837dfc00    cmp dword [ebp-04],+00
[00001456](02) 7402        jz 0000145a
[00001458](02) ebfe        jmp 00001458
[0000145a](02) 8be5        mov esp,ebp
[0000145c](01) 5d          pop ebp
[0000145d](01) c3          ret
Size in bytes:(0035) [0000145d]
```

```
_main()
[0000146b](01) 55          push ebp
[0000146c](02) 8bec         mov ebp,esp
[0000146e](05) 683b140000  push 0000143b
[00001473](05) e8c3ffffff  call 0000143b
[00001478](03) 83c404      add esp,+04
[0000147b](02) 33c0        xor eax,eax
[0000147d](01) 5d          pop ebp
[0000147e](01) c3          ret
Size in bytes:(0020) [0000147e]
```

machine address	stack address	stack data	machine code	assembly language
[0000146b]	[00102428]	[00000000]	55	push ebp
[0000146c]	[00102428]	[00000000]	8bec	mov ebp,esp
[0000146e]	[00102424]	[0000143b]	683b140000	push 0000143b
[00001473]	[00102420]	[00001478]	e8c3ffffff	call 0000143b
[0000143b]	[0010241c]	[00102428]	55	push ebp
[0000143c]	[0010241c]	[00102428]	8bec	mov ebp,esp
[0000143e]	[00102418]	[00000000]	51	push ecx
[0000143f]	[00102418]	[00000000]	8b4508	mov eax,[ebp+08]
[00001442]	[00102414]	[0000143b]	50	push eax
[00001443]	[00102414]	[0000143b]	8b4d08	mov ecx,[ebp+08]
[00001446]	[00102410]	[0000143b]	51	push ecx
[00001447]	[0010240c]	[0000144c]	e8affcffff	call 000010fb

// push P  
 // call P with argument on stack  
 // enter executed P  
 // load eax with argument to P  
 // push P from eax  
 // load ecx with argument to P  
 // push P from ecx  
 // call executed H with arguments on stack

```

H: Begin Simulation      Execution Trace Stored at:1124d
Address_of_H:10fb
[0000143b][001124c0][001124c4] 55      push ebp          // enter emulated P
[0000143c][001124c0][001124c4] 8bec     mov ebp,esp
[0000143e][001124bc][00102490] 51      push ecx
[0000143f][001124bc][00102490] 8b4508   mov eax,[ebp+08] // load eax with argument to P
[00001442][001124b8][0000143b] 50      push eax         // push P from eax
[00001443][001124b8][0000143b] 8b4d08   mov ecx,[ebp+08] // load ecx with argument to P
[00001446][001124b4][0000143b] 51      push ecx         // push P from ecx
[00001447][001124b0][0000144c] e8affcffff call 000010fb   // call emulated H with arguments on stack
H: Infinitely Recursive Simulation Detected Simulation Stopped

```

**When simulating halt decider H(P,P) simulates its input it can see that:**

- (1) Function H() is called from P().
- (2) With the same arguments to H().
- (3) With no instructions in P preceding its invocation of H(P,P) that could escape repeated simulations.

The above shows that the simulated P cannot possibly (reaches its "return" instruction and) terminate normally. H(P,P) simulates its input then P calls H(P,P) to simulate itself again. When H sees that this otherwise infinitely nested simulation would never end it aborts its simulation of P and rejects P as non-halting.

```

[0000144c][00102418][00000000] 83c408   add esp,+08     // return to executed P
[0000144f][00102418][00000000] 8945fc   mov [ebp-04],eax // load Halt_Status with return value
[00001452][00102418][00000000] 837dfc00 cmp dword [ebp-04],+00 // if Halt_Status == 0
[00001456][00102418][00000000] 7402     jz 0000145a    // goto 0000145a
[0000145a][0010241c][00102428] 8be5     mov esp,ebp
[0000145c][00102420][00001478] 5d      pop ebp
[0000145d][00102424][0000143b] c3      ret            // return from executed P to main
[00001478][00102428][00000000] 83c404   add esp,+04
[0000147b][00102428][00000000] 33c0     xor eax,eax    // set eax to 0
[0000147d][0010242c][00000018] 5d      pop ebp
[0000147e][00102430][00000000] c3      ret            // return from main to operating system
Number of Instructions Executed(998) == 15 Pages

```

## Appendix (Simulating halt decider applied to Peter Linz proof)

The following is the same idea as shown above this time it is applied to the Peter Linz Halting Problem proof. It can only be understood within the context of this proof.

A simulating halt decider (SHD) computes the mapping from its inputs to its own final states on the basis of the behavior of its correctly simulated input.

All of the conventional halting problem counter-example inputs are simply rejected by a simulating halt decider as non-halting because they fail to meet the Linz definition of halting:

**computation that halts ...** the Turing machine will halt whenever it enters a final state.  
(Linz:1990:234)

**USENET comp.theory: On 4/11/2022 3:19 PM, Malcolm McLean wrote:**

- > PO's idea is to have a simulator with an infinite cycle detector.
- > You would achieve this by modifying a UTM, so describing it as
- > a "modified UTM", or "acts like a UTM until it detects an infinite
- > cycle", is reasonable. And such a machine is a fairly powerful
- > halt decider. Even if the infinite cycle detector isn't very
- > sophisticated, it will still catch a large subset of non-halting
- > machines.

The following simplifies the syntax for the definition of the Linz Turing machine  $\hat{H}$ . There is no need for the infinite loop after  $H.qy$  because it is never reached. The halting criteria has been adapted so that it applies to a simulating halt decider (SHD).

$\hat{H}.q_0 \langle \hat{H} \rangle \vdash^* H \langle \hat{H} \rangle \langle \hat{H} \rangle \vdash^* \hat{H}.qy$

If the correctly simulated input  $\langle \hat{H} \rangle \langle \hat{H} \rangle$  to  $H$  would reach its own final state of  $\langle \hat{H}.qy \rangle$  or  $\langle \hat{H}.qn \rangle$ .

$\hat{H}.q_0 \langle \hat{H} \rangle \vdash^* H \langle \hat{H} \rangle \langle \hat{H} \rangle \vdash^* \hat{H}.qn$

If the correctly simulated input  $\langle \hat{H} \rangle \langle \hat{H} \rangle$  to  $H$  would never reach its own final state of  $\langle \hat{H}.qy \rangle$  or  $\langle \hat{H}.qn \rangle$ .

When  $\hat{H}$  is applied to  $\langle \hat{H} \rangle$  // subscripts indicate unique finite strings  
 $\hat{H}$  copies its input  $\langle \hat{H}_0 \rangle$  to  $\langle \hat{H}_1 \rangle$  then  $H$  simulates  $\langle \hat{H}_0 \rangle \langle \hat{H}_1 \rangle$

Then these steps would keep repeating: (unless their simulation is aborted)

$\hat{H}_0$  copies its input  $\langle \hat{H}_1 \rangle$  to  $\langle \hat{H}_2 \rangle$  then  $H_0$  simulates  $\langle \hat{H}_1 \rangle \langle \hat{H}_2 \rangle$

$\hat{H}_1$  copies its input  $\langle \hat{H}_2 \rangle$  to  $\langle \hat{H}_3 \rangle$  then  $H_1$  simulates  $\langle \hat{H}_2 \rangle \langle \hat{H}_3 \rangle$

$\hat{H}_2$  copies its input  $\langle \hat{H}_3 \rangle$  to  $\langle \hat{H}_4 \rangle$  then  $H_2$  simulates  $\langle \hat{H}_3 \rangle \langle \hat{H}_4 \rangle$ ...

Since we can see that the simulated input:  $\langle \hat{H}_0 \rangle$  to  $H$  would never reach its own final state of  $\langle \hat{H}_0.qy \rangle$  or  $\langle \hat{H}_0.qn \rangle$  we know that it is non-halting.

**Linz, Peter 1990.** An Introduction to Formal Languages and Automata. Lexington/Toronto: D. C. Heath and Company. (317-320) **this paper copyright 2022 by PL Olcott**