

Rebutting the Sipser Halting Problem Proof

MIT Professor Michael Sipser has agreed that the following verbatim paragraph is correct (he has not agreed to anything else in this paper):

If simulating halt decider H correctly simulates its input D until H correctly determines that its simulated D would never stop running unless aborted then H can abort its simulation of D and correctly report that D specifies a non-halting sequence of configurations.

A simulating halt decider computes the mapping from its input finite strings to an accept or reject state on the basis of the actual behavior specified by this input as measured by its correct simulation of this input. The following shows how the correct value for the D and $\langle D \rangle$ diagonal in Sipser's Figure 4.6 is accept.

We start with Sipser's definitions of H and D:

On input $\langle M, w \rangle$, where M is a TM and w is a string, H halts and accepts if M accepts w. Furthermore, H halts and rejects if M fails to accept w.

In other words, we assume that H is a TM, where

$$H(\langle M, w \rangle) = \begin{cases} \text{accept} & \text{if } M \text{ accepts } w \\ \text{reject} & \text{if } M \text{ does not accept } w \end{cases}$$

Now we construct a new Turing machine D with H as a subroutine. This new TM calls H to determine what M does when the input to M is its own description $\langle M \rangle$. Once D has determined this information, it does the opposite. That is, it rejects if M accepts and accepts if M does not accept.

$$D(\langle M \rangle) = \begin{cases} \text{accept} & \text{if } M \text{ does not accept } \langle M \rangle \\ \text{reject} & \text{if } M \text{ accepts } \langle M \rangle \end{cases} \quad (\text{Sipser 1997:165})$$

We encode the Sipser D and define the behavior of Sipser H as C functions.

```
//  
// sipser_H returns 1 when its input would halt and return 1  
// otherwise sipser_H returns 0  
//  
int sipser_D(int (*M)())  
{  
    if ( sipser_H(M, M) )  
        return 0;  
    return 1;  
}  
  
int main()  
{  
    output((char*)"Input_Halts = ", D(D));  
}
```

When H correctly simulates D it finds that D remains stuck in infinitely recursive simulation:

- (a) D calls H that simulates D with an x86 emulator
- (b) that calls a simulated H that simulates D with an x86 emulator
- (c) that calls a simulated H that simulates D with an x86 emulator ...

Until the executed H recognizes this repeating state, aborts its simulation of D and returns 0.

Complete halt deciding system (Visual Studio Project) Sipser version.

- (a) x86utm operating system
 - (b) x86 emulator adapted from libx86emu to compile under Windows
 - (c) Several halt deciders and their sample inputs contained within Halt7.c
 - (d) The execution trace of Sipser_H applied to Sipser_D is shown in Halt7_Sipser.txt
- https://liarparadox.org/2022_10_08.zip

D calls simulating halt decider H which computes the mapping from its input D to an accept or reject state on the basis of the behavior of its correct simulation of D. When H correctly determines that this simulated input would remain stuck in recursive simulation H aborts this simulation and reports non-halting by returning 0. When D reverses this decision it returns 1. This is used to correctly fill in the “?” in the Sipser Figure 4.6 (see below) with “accept”.

Simulating halt decider H recognizes instances of recursive simulation using the same criteria that it uses in its dynamic behavior pattern that recognizes infinite recursion:

```
void Infinite_Recursion(u32 N)
{
    Infinite_Recursion(N);
}
```

	⟨M ₁ ⟩	⟨M ₂ ⟩	⟨M ₃ ⟩	⟨M ₄ ⟩ ...	⟨D⟩ ...
M ₁	<u>accept</u>	reject	accept	reject	accept
M ₂	accept	<u>accept</u>	accept	accept	accept
M ₃	reject	reject	<u>reject</u>	reject	reject
M ₄	accept	accept	reject	<u>reject</u>	accept
...					
D	reject	reject	accept	accept	<u>?</u>
...					

Figure 4.6 (Sipser 1997:167)

Sipser, Michael 1997. Introduction to the Theory of Computation. Boston: PWS Publishing Company (165-167)

Appendix

```
int Sipser_D(int (*M)())
{
    if ( Sipser_H(M, M) )
        return 0;
    return 1;
}

int main()
{
    output((char*)"Input_Halts = ", D(D));
}
```

```
_Sipser_D()
[000012ae] 55          push ebp
[000012af] 8bec         mov ebp,esp
[000012b1] 8b4508       mov eax,[ebp+08]
[000012b4] 50          push eax
[000012b5] 8b4d08       mov ecx,[ebp+08]
[000012b8] 51          push ecx
[000012b9] e880fdffff   call 0000103e
[000012be] 83c408       add esp,+08
[000012c1] 85c0        test eax,eax
[000012c3] 7404        jz 000012c9
[000012c5] 33c0        xor eax,eax
[000012c7] eb05        jmp 000012ce
[000012c9] b801000000   mov eax,00000001
[000012ce] 5d          pop ebp
[000012cf] c3          ret
Size in bytes:(0034) [000012cf]
```

Sipser_H: Begin Simulation Execution Trace Stored at:111fa8

machine address	stack address	stack data	machine code	assembly language
[000012ae]	[00111f94]	[00111f98]	55	push ebp // Begin Sipser_D
[000012af]	[00111f94]	[00111f98]	8bec	mov ebp,esp
[000012b1]	[00111f94]	[00111f98]	8b4508	mov eax,[ebp+08]
[000012b4]	[00111f90]	[000012ae]	50	push eax // push Sipser_D
[000012b5]	[00111f90]	[000012ae]	8b4d08	mov ecx,[ebp+08]
[000012b8]	[00111f8c]	[000012ae]	51	push ecx // push Sipser_D
[000012b9]	[00111f88]	[000012be]	e880fdffff	call 0000103e // call Sipser_H

Sipser_H: Infinitely Recursive Simulation Detected Simulation Stopped

We can see that the first seven instructions of Sipser_D simulated by Sipser_H precisely match the first seven instructions of the x86 source-code of Sipser_D. This conclusively proves that these instructions were simulated correctly.

Anyone sufficiently technically competent in the x86 programming language will agree that the above execution trace of Sipser_D simulated by Sipser_H shows that Sipser_D will never stop running unless Sipser_H aborts its simulation of Sipser_D.