

Refuting the Sipser Halting Problem Diagonalization Argument

Every machine that halts in a reject state is a halting computation. When machine D is inserted into Figure 4.5 deriving Figure 4.6 the fact that a reject state is a halting computation is ignored. This makes the values at $\langle D, \langle M_1 \rangle \rangle$ and $\langle D, \langle M_1 \rangle \rangle$ in Figure 4.6 incorrect. When machine D is inserted into both Figure 4.4 and Figure 4.5 correctly (figures 4.4b and 4.5a respectively) the contradiction is eliminated.

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle \dots$
M_1	accept		accept	
M_2	accept	accept	accept	accept
M_3				
M_4	accept	accept		

Original Figure 4.4

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle \dots$
M_1	accept	~halt	accept	~halt
M_2	accept	accept	accept	accept
M_3	~halt	~halt	~halt	~halt
M_4	accept	accept	~halt	~halt

Figure 4.4a (converted from Figure 4.4 making ~halt assumption explicit)

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle \dots$
M_1	<u>accept</u>	reject	accept	reject
M_2	accept	<u>accept</u>	accept	accept
M_3	reject	reject	<u>reject</u>	reject
M_4	accept	accept	reject	<u>reject</u>

Original Figure 4.5 (underlining added)

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle \dots$	$\langle D \rangle \dots$
M_1	accept	~halt	accept	~halt	DC
M_2	accept	accept	accept	accept	DC
M_3	~halt	~halt	~halt	~halt	DC
M_4	accept	accept	~halt	~halt	DC
D	reject	reject	accept	accept	reject

Figure 4.4b (Insert D into Figure 4.4a)

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle \dots$	$\langle D \rangle \dots$
M_1	<u>accept</u>	reject	accept	reject	DC
M_2	accept	<u>accept</u>	accept	accept	DC
M_3	reject	reject	<u>reject</u>	reject	DC
M_4	accept	accept	reject	<u>reject</u>	DC
D	accept	accept	accept	accept	<u>accept</u>

Figure 4.5a (Insert D into Figure 4.5)

The above refutation of the Sipser diagonalization proof applies to all halting problem diagonalization proofs. Sipser was chosen as a widely available and very clear proof.

The Sipser proof was the basis for this superb lecture by Professor Dan Gusfield of UC Davis:

L15: Proof by Diagonalization that ATM (Halting Problem) is Not Decidable
Dec 12, 2012 <https://www.youtube.com/watch?v=jM6osxSX9GA>

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The following portions of pages 166-167 are directly relevant to the rebuttal.

Sipser, Michael 1997. Introduction to the Theory of Computation. Boston: PWS Publishing Company (165-167)

Where is the diagonalization in the proof of Theorem 4.9? It becomes apparent when you examine tables of behavior for TMs H and D . In these tables we list all TMs down the rows, M_1, M_2, \dots and all their descriptions across the columns, $\langle M_1 \rangle, \langle M_2 \rangle, \dots$. The entries tell whether the machine in a given row accepts the input in a given column. The entry is *accept* if the machine accepts the input but is blank if it rejects or loops on that input. We made up the entries in the following figure to illustrate the idea.

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle$	\dots
M_1	<i>accept</i>		<i>accept</i>		
M_2	<i>accept</i>	<i>accept</i>	<i>accept</i>	<i>accept</i>	
M_3					\dots
M_4	<i>accept</i>	<i>accept</i>			
\vdots			\vdots		

FIGURE 4.4

Entry i, j is *accept* if M_i accepts $\langle M_j \rangle$

In the following figure the entries are the results of running H on inputs corresponding to Figure 4.4. So if M_3 does not accept input $\langle M_2 \rangle$, the entry for row M_3 and column $\langle M_2 \rangle$ is *reject* because H rejects input $\langle M_3, \langle M_2 \rangle \rangle$.

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle$	\dots
M_1	<i>accept</i>	<i>reject</i>	<i>accept</i>	<i>reject</i>	
M_2	<i>accept</i>	<i>accept</i>	<i>accept</i>	<i>accept</i>	\dots
M_3	<i>reject</i>	<i>reject</i>	<i>reject</i>	<i>reject</i>	
M_4	<i>accept</i>	<i>accept</i>	<i>reject</i>	<i>reject</i>	
\vdots			\vdots		

FIGURE 4.5

Entry i, j is the value of H on input $\langle M_i, \langle M_j \rangle \rangle$

In the following figure, we added D to Figure 4.5. By our assumption, H is a TM and so is D . Therefore it must occur on the list M_1, M_2, \dots of all TMs. Note that D computes the opposite of the diagonal entries. The contradiction occurs at the point of the question mark where the entry must be the opposite of itself.

	$\langle M_1 \rangle$	$\langle M_2 \rangle$	$\langle M_3 \rangle$	$\langle M_4 \rangle$	\dots	$\langle D \rangle$	\dots
M_1	<u><i>accept</i></u>	<i>reject</i>	<i>accept</i>	<i>reject</i>		<i>accept</i>	
M_2	<i>accept</i>	<u><i>accept</i></u>	<i>accept</i>	<i>accept</i>	\dots	<i>accept</i>	\dots
M_3	<i>reject</i>	<i>reject</i>	<u><i>reject</i></u>	<i>reject</i>	\dots	<i>reject</i>	\dots
M_4	<i>accept</i>	<i>accept</i>	<i>reject</i>	<u><i>reject</i></u>		<i>accept</i>	
\vdots			\vdots		\ddots		
D	<i>reject</i>	<i>reject</i>	<i>accept</i>	<i>accept</i>		<u>?</u>	
\vdots			\vdots		\ddots		\ddots

FIGURE 4.6

If D is in the figure, a contradiction occurs at “?”