

The Design of a Formal System of Science: A Compelling Foundation for Physics

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Abstract

A formal system of science is presented as a candidate for a more powerful foundation of physics. Rather than starting with an axiomatic theory of physics as is typically done (quantum mechanics, general relativity, classical mechanics, etc), we instead had the idea to start at the level of science by designing a formal system of science. Our goal was then to attempt, as a challenge, to derive the laws of physics as a theorem of the scientific method. This exercise turned out to be quite the cleansing exercise; correcting definitions and resolving open problems spawning philosophy, mathematics, science (philosophy of) and physics. Some of our notable results include a definition of experimental data (done purely mathematically), a definition of the observer (again done purely mathematically), and a derivation of the laws of physics as the product of the scientific method (and thus, we allege, done uncontroversially). This group of results constitutes the primary items of modern theoretical physics currently missing mathematical formalisation. The first part of the paper consists of constructing an experimental basis for pure mathematics. For this we employ halting programs. Since halting programs are arbitrarily complex and subject to the halting problem, we leverage modern notions of mathematical incompleteness to create a formulation of mathematics conducive to experimental methods in which one must conduct halting experiments to make progress. In this setup, the observer is defined as a measure space over all possible experiments and maximizing the entropy of its measure produces a 'quantum-computation superset' of quantum physics that supports quantum gravity. Finally, testable predictions are proposed.

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1 Introduction

In classical philosophy an axiom is a statement which is self-evidently true such that it is accepted without controversy or question. But this definition has been retired in modern usage. Any so-called "self-evident" axiom can also be posited to be false and either choice of its truth-value yields a different model; the archetypal example being the parallel line postulate of Euclid, allowing for hyperbolic/spherical geometry when it is false. Consequently, in modern logic an axiom is demoted to simply be a starting point for a premise, and in mathematics an axiom is a sentence of a language that is held to be true by definition.

A long standing goal of philosophy has been to find necessarily true principles that could be used as the basis of knowledge. For instance, the universal doubt method of Descartes had such a goal in mind. The 'justified true belief' theory of epistemology is another attempt with a similar goal. But, so far, all such attempts have flaws and loopholes, the elimination of which is assumed, at

best, to reduce the theory to a handful of statements, rendering it undesirable as a foundation to all knowledge.

In epistemology, the Gettier problem[1] is a well known objection to the belief that knowledge is that which is both true and justified, relating to a family of counter-examples. All such counter-examples rely on the same loophole: if the justification is not 'air-tight' then there exists a case where one is right by pure luck, even if the claim were true and believed to be justified. For instance, if one glances at a field and sees a shape in the form of a dog, one might think he or she is justified in the belief that there is a dog in the field. Now suppose there is a dog elsewhere in the field, but hidden from view. The belief "there is a dog in the field" is now justified and true, but it is not knowledge because it is only true by pure luck.

Richard Kirkham[2] proposed to add the criteria of infallibility to the justification. What used to be "justified true belief" would now become "infallible true belief". This eliminates the loophole, but it is an unpopular solution because adding it is assumed to reduce epistemology to radical skepticism in which almost nothing is knowledge, thus rendering epistemology non-comprehensive.

Here, we will adopt the insight of Kirkham regarding the requirement of infallibility, whilst also resolving the non-comprehensiveness objection. To do so, we will structure our statements such that they are individually infallible, yet as a group form a Turing complete language. Turing completeness will guarantee comprehensiveness. Our tool of choice will be halting programs. As we will see, halting programs carry all desired features to make this possible. Using them, we will be able to tackle epistemological knowledge using infallible statements, as put forward by Kirkham. That may sound impressive, but there is a catch. Although we can describe the knowledge once it is acquired, acquiring it in the first place will be difficult, in some case even arbitrarily difficult. After-all, halting programs are of course subject to the halting problem. Indeed, we still have to identify those program that halts from those that do not, and because of the halting problem, as there exists no general algorithm able to do so, the system will be irreducibly experimental (more on that later).

With these two features, the sum-total of epistemological knowledge becomes well-defined and is in fact recursively enumerable. The implications of such a construction should not be underestimated as this was a long standing goal of the ancients of philosophy, even if attempts to crack it have generally fallen out of favour in modern times. In any case, our formal system of science will thoroughly exploit this construction.

Finally, let us state that attempts to find a logical basis for epistemology have been made ad nauseam in the past but they failed for primarily two reasons. First, they were attempted before Gödel-type theorems were known and appreciated, and attempts were directed at constructing *decidable* logical bases. Secondly, instead of directing efforts to recursively enumerable bases following the discovery of said theorems, efforts simply felt out of favour. However, it is possible to construct a logical basis for epistemology provided that such basis is recursively enumerable (and not decidable), and further the limitations induced by recursive enumeration ought to instead be seen as an *opportunity*; in this

case, to create a formal system to map out epistemology, such that it may serve as the foundation to a formalization of science.

For more general information regarding the connection between mathematics, science and programs, we recommend the works of Gregory Chaitin[3, 4, 5]. A familiarity with his work is assumed. Let us now continue.

1.1 Halting Programs as Knowledge

How do we construct an infallible statement, so that it qualifies as an epistemological statement in the sense of Kirkham?

Let us take the example of a statement that may appear as an obvious true statement such as " $1 + 1 = 2$ ", but is in fact not infallible. Here, we will provide the correct definition of an infallible statement, but equally important, such that the set of all such statements is Turing complete, thus forming a language of maximum expressive power.

Specifically, the sentence " $1 + 1 = 2$ " halts on some Turing machine, but not on others and thus is not infallible. Instead consider the sentence $PA \vdash [s(0) + s(0) = s(s(0))]$ to be read as "Peano's axioms prove that $1 + 1 = 2$ ". Such a statement embeds as a prefix the set of axioms in which it is provable. One can deny that $1 + 1 = 2$ (for example, an adversary could claim binary numbers, in which case $1 + 1 = 10$), but if one specifies the exact axiomatic basis in which the claim is provable, said adversary would find it harder to find a loophole to fail the claim. Nonetheless, even with this improvement, an adversary can fail the claim by providing a Turing machine for which $PA \vdash [s(0) + s(0) = s(s(0))]$ does not halt.

If we use the tools of theoretical computer science we can produce statements free of all loopholes, thus ensuring they are infallible. Those statements become programs:

Definition 1 (Halting Program). *Let \mathbb{L} be the set of all sentences with alphabet Σ . An halting program f is a pair (TM, p) of sentences from $\mathbb{L} \times \mathbb{L}$ such that a universal Turing machine UTM halts for it:*

$$\text{iff } UTM(TM, p) \text{ halts, then } f = (TM, p) \text{ is a halting program} \quad (1)$$

A universal Turing machine UTM which takes a Turing machine TM and a sentence p as inputs, will halt if and only if p halts on TM. Thus a claim that p halts on TM, if true, is an halting program because it is verifiable on all universal Turing machines.

The second objection is that the infallibility requirement is too strong making makes epistemology non-comprehensive, only able to tackle a handful of statements. However, the set of all halting programs constitutes the entire domain of the universal Turing machine, and thus the expressive power of halting programs must be on par with any Turing complete language. Since there exists no greater expressive power for a formal language than that of Turing completeness, then no reduction takes place. The resulting construction is both element-wise infallible, and comprehensive as a set.

1.2 Halting Manifests

We will now use halting programs to redefine the foundations of mathematics in terms of *halting manifests*, replacing *formal axiomatic systems*.

Before we do so, let us build up the intuition. In principle, one can use any Turing complete structure to re-define all of mathematics. The task is not particularly difficult but the work can nonetheless be substantial. One generally has to build a translator between the two formulation, whose existence is interpreted as a proof of equivalence. For instance, one can write all of mathematics using the english language (if one were so included), or with using set theory (with arbitrary set equipment), or category theory, or using a computer language such as c++, or using arithmetic with multiplication, etc. If the language is Turing complete, then it is as expressive as any other Turing complete language, and a translator is guaranteed to exist. So why pick a particular basis over another? This is often due to other conveniences and constraints than pure expressive power. For instance, sets allow us to intuitively express a very large class of mathematical problems quite conveniently. Typical selection criterions are: can we express the problem at hand clearly?, elegantly?, are the solutions also clear and easier to formulate, than in alternative languages?

Here we will use and introduce the halting manifest formulation of mathematics, and, as we will see, its advantages are stunning. A halting manifest comprises a group of programs known to halt, and this group of programs defines a specific instance of mathematical knowledge.

Definition 2 (Halting Manifest). *Let $\mathbb{D} = \text{Dom}(\text{UTM})$ be the set of all halting programs for a given universal Turing machine. A halting manifest, or simply a manifest, \mathbf{m} of n halting program is an element of the n -fold Cartesian product of \mathbb{D} :*

$$\mathbf{m} \in \mathbb{D}^n, \quad \mathbf{m} := ((\text{TM}_1, p_1), \dots, (\text{TM}_n, p_n)) \quad (2)$$

The tuple, in principle, can be empty $\mathbf{m} := ()$, finite $n \in \mathbb{N}$ or countably infinite $n = \infty$.

- *A manifest contains some, but not necessarily all, programs of the domain of the universal Turing machine.*

The programs comprising a manifest replace the normal role of both axioms and theorems and instead form a single verifiable atomic concept constituting a unit of mathematical knowledge. Let us explicitly point out the difference between the literature definition of a formal system and ours: for the former, its theorems are a subset of the sentences of \mathbb{L} provable from the axioms — whereas for a manifest, its elements are pairs of $\mathbb{L} \times \mathbb{L}$ which halts on a UTM.

Let us now explore some of the advantages of using manifests versus formal axiomatic systems. Manifests are more conducive to a description of the scientific process, including the accumulation of experimental data, than formal axiomatic systems are. Let us take an example. Suppose we wish to represent in

real-time, and with live updates, the set of all knowledge produced by a group of 50,000ish (and growing!) mathematicians working in a decentralized manner (perhaps from their offices) over the course of at least many decades, and perhaps even for an indefinite amount of time into the future. Some of the work they produced may build on each others', but it will also be the case that part of their work is incompatible. For instance, some might find contradictions in their assumptions and abandon large segments of their work. As one learns primarily from his or her errors, we may wish to catalogue these contradictions for posterity. Let us first try with formal axiomatic systems. Finding the 'correct' and singular formal axiomatic system to describe the totality of what they have discovered, including abandoned work and contradictions, will be quite a challenge. One challenge occurs whenever a new contradiction is found, as one would need to further isolate it within a wrapper of para-consistent logic, before inclusion within the all-encompassing formal axiomatic system. Another challenge occurs when mathematicians invent new, possibly more elegant, axiomatic basis outright. One would constantly need to adjust his or her proposed all-encompassing formal axiomatic system to account for new discoveries as they are made. Such an axiomatic basis would eventually grow to an unmaintainable level, not unlike the spaghetti codes of the early days of software engineering. And we have not even mentioned the problems spawned by general incompleteness theorems such as those of Gödel and Gregory Chaitin, and the negative resolution to Hilbert's second problem! What if someone proves a statement that is not provable from the selected axiomatic basis; in this case re-adjustments are perpetually necessary. As mathematicians are a creative bunch, one would never be able to settle on a final axiomatic system as they could always decide to explore a sector of mathematical space not covered by the current system. Comparatively, using a manifest, the task is much easier: One simply need to push each new discovery at the end of the manifest; no adjustment is ever required after insertion, we never run out of space, and halting programs do not undermine each other even if they represent a contradiction. A manifest is conceptually similar to an empirical notebook of raw mathematical knowledge.

Formal axiomatic systems do not excel at pure description because they are be more akin to an *interpretation* of mathematical knowledge based on a preference of some patterns or tools (we like sets, thus ZFC!, or we prefer categories, thus category theory!). New knowledge will eventually force one to change this preference, or at least to consider alternative preferences. Not so with manifests! Manifests are the true starting point of the logical inquiry as they represent an infallible and non-preferential description of mathematical knowledge.

We will now explore the concept more rigorously.

Note on the notation: we will designate $f_i = (\text{TM}_i, p_i)$ as an halting program element of \mathbf{m} , and $\text{proj}_1(f_i)$ and $\text{proj}_2(f_i)$ designate the first and second projection of the pair f_i , respectively. Thus $\text{proj}_1(f_i)$ is the TM_i associated with f_i , and $\text{proj}_2(f_i)$ is the input p_i associated with f_i . If applied to a tuple or set of pairs, then $\text{proj}_1(\mathbf{m})$ returns the set of all p in \mathbf{m} and $\text{proj}_2(\mathbf{m})$ returns the set of all TM in \mathbf{m} .

Theorem 1 (Incompleteness Theorem). *Let \mathbf{m} be a manifest. If $\mathbf{m} = \text{Dom}(\text{UTM})$, then \mathbf{m} is recursively enumerable (and non-decidable). The proof follows from the domain of a universal Turing machine being non-decidable.*

Definition 3 (Premises). *Let \mathbf{m} be a manifest. The premises in \mathbf{m} are defined as the set of all TM in \mathbf{m} :*

$$\mathbb{P} := \text{proj}_1(\mathbf{m}) \quad (3)$$

Definition 4 (Theorems). *Let \mathbf{m} be a manifest. The theorems of \mathbf{m} are defined as the set of all p in \mathbf{m} :*

$$\mathbb{T} := \text{proj}_2(\mathbf{m}) \quad (4)$$

Definition 5 (Spread (of a theorem)). *The set of all premises in \mathbf{m} in which a theorem is repeated is called the spread of the theorem. For instance if $\mathbf{m} = (\text{TM}_1, p_1), (\text{TM}_2, p_1)$, then the spread of p_1 is $\{\text{TM}_1, \text{TM}_2\}$.*

Definition 6 (Scope (of a premise)). *The set of all theorems in \mathbf{m} in which a premise is repeated is called the scope of the premise. For instance if $\mathbf{m} = \{(\text{TM}_1, p_1), (\text{TM}_1, p_2)\}$, then the scope of TM_1 is $\{p_1, p_2\}$.*

1.2.1 Connection to Formal Axiomatic Systems

We can, of course, connect our halting-manifest formulation to the standard formal axiomatic system (FAS) formulation:

Definition 7 (Enumerator (of a FAS)). *A function $\text{enumerator}_{\text{FAS}}$ is an enumerator for FAS if it recursively enumerates the theorems of FAS. For instance:*

$$\text{enumerator}_{\text{FAS}}(s) = \begin{cases} 1 & \text{FAS} \vdash s \\ \#/\text{does-not-halt} & \text{otherwise} \end{cases} \quad (5)$$

Definition 8 (Domain (of FAS)). *Let FAS be a formal axiomatic system, let \mathbf{m} be a manifest and let $\text{enumerator}_{\text{FAS}}$ be a function which recursively enumerates the theorems of FAS. Then the domain of FAS, denoted as $\text{Dom}(\text{FAS})$, is the set of all sentences $s \in \mathbb{L}$ which halts for $\text{enumerator}_{\text{FAS}}$.*

Definition 9 (Formal Axiomatic Representation). *Let FAS be a formal axiomatic system, let \mathbf{m} be a manifest and let $\text{enumerator}_{\text{FAS}}$ be a function which recursively enumerates the theorems of FAS. Then FAS is a formal axiomatic representation of \mathbf{m} iff:*

$$\text{Dom}(\text{FAS}) = \text{proj}_2(\mathbf{m}) \quad (6)$$

Definition 10 (Factual Isomorphism). *Two formal axiomatic systems FAS_1 and FAS_2 are factually-isomorphic if and only if $\text{Dom}(\text{FAS}_1) = \text{Dom}(\text{FAS}_2)$.*

1.2.2 Axiomatic Information

The first result of interest will be the introduction of axiomatic information. We can understand the elements of any particular manifest as having been 'picked', in some sense, from the set of all possible halting programs. If the pick is random and described as a probability measure ρ , we can quantify the information of the pick using the entropy.

Definition 11 (Axiomatic Information). *Let $\mathbb{D} = \text{Dom}(\text{UTM})$ be the domain of a universal Turing machine. Then, let $\rho : \mathbb{D} \rightarrow [0, 1]$ be a probability measure over \mathbb{D} . Finally, let \mathbf{m} be a manifest subset of \mathbb{D}^n . The axiomatic information of a single element of \mathbf{m} is quantified as the entropy of ρ :*

$$S = - \sum_{q \in \mathbb{D}} \rho(q) \ln \rho(q) \quad (7)$$

For instance, a well-known (non-computable) probability measure regarding a sum of prefix-free programs is the Halting probability[6] of computer science:

$$\Omega = \sum_{p \in \text{Dom}(\text{UTM})} 2^{-|p|} \implies \rho(p) = 2^{-|p|} \quad (8)$$

The quantity of axiomatic information of a given manifest (and especially its maximization), rather than any particular set of axioms, will be the primary quantity of interest for the production of a maximally informative theory in this framework. A strategy to gather mathematical knowledge which picks halting programs according to the probability measure which maximizes the entropy will be a maximally informative strategy.

1.3 Discussion — The Mathematics of Knowledge

Each element of a manifest is a program-input pair representing an algorithm which is known to produce a specific result. Let us see a few examples.

How does one know how to tie one's shoes? One knows the algorithm required to produce a knot in the laces of the shoe. How does one train for a new job? One learns the internal procedures of the shop, which are known to produce the result expected by management. How does one impress management? One learns additional skills outside of work and applies them at work to produce results that exceed the expectation of management. How does one create a state in which there is milk in the fridge? One ties his shoes, walks to the store, pays for milk using the bonus from his or her job, then brings the milk back home and finally places it in the fridge. How does a baby learn about object permanence? One plays peak-a-boo repeatedly with a baby, until it ceases to amuse the baby — at which point the algorithm which hides the parent, then shows him or her again, is learned as knowledge. How does one untie his shoes? One simply pulls on the tip of the laces. How does one untie

his shoes if, after partial pulling, the knot accidentally tangles itself preventing further pulling? One uses his fingers or nails to untangle the knot, and then tries pulling again.

Knowledge can also be in more abstract form — for instance in the form of a definition that holds for a special case. How does one know that a specific item fits a given definition of a chair? One iterates through all properties referenced by the definition of the chair, each step confirming the item has the given property — then if it does for all properties, it is known to be a chair according to the given definition.

In all cases, knowledge is an algorithm along with an input, such that the algorithm halts for it, lest it is not knowledge. The set of all known pairs forms a manifest.

Let us consider a few edge cases. What if a manifest contains both "A" and "not A" as theorems? For instance, consider:

$$\mathbf{m} := ((\text{TM}_1, A), (\text{TM}_1, \neg A)) \quad (9)$$

Does allowing contradictions at the level of the theorems of \mathbf{m} create a problem? Should we add a few restrictions to avoid this unfortunate scenario? Let us try an experiment to see what happens — specifically, let me try to introduce $A \wedge \neg A$ into my personal manifest, and then we will evaluate the damage I have been subjected to by this insertion. Consider the following program TM_1 :

1. If (p="A" or p="not A") then
2. return 1;
3. else (loop())

It thus appears that I can have knowledge that the above program halts for both "A" and "not A" and still survive to tell the tale. A-priori, the sentences "A" and "not A" are just symbols. Our reflex to attribute the law of excluded middle to these sentences requires the adoption of a deductive system. This occurs one step further at the selection of a specific formal axiomatic representation of the manifest, and not at the level of the manifest itself.

The only inconsistency that would create problems for this framework would be a proof that a given halting program both [HALTS] and [NOT HALTS] on a UTM. By definition of a UTM, this cannot happen lest the machine was not a UTM to begin with. Thus, we are expected to be safe from such contradictions.

Now, suppose one has a sizeable manifest which may contain a plurality of pairs:

$$\mathbf{m} := ((\text{TM}_1, p_1), (\text{TM}_2, \neg p_1), (\text{TM}_1, p_2), (\text{TM}_2, p_1), (\text{TM}_2, \neg p_3)) \quad (10)$$

Here, the negation of some, but not all, is also present across the pairs: in this instance, the theorems p_1 and p_3 are negated but for different premises. What

interpretation can we give to such elements of a manifest? For our example, let us call the sentences p_1, p_2, p_3 the various flavours of ice cream. It could be that the Italians define ice cream in a certain way, and the British define it in a slightly different way. Recall that halting programs are pairs which contain a premise and a theorem. The premise contains the 'definition' under which the flavour qualifies as real ice cream. A flavour with a large spread is considered real ice cream by most definitions (i.e. vanilla or chocolate ice cream), and one with a tiny spread would be considered real ice cream by only very few definitions (i.e. tofu-based ice cream). Then, within this example, the presence of p_1 and its negation simply means that tofu-based ice cream is ice cream according to one definition, but not according to another.

Reality is of a complexity such that a one-size-fits-all definition does not work for all concepts, and further competing definitions might exist: a chair may be a chair according to a certain definition, but not according to another. The existence of many definitions for one concept is a part of reality, and the mathematical framework which correctly describes its manifest ought to be sufficiently flexible to handle this, without itself exploding into a contradiction.

Even in the case where both A and its negation $\neg A$ were to be theorems of \mathbf{m} while also having the same premise, is still knowledge. It means one has verified that said premise is inconsistent. One has to prove to oneself that a given definition is inconsistent by trying it out against multiple instances of a concept, and those 'trials' are all part of the manifest.

1.4 Discussion — Idempotency, and Infinite Regress

Suppose one defines a quantum theory referencing a number of axioms; perhaps both a) the Born rule and b) the wave-function collapse are axiomatic by definition. Now, if one is curious, one may naturally wonder: "is there a more fundamental explanation for these axioms?". Now suppose one successfully invents an axiomatic non-local hidden variable theory which does derive those rules, allegedly, from "*first principles*". This may provide us with a small dopamine hit at first, but surely the next generation will eventually begin to wonder the same: "is there a more fundamental explanation for these axioms?". The cycle simply repeats itself — we ask rhetorically; have we truly made progress?

In general all formal axiomatic systems are subject to a perpetual realization of the infinite regress problem. The reason is simply that any FAS is logically entailed by infinitely-many other FAS, thus one is free to perpetually replace a given FAS by another, believed to be more fundamental, FAS. Even if one were to invent an axiomatic basis able to derive a) and b) as theorems, and somehow show it to be an equivalent representation of quantum mechanics, then the new construction will import its axiomatic basis, and said axiomatic basis will remain the subject of the infinite regress problem. One may perhaps fool himself or herself into believing a more fundamental theory has been achieved, however one has merely identified a new element out the set of all equivalent FAS for a given state of knowledge.

Now, let us try with halting manifests and compare. Can one, as a challenge,

produce a "more fundamental" halting manifest for a given manifest? Before we investigate, let us note that a manifest uniquely represents a specific instance of mathematical knowledge. To change the manifest with another, entails to represent a different instance of mathematical knowledge. For instance, say one has $\mathbf{m}_a := (f_1, f_2)$. Now suppose one wishes to produce a 'more fundamental' manifest of \mathbf{m}_a , calling it \mathbf{m}_b . Well either $\mathbf{m}_a = \mathbf{m}_b$ at which point neither are more or less fundamental (they are the same), or $\mathbf{m}_a \neq \mathbf{m}_b$ at which point they describe different system. We thus say that halting manifests are *idempotent* representations of knowledge.

The relationships and differences can be visualized by noting that a manifest admits infinitely-many possible formal axiomatic representations, whereas a formal axiomatic system admits a single manifest as a representation of its domain. Idempotency thus occurs at and only at the level of the manifest representation of mathematics. Finally, the infinite regress problem immediately terminates for knowledge-bases due to this idempotency.

2 The Formal System of Science

We now assign to our halting-manifest formulation of mathematics the interpretation of a purely mathematical system of science.

2.1 Axiomatic Foundation of Science

The fundamental object of study of science is not the electron, the quark or even super-strings, but the reproducible experiment. An experiment represents an 'atom' of experimentally verifiable knowledge.

Definition 12 (Experiment). *Let $f = (\text{TM}, p)$ be a pair comprised of two sentences of a language \mathbb{L} . The first sentence, TM, is called the protocol. The second sentence, p , is called the hypothesis. Let UTM be a universal Turing machine. If $\text{UTM}(\text{TM}, p)$ halts then the pair (TM, p) is said to be an experiment. In this case, we say that the protocol verifies the hypothesis. If $\text{UTM}(\text{TM}, p)$ does not halt, we say that the pair fails to verify the hypothesis.*

$$\text{UTM}(\text{TM}, p) \begin{cases} \text{HALT} & \implies \text{the experiment verifies } p \\ \neg \text{HALT} & \implies \text{the pair fails verification} \end{cases} \quad (11)$$

Of course, in the general case, as per the halting problem there exists no decidable function which can determine which pair is an experiment and which pair fails verification. In the general case, one must try them out to see which one halts — this is why they are called experiments.

Thesis 1 (Formal Equivalence Thesis). *As you would notice, our definition of experiments is the same as that of halting programs. Thus, we have merely rebranded halting programs with a new name. The formal equivalence thesis is*

a claim that the set of all possible experiments (including physical experiments) and the set of all halting programs are the same. The thesis is proven by its first proposition; formal reproducibility which necessitates this equivalence. The next propositions provide insight and help to fix the intuition.

Proposition 1 (Formal Reproducibility). *Experiments are formally reproducible.*

Proof. Let UTM and UTM' each be a universal Turing machine. For each pair UTM(TM, p) which halts on UTM, there exists a computable function, called an encoding function, which maps said pairs as $encode(TM, p) \rightarrow (TM', p')$ such that (TM', p') halts for UTM'. The existence of such function is guaranteed by (and equivalent to) the statement that any UTM can simulate any other. \square

An experiment, so defined, is formally reproducible. I can transmit, via fax or other telecommunication medium, the pair (TM, p) to another experimentalist, and I would know with absolute certainty that he or she has everything required to reproduce the experiment to perfection.

In the peer-reviewed literature, the typical requirement regarding the reproducibility of an experiment is that an expert of the field be able to reproduce the experiment, and this is of course a much lower standard than formal reproducibility which is a mathematically precise definition. Here, for the protocol TM to be a Turing machine, the protocol must specify all steps of the experiment including the complete inner workings of any instrumentation used for the experiment. The protocol must be described as an effective method equivalent to an abstract computer program. Should the protocol fail to verify the hypothesis, the entire experiment (that is the group comprising the hypothesis, the protocol and its complete description of all instrumentation) is rejected. For these reasons and due to the generality of the definition, I conjecture that the above definition is the only possible definition of the experiment that is formally reproducible (as opposed to say "sufficiently reproducible for practical purposes"), yet universal as a Turing complete language.

Formal reproducibility supports the thesis as follows: only halting programs are formally reproducible. Thus, by necessity, if a physical experiment is to be formally reproducible, and only halting programs are formally reproductive, then it must be equivalent to a halting program.

Proposition 2 (Formal Falsifiability). *To falsify a scientific theory one must produce a solution (to its equations) or a theorem (to its axioms). Both the product of a solution or the proof of a theorem are mathematical knowledge and are thus representable as halting programs.*

For instance, one can equate the set of all (non-relativistic) quantum mechanical experiments to the set of all solutions of the Schrödinger equation. The latter are programs (i.e. numerical solutions to the equation are typically obtained by computers, or by hand), and they ought to map one-to-one with the former, lest the equation is falsified. This correspondance of course must hold true for any theory to be qualified as a non-falsified scientific theory.

Proposition 3 (Recursive Enumeration in Nature). *If one can construct a computer in nature and instruct it to recursively enumerate halting programs, then by necessity an experimentally-complete physical theory (which would have account for the inner-workings of said computer) must share its domain with that of the halting programs as they are Turing complete. Nature can, at most, enumerate all halting programs.*

This is of course not yet the laws of physics (Axiom 1 will be required for that) but it is nonetheless a necessary mapping between halting programs and physical experiments, enforced strictly by us requiring science to be falsifiable and reproducible. If one holds the philosophical opinion that qualifying nature as a tuple of experiments (which are in turn halting programs) is too restrictive to be palatable, then one's only way out is to relax falsifiability or reproducibility by allowing not falsifiable elements in a theory, or irreproducible experiments in nature... an import which, to us, seems even more unpalatable. Finally, please note that the apparent reductionism to the halting programs is just that: apparent — we have already shown them to be epistemologically comprehensive and thus no reduction in *knowledge* takes place as a result of the equivalence thesis. One must appreciate what Turing completeness entails to realize that comprehensiveness is necessarily maintained.

We now introduce a few more definitions to complete the formal system of science.

Definition 13 (Domain of science). *We note \mathbb{D} as the domain (Dom) of science. We can define \mathbb{D} in reference to a universal Turing machine UTM as follows:*

$$\mathbb{D} := \text{Dom}(\text{UTM}) \quad (12)$$

Thus, for all pairs of sentences (TM, p) , if $\text{UTM}(\text{TM}, p)$ halts, then $(\text{TM}, p) \in \mathbb{D}$. It follows that all experiments are elements of the domain of science.

Definition 14 (Experimental Space). *Let \mathbf{m} be a manifest comprised of n halting programs, and let $\mathbb{M} = \bigcup_{i=1}^n \text{proj}_i(\mathbf{m})$ be the set comprised of the halting programs of \mathbf{m} . The experimental space \mathbb{E} of \mathbf{m} is the "powertuple" of \mathbf{m} :*

$$\mathbb{E} := \bigcup_{i=0}^n (\mathbb{M})^i \quad (13)$$

- *Put simply, experimental space is the set of all possible manifests (including the empty manifest).*
- *Conceptually, a powertuple is similar to a powerset where the notion of the set is replaced by that of the tuple.*
- *All elements of an experimental space are manifests, and all "sub-tuples" of a manifest are elements of its experimental space.*

Definition 15 (Scientific method). *An algorithm which recursively enumerates experiments, is called a scientific method.*

Theorem 2 (Scientific method (Existence of)). *Existence of the scientific method.*

Proof. Consider a dovetail program scheduler which works as follows.

1. Sort all pairs of sentences of $\mathbb{L} \times \mathbb{L}$ in shortlex. Let the ordered pairs $(TM_1, p_1), (TM_2, p_1), (TM_1, p_2), (TM_2, p_2), (TM_3, p_1), \dots$ be the elements of the sort.
2. Take the first element of the sort, $UTM(TM_1, p_1)$, then run it for one iteration.
3. Take the second element of the sort, $UTM(TM_2, p_1)$, then run it for one iteration.
4. Go back to the first element, then run it for one more iteration.
5. Take the third element of the sort, $UTM(TM_1, p_2)$, then run it for one iteration.
6. Continue with the pattern, performing iterations one by one, with each cycle adding a new element of the sort.
7. Make note of any pair (TM_i, p_j) which halts.

This scheduling strategy is called dovetailing and allows one to enumerate the domain of a universal Turing machine recursively, without getting stuck by any singular program that may not halt. Progress will eventually be made on all programs... thus producing a recursive enumeration. \square

Dovetailing is of course a simple/non-creative approach to the scientific method. The point here was only to show existence of such an algorithm, not to find the optimal one.

2.1.1 Classification of Scientific Theories

Definition 16 (Scientific Theory). *Let \mathbf{m} be a manifest and let ST be a formal axiomatic system. If*

$$\text{proj}_2(\mathbf{m}) \cap \text{Dom}(\text{ST}) \neq \emptyset \quad (14)$$

then ST is a scientific theory of \mathbf{m} .

Definition 17 (Empirical Theory). *Let \mathbf{m} be a manifest and let ST be a scientific theory. If*

$$\text{proj}_2(\mathbf{m}) = \text{Dom}(\text{ST}) \quad (15)$$

then ST is an empirical theory of \mathbf{m} .

Definition 18 (Scientific Field). *Let \mathbf{m} be a manifest and let ST be a scientific theory. If*

$$\text{Dom}(\text{ST}) \subset \text{proj}_2(\mathbf{m}) \quad (16)$$

then ST is a scientific field of \mathbf{m} .

Definition 19 (Predictive Theory). *Let \mathbf{m} be a manifest and let ST be a scientific theory. If*

$$\text{proj}_2(\mathbf{m}) \subset \text{Dom}(\text{ST}) \quad (17)$$

then ST is a predictive theory of \mathbf{m} .

Specifically, the predictions of ST are given as follows:

$$\mathbb{S} := \text{Dom}(\text{ST}) \setminus \mathbf{m} \quad (18)$$

Scientific theories that are predictive theories are supported by experiments, but may diverge outside of this support.

2.1.2 The Fundamental Theorem of Science

With these definitions we can prove, from first principle, that the possibility of falsification is a necessary consequence of the scientific method.

Theorem 3 (The Fundamental Theorem of Science). *Let \mathbf{m}_1 and \mathbf{m}_2 be two manifests, such that the later includes and is larger than the former: $\mathbf{m}_1 \subset \mathbf{m}_2$. If ET_2 is an empirical theory of \mathbf{m}_2 , then it follows that ET_2 is a predictive theory of \mathbf{m}_1 . Finally, up to factual-isomorphism, $\text{Dom}(\text{ET}_2)$ has measure 0 over the set of all distinct domains spawned by the predictive theories of \mathbf{m}_2 .*

Proof. $\text{Dom}(\text{ET}_2)$ is unique. Yet, the number of distinct domains spawned by the set of all possible predictive theories of \mathbf{m}_1 is infinite. Finally, the measure of one element of an infinite set is 0. \square

Consequently, the fundamental theorem of science leads to the concept of falsification, as commonly understood in the philosophy of science and as given in the sense of Popper. It is (almost) certain that a predictive scientific theory will eventually be falsified.

2.2 Axiomatic Foundation of Physics

Recall that earlier we used a dovetailing algorithm in Theorem 2 as an implementation of the scientific method, and we claimed that although it was a possible strategy, it was not necessarily the optimal one. So what then is the optimal one? Do we want to optimize for an efficient algorithm, perhaps the most elegant one, or the one that uses the least amount of memory, etc., but thinking in those terms would be a trap — we must think a bit more abstractly

than postulating or arguing for a specific properties of the algorithm which implements the strategy. We must optimize for axiomatic information.

The best strategy will be to *maximize the axiomatic information* gained from the scientific method as it produces a manifest, and this means, in the technical sense, to maximize the entropy of a probability measure on experimental space.

To have information (in the information-theoretical sense) one must work with a probability measure:

Axiom 1 (Observer). *An observer, denoted as \mathcal{O} , is a measure space over experimental space:*

$$\mathcal{O} := (\mathbf{m}, \mathbb{E}, \rho : \mathbb{E} \rightarrow [0, 1]) \quad (19)$$

where ρ is a probability measure, \mathbf{m} is a manifest, and \mathbb{E} is the experimental space of \mathbf{m} . The definition is reminiscent of a measure space in measure theory, with the important difference that the elements of \mathbb{E} are tuples and not sets. The consequences of this difference will be investigated in the main result section, where a prescription to derive a measure from this triplet is provided.

An observer is a specialization of the definition of the manifest in the sense that a manifest is an element of experimental space, and the observer admits a *measure* over experimental space. Note that typically in physics, the observer (which is not mathematically integrated into the formalism... leading to a family of open problems regarding the 'observer effect') is associated to a random selection of an element from a set of possible observations. This 'effect' will eventually be revealed to be a consequence of the present definition, inherited from the requirement of maximizing axiomatic information. Although it is still obscure at this point, this definition will perfectly coincide with what we understand an observer to be in quantum physics, and also to what we understand it to be in general relativity.

Thesis 2 (Model of Physics). *The probability measure that maximizes the entropy of \mathcal{O} constitutes the laws of physics.*

The laws of physics are a specialization of the definition of the observer, in the sense that an observer admits a measure over experimental space, and the laws of physics are the *entropy-maximized measure* of the same space.

2.2.1 Example

I have been told, despite the ultimate simplicity of the previous axiom and model, this is too abstract (and unconventional) to fix the intuition correctly. Therefore, let me provide a simple example to illustrate the concept.

How do I learn to tie my shoes?

Answer: I practice until I figure out the algorithm. Each practice run is a sequence of steps which produced a result. From these trials I may have learned 50 ways in which a knot is produced, and maybe 15 in which it is not. Each

execution is a series of steps that, if defined and catalogued well enough to be formally reproducible, maps to a halting program. The group of all trials both successful and unsuccessful constitute a manifest.

Now, what most do is try to come up with a physical theory of knot-making hoping to get some physics out of it by constructing an axiomatic basis able to account for the laces, their degrees of freedom, the movement of the hands, the forces involved, etc.

But this method is SCIENTIFICALLY NAIVE (naive being used in the mathematical sense of 'not formal'), because it is perpetually falsified as more experiments are produced, as such will eventually contradict the assumptions that went into the proposed axiomatic basis... and ultimately revealing said axiomatic basis was simply pulled out of nowhere substantial.

The correct method is as done in this paper; resulting in a equivalent representation of the steps of the programs part of the manifest... except said equivalent representation is now derived under maximum entropy conditions, resulting in a maximally informative representation of the construction of the manifest and execution of its programs. This representation of the elements of knowledge and how the steps are performed results in an extended theory quantum computations which embeds gravity effortlessly. The laws of physics describes (and limits) how computation can occur for the construction of a manifest.

The "degrees of freedom, movement of the hands, and forces involved" in knot tying are recovered automatically from the resulting physical theory, as the passage from state LACE UNTIED to state LACE TIED is executed via quantum computations.

Consequently, we do not naively postulate the laws of physics, rather we derive them!

2.2.2 Additional Comments

Let us note the following comments:

1. (One Observer): The system assumes only one observer (at least axiomatically). Other observers, if they exist, ought to be derived by the lone postulated observer from inspecting the manifest (as is everything else).
2. (Randomness): The definition of an observer is introduced as an axiom (and not a definition) because randomness is not derivable by math, and the definition references randomness. Randomness is a consequence implied by the idempotence of the 'given': a halting manifest. Since no more fundamental representation of the manifest are possible within mathematics, a more fundamental construction consists of attributing the 'random' designator to each element of axiomatic information. In this case, as it contains randomness, a concept unknown to mathematics, the resulting model is a physical theory.

3. (Minimalistic Assumptions): By design, the probability measure utilized is the measure which assumes the least possible about the space (i.e. max entropy). It is therefore the 'philosophically and axiomatically safest' measure of this space.
4. (Implied inviolability): Maximizing the entropy in this case coincides with maximizing the randomness of selecting alternative experiments to what has happened — thus represents the most permissive model of experimental space. The perceived inviolability, by the observer, of said model are an implicit consequence of this maximization. One cannot violate that which is maximally permissive.
5. (Information maximum): The laws of physics, defined as such, represent the information-theoretical maximum of interpreting the manifest as a known message, in the Claude Shannon information theoretical sense, from reality, selected from the set of all possible realities.

3 Main Result

Let us now use these definitions to derive the laws of physics from first principle, and show that they are at least the same as the laws of physics, and then investigate extended results and predictions.

Our starting point will be the definition of the observer, then we will maximize the entropy of ρ using the method of the Lagrange multipliers. We recall that our definition of the observer is:

$$\mathcal{O}_i := (\mathbf{m}_i, \mathbb{E}, \rho : \mathbb{E} \rightarrow [0, 1]) \quad (20)$$

where \mathbf{m}_i is a n -tuple, \mathbb{E} is a "powertuple" and ρ is a (probability) measure over \mathbb{E} .

Note the similarity between our definition of the observer to that of a measure space. Comparatively, the definition of a measure space is:

$$M := (X, \Sigma, \mu(X)) \quad (21)$$

where X is a set, Σ is (often) taken to be the powerset of X , and μ is a measure over Σ . The difference is simply that sets have been replaced by tuples. Consequently, we must adapt the definition of a measure space from set to tuples. To do so, we will use the following prescription:

1. We assign a non-negative number to each element of \mathbf{m} .
2. We then equip said numbers with the addition operation, converting the construction to a vector space.

3. We maximize the entropy of a single halting program under the effect of constraints, by using the method of the Lagrange multipliers.
4. We prescribe that any and all constraints on said entropy must remain invariant with respect to a change of basis of said vector space.
5. Finally, we use the tensor product n-times over said vector space to construct probability measure of n-tuples of halting programs.

Explicitly, we maximize the entropy:

$$S = - \sum_{f \in \mathbb{M}} \rho(f) \ln \rho(f) \quad (22)$$

subject to these constraints:

$$\sum_{f \in \mathbb{M}} \rho(f) = 1 \quad (23)$$

$$\sum_{f \in \mathbb{M}} \rho(f) \operatorname{tr} \mathbf{M}(f) = \operatorname{tr} \overline{\mathbf{M}} \quad (24)$$

where $\mathbf{M}(f)$ are a matrix-valued maps from \mathbb{M} to $\mathbb{C}^{n \times n}$ representing the linear transformations of the vector space and where $\overline{\mathbf{M}}$ is a element-by-element average matrix.

Usage of the trace of a matrix as a constraint imposes an invariance with respect to a similarity transformation, accounting for all possible linear reordering of the elements of the tuples of the sum, thus allowing the creation of a measure of a tuple or group of tuples form within a space of tuples, invariantly with respect to the order of the elements of the tuples. Similarity transformation invariance on the trace is the result of this identity:

$$\operatorname{tr} \mathbf{M} = \operatorname{tr} \mathbf{B} \mathbf{M} \mathbf{B}^{-1} \quad (25)$$

We now use the Lagrange multiplier method to derive the expression for ρ that maximizes the entropy, subject to the above mentioned constraints. Maximizing the following equation with respect to ρ yields the answer:

$$\mathcal{L} = -k_B \sum_{f \in \mathbb{M}} \rho(f) \ln(f) + \alpha \left(1 - \sum_{f \in \mathbb{M}} \rho(f) \right) + \tau \left(\operatorname{tr} \overline{\mathbf{M}} - \sum_{f \in \mathbb{M}} \rho(f) \operatorname{tr} \mathbf{M}(f) \right) \quad (26)$$

where α and τ are the Lagrange multipliers. The explicit derivation is made available in Annex B. The result of the maximization process is:

$$\rho(f, \tau) = \frac{1}{Z(\tau)} \det \exp -\tau \mathbf{M}(f) \quad (27)$$

where

$$Z(\tau) = \sum_{f \in \mathbb{M}} \det \exp -\tau \mathbf{M}(f) \quad (28)$$

3.1 Prior

No good probability measure is complete without a prior. The prior, which accounts for an arbitrary preparation of the ensemble, ought to be—for purposes of preserving the scope of the theory—of the same kind as the elements of the probability measure. Let us thus introduce the prior as the map $\mathbf{P} : \mathbb{Q} \rightarrow \mathbb{C}^{n \times n}$ and inject it into the probability measure as well as into the partition function:

$$\rho(f) = \frac{1}{Z} \det \exp (\mathbf{P}(f)) \det \exp (-\tau \mathbf{M}(f)) \quad (29)$$

where

$$Z = \sum_{f \in \mathbb{M}} \det \exp (\mathbf{P}(f)) \det \exp (-\tau \mathbf{M}(f)) \quad (30)$$

3.2 Overview

We will begin with a small overview connecting the main result to its origins as a manifest, via a passage into quantum computation, then we will expand the connection to physics into two parts:

Part 1: We will show that a special case of this result is an equivalent representation of standard quantum mechanics, and discuss the implications of this formulation.

Part 2: We will show that this result, in the general case, is a quantum theory of gravity which adheres to both 1) the axioms of quantum mechanics and 2) to the theory of general relativity.

3.2.1 Quantum Computing (Recall)

Let us begin by reviewing the basics of quantum computation. One starts with a state vector:

$$|\psi_a\rangle = \begin{pmatrix} 0 \\ \vdots \\ n \end{pmatrix} \quad (31)$$

Which evolves unitarily to a final state:

$$|\psi_b\rangle = U_0 U_1 \dots U_m |\psi_a\rangle \quad (32)$$

Clever use of the unitary transformations, often arranged as simple 'gates', allows one to execute a program. The input to the program is the state $|\psi_a\rangle$ and the output is the state $|\psi_b\rangle$. One would note that, so defined and if the sequence of unitary transformation is finite, such a program must always halt, and thus its complexity must be bounded. One can however get out of this predicament by taking the final state $|\psi_b\rangle$ to instead be an intermediary state, and then to add more gates in order continue with a computation:

$$\text{step 1} \quad |\psi_b\rangle = U_0 U_1 \dots U_p |\psi_a\rangle \quad (33)$$

$$\text{step 2} \quad |\psi_c\rangle = U'_0 U'_1 \dots U'_q |\psi_b\rangle \quad (34)$$

$$\vdots \quad (35)$$

$$\text{step k} \quad |\psi_{k'}\rangle = U'_0 U'_1 \dots U'_v |\psi_k\rangle \quad (36)$$

$$\vdots \quad (37)$$

For a quantum computation to simulate a universal Turing machine it must be able to add more steps until a halting state is reached (or continue to add steps indefinitely if the program never halts). But note, that each step is itself a completed program, and further it is the case that each step can be infinitely divided.

3.2.2 Sketch

The linear transformations of our main result are here interpreted in the same manner as those used in quantum computations, but extended to the general linear group.

We begin by splitting the probability measure into a first step, which is linear with respect to a 'probability amplitude', and a second which connects the amplitude to the probability. We thus write the probability measure as:

$$\rho(f, \tau) = \frac{1}{Z} \det \psi(f, \tau) \quad (38)$$

where

$$\psi(f, \tau) = \exp(\mathbf{P}(f)) \exp(-\tau \mathbf{M}(f)) \quad (39)$$

Here, the determinant is interpreted as a generalization of the Born rule and reduces to exactly it when \mathbf{M} is the matrix representation of the complex

numbers. In the general case where \mathbf{M} are arbitrary $n \times n$ matrices, $\psi(f, \tau)$ will be called the *general linear probability amplitude*.

We can write $\psi(f, \tau)$ as a column vector:

$$\boldsymbol{\psi} := |\psi\rangle := \begin{pmatrix} \psi(f_1, \tau) \\ \psi(f_2, \tau) \\ \vdots \\ \psi(f_n, \tau) \end{pmatrix} = \begin{pmatrix} \psi_1 \\ \psi_2 \\ \vdots \\ \psi_n \end{pmatrix} \quad (40)$$

Paths will be constructed by chaining transformations on those vectors:

$$|\psi_b\rangle = \underbrace{\mathbf{T}_1 \mathbf{T}_2 \dots \mathbf{T}_n}_{\text{computing steps}} |\psi_a\rangle \quad (41)$$

As more transformations are chained, progressively richer manifests are constructed. Paths in experimental space are realized by completing the missing computational steps required for a starting-point manifest to be the end-point manifest.

Comparatively, quantum mechanical computations are simply a special cases when the transformations are unitary:

$$|\psi_b\rangle = \underbrace{\mathbf{U}_1 \mathbf{U}_2 \dots \mathbf{U}_n}_{\text{computing steps}} |\psi_a\rangle \quad (42)$$

3.2.3 Matrix-Valued Vector and Transformations

To work with the general linear probability amplitude, we will use vectors whose elements are matrices. An example of such a vector is:

$$\mathbf{v} = \begin{pmatrix} \mathbf{M}_1 \\ \vdots \\ \mathbf{M}_m \end{pmatrix} \quad (43)$$

Likewise a linear transformation of this space will be expressed as a matrix of matrices:

$$\mathbf{T} = \begin{pmatrix} \mathbf{M}_{00} & \dots & \mathbf{M}_{0m} \\ \vdots & \ddots & \vdots \\ \mathbf{M}_{m0} & \dots & \mathbf{M}_{mm} \end{pmatrix} \quad (44)$$

Note: The scalar element of the vector space are given as:

$$a\mathbf{v} = \begin{pmatrix} a\mathbf{M}_1 \\ \vdots \\ a\mathbf{M}_m \end{pmatrix} \quad (45)$$

3.2.4 Sums of Paths

So far, the sums of programs we have used were over manifests comprised of a single program each. How do we extend this to manifests containing multiple programs? We have to use a Cartesian product on the sets of manifests and a tensor product on the probability amplitudes. For instance, let us consider the following sets of halting programs:

$$\mathbb{M}_1 = \{(f_{1a}), (f_{1b})\} \quad (46)$$

$$\mathbb{M}_2 = \{(f_{2a}), (f_{2b})\} \quad (47)$$

The Cartesian product produces manifests comprised of two elements:

$$\mathbf{m} \in \mathbb{M}_1 \times \mathbb{M}_2 = \{(f_{1a}, f_{2a}), (f_{1a}, f_{2b}), (f_{1b}, f_{2a}), (f_{1b}, f_{2b})\} \quad (48)$$

At the level of the probability amplitude, the Cartesian product of sets translates to the tensor product. For instance, we start with a wave-function of one program;

$$\psi_1 = \begin{pmatrix} \exp \mathbf{P}(f_{1a}) \\ \exp \mathbf{P}(f_{1b}) \end{pmatrix} \quad (49)$$

Adding a program-step via a linear transformation produces:

$$\mathbf{T}\psi_1 = \begin{pmatrix} T_{00} \exp \mathbf{P}(f_{1a}) + T_{01} \exp \mathbf{P}(f_{1b}) \\ T_{10} \exp \mathbf{P}(f_{1a}) + T_{11} \exp \mathbf{P}(f_{1b}) \end{pmatrix} \quad (50)$$

If we tensor product this wave-function:

$$\psi_2 = \begin{pmatrix} \exp \mathbf{P}(f_{2a}) \\ \exp \mathbf{P}(f_{2b}) \end{pmatrix} \quad (51)$$

along with a program-step:

$$\mathbf{T}'\psi_2 = \begin{pmatrix} T'_{00} \exp \mathbf{P}(f_{2a}) + T'_{01} \exp \mathbf{P}(f_{2b}) \\ T'_{10} \exp \mathbf{P}(f_{2a}) + T'_{11} \exp \mathbf{P}(f_{2b}) \end{pmatrix} \quad (52)$$

Then the tensor product of these states produces the probability measure of a manifest as follows:

$$\mathbf{T}\psi_1 \otimes \mathbf{T}'\psi_2 = \begin{pmatrix} (T_{00} \exp \mathbf{P}(f_{1a}) + T_{01} \exp \mathbf{P}(f_{1b}))(T'_{00} \exp \mathbf{P}(f_{2a}) + T'_{01} \exp \mathbf{P}(f_{2b})) \\ (T_{00} \exp \mathbf{P}(f_{1a}) + T_{01} \exp \mathbf{P}(f_{1b}))(T'_{10} \exp \mathbf{P}(f_{2a}) + T'_{11} \exp \mathbf{P}(f_{2b})) \\ (T_{10} \exp \mathbf{P}(f_{1a}) + T_{11} \exp \mathbf{P}(f_{1b}))(T'_{00} \exp \mathbf{P}(f_{2a}) + T'_{01} \exp \mathbf{P}(f_{2b})) \\ (T_{10} \exp \mathbf{P}(f_{1a}) + T_{11} \exp \mathbf{P}(f_{1b}))(T'_{10} \exp \mathbf{P}(f_{2a}) + T'_{11} \exp \mathbf{P}(f_{2b})) \end{pmatrix} \quad (53)$$

Now, each element of the resulting vector is a manifest of two programs, but its probability is a sum over a path. One can repeat the process n times, and even take to limit to infinity.

4 Foundation of Physics (Standard QM)

Remarkably, the main result, even if it is barely three lines, is able to produce, as theorems, all axioms of quantum mechanics –including the famously elusive wave-function collapse mechanism, and the origin of the Born rule– thus it constitute, necessarily, a more fundamental formalism of QM that all other alternatives requiring axiomatic definitions. Finally, it necessitate, implies and guarantees not only ensemble interpretation of quantum mechanics, but also derives it from maximum entropy principles. It thus offers a ”total solution” to the foundations, origins and interpretation of QM.

The study of the main result will be in two parts. In the first part, to recover QM, we will reduce the expressivity of the constraint of the main result to that of the matrix representation of the complex number, as opposed to a general matrix. In this special case, we will recover standard QM without gravity. In the section following this one, we will investigate the unadulterated version of the main result and show that it produces a complete theory of quantum gravity adhering both to 1) all axioms of QM, and also 2) is a complete formulation of general relativity, thus producing what we interpret as a theory of quantum gravity.

4.1 Born Rule

Definition 20 (Born Rule). *The standard definition of the Born rule connects the probability to the complex norm of the wave-function:*

$$\rho = |\psi|^2 \quad (54)$$

The Born rule was postulated by Born in 1926, but attempts to derive it from first principles has been elusive. Notable proposals are those of Andrew M. Gleason[7], Kastner, R. E[8], and Lluís Masanes and Thomas Galley[9], and others.

In our formalism the determinant rule is the result of maximizing the entropy of a measure over a tuple. In the case of a matrix representation of the complex

number, the determinant rule becomes the Born rule. Indeed, we first note that a complex number is represented by a matrix as follows:

$$a + ib \cong \begin{pmatrix} a & -b \\ b & a \end{pmatrix} \quad (55)$$

Then, we note that its determinant is the same as the complex norm:

$$\det \begin{pmatrix} a & -b \\ b & a \end{pmatrix} = a^2 + b^2 \quad (56)$$

Finally, the determinant rule reduces to the Born rule for the complex case:

$$\rho(f, \tau) = \frac{1}{Z} \det \left(\exp -\beta \begin{pmatrix} a(f) & 0 \\ 0 & a(f) \end{pmatrix} \exp -\tau \begin{pmatrix} 0 & -b(f) \\ b(f) & 0 \end{pmatrix} \right) \quad (57)$$

$$= \frac{1}{Z} \exp -\beta 2a(f) \quad (58)$$

where $\exp -\beta 2a(f)$ is the preparation of the ensemble (the prior) and Z is the normalization constants. In the case where $\beta = 1/k_B T$, then it is a thermal preparation. The Born rule is thus revealed to be the probability measure derived from maximizing the entropy of the selection of a tuple from a set of tuples, and in this sense is analogous to the Gibbs measure in statistical physics.

4.2 Axioms of QM

4.2.1 Dirac-Von Neumann Axioms (State vector)

The standard Dirac-Von Neumann Axioms are recovered as theorems. This becomes apparent if we split the probability measure into a two-step process:

$$\rho(f, \tau) = \det \psi(f, \tau) |_{\psi \in \mathbb{C}} = |\psi(f, \tau)|^2 \quad (59)$$

where

$$\psi(f, \tau) = \frac{1}{\sqrt{Z}} \left(\exp -\beta \begin{pmatrix} a(f) & 0 \\ 0 & a(f) \end{pmatrix} \exp -\tau \begin{pmatrix} 0 & -b(f) \\ b(f) & 0 \end{pmatrix} \right) \quad (60)$$

The formalism in terms of Hilbert space is obtained simply by taking ψ rather than ρ as the object of study. To show this is almost trivial; $\psi \in \mathbb{C}$ and ρ is given as the complex norm, therefore the ψ are the unit vectors of a complete complex vector space.

4.2.2 Dirac-Von Neumann Axioms (Expectation Value)

In statistical physics, an observable is simply a real value tied to each element of the probability measure:

$$\overline{O} = \sum_{q \in \mathbb{Q}} \rho(q) O(q) \quad (61)$$

Applied to our matrix-based constraints, this definition becomes that of a self-adjoint operator. And the expectation is a sum as follows:

$$\overline{O} = \sum_{q \in \mathbb{Q}} \det \psi(q) O(q) |_{\psi \in \mathbb{C}} = \sum_{q \in \mathbb{Q}} \psi(q)^* \psi(q) O(q) \quad (62)$$

and is thus the same definition as that of the expectation value of a self-adjoint operator acting on the unit vectors of a complex Hilbert space:

$$\sum_{q \in \mathbb{Q}} \psi(q)^* \psi(q) O(q) = \langle \psi | O | \psi \rangle \quad (63)$$

4.2.3 Time-Evolution (Schrödinger equation)

The last axiom of QM is usually a statement that the time-evolution of a state vector ψ is given by the Schrödinger equation. To derive it from our framework, that we can write:

$$\psi(f, \tau) = \frac{1}{\sqrt{Z}} \left(\exp -\beta \begin{pmatrix} a(f) & 0 \\ 0 & a(f) \end{pmatrix} \exp -\tau \begin{pmatrix} 0 & -b(f) \\ b(f) & 0 \end{pmatrix} \right) \quad (64)$$

as:

$$\psi(f, \tau) = \frac{1}{\sqrt{Z}} \exp -\tau i b(f) \quad (65)$$

where, to simplify, we have taking the prior to be unity ($a(f) = 0$). This is the familiar form of the quantum mechanical unitary evolution operator, where we rewrite $\tau \rightarrow t$ and $b \rightarrow H$. The Taylor expansion to the first linear term: $U(\delta t) \approx 1 - i\delta t H$. Then:

$$|\psi(t + \delta t)\rangle = U(\delta t) |\psi(t)\rangle \approx (1 - i\delta t H) |\psi\rangle \quad (66)$$

$$\implies |\psi(t + \delta t)\rangle - |\psi\rangle \approx -i\delta t H |\psi\rangle \quad (67)$$

$$\implies i \frac{|\psi(t + \delta t)\rangle - |\psi\rangle}{\delta t} \approx H |\psi\rangle \quad (68)$$

$$\implies i \frac{\partial |\psi(t)\rangle}{\partial t} = H |\psi\rangle \quad (69)$$

which is the time-dependant Schrödinger equation.

5 Foundation of Physics (Quantum Gravity)

In this section we will show that the probability measure here-in derived adheres to 1) the axioms of quantum mechanics, and 2) is simultaneously a gauge-theoretical theory of gravitation.

We will introduce an *algebra of natural states* and we will use it to classify the linear transformations on said amplitude. We will start with the 2D case, then the 4D case. In all cases, the probability amplitude transforms linearly with respect to general linear transformations and the probability measure, obtained from the determinant, is positive-definite.

5.1 Algebra of Natural States, in 2D

The notation of our upcoming definitions will be significantly improved if we use a geometric representation for matrices. Let us therefore introduce a geometric representation of 2×2 matrices.

5.1.1 Geometric Representation of 2×2 matrices

Let $\mathbb{G}(2, \mathbb{R})$ be the two-dimensional geometric algebra over the reals. We can write a general multi-vector of $\mathbb{G}(2, \mathbb{R})$ as follows:

$$\mathbf{u} = A + \mathbf{X} + \mathbf{B} \quad (70)$$

where A is a scalar, \mathbf{X} is a vector and \mathbf{B} is a pseudo-scalar. Each multi-vector has a structure-preserving (addition/multiplication) matrix representation. Explicitly, the multi-vectors of $\mathbb{G}(2, \mathbb{R})$ are represented as follows:

Definition 21 (Geometric representation of a matrix (2×2)).

$$A + X\hat{\mathbf{x}} + Y\hat{\mathbf{y}} + B\hat{\mathbf{x}} \wedge \hat{\mathbf{y}} \cong \begin{pmatrix} A + X & -B + Y \\ B + Y & A - X \end{pmatrix} \quad (71)$$

And the converse is also true, each 2×2 real matrix is represented as a multi-vector of $\mathbb{G}(2, \mathbb{R})$.

We can define the determinant solely using constructs of geometric algebra[10].

Definition 22 (Clifford conjugate (of a $\mathbb{G}(2, \mathbb{R})$ multi-vector)).

$$\mathbf{u}^\dagger := \langle \mathbf{u} \rangle_0 - \langle \mathbf{u} \rangle_1 - \langle \mathbf{u} \rangle_2 \quad (72)$$

Then the determinant of \mathbf{u} is:

Definition 23 (Geometric representation of the determinant (of a 2×2 matrix)).

$$\begin{aligned} \det &: \mathbb{G}(2, \mathbb{R}) \longrightarrow \mathbb{R} \\ \mathbf{u} &\longmapsto \mathbf{u}^\dagger \mathbf{u} \end{aligned} \quad (73)$$

For example:

$$\det \mathbf{u} = (A - \mathbf{X} - \mathbf{B})(A + \mathbf{X} + \mathbf{B}) \quad (74)$$

$$= A^2 - X^2 - Y^2 + B^2 \quad (75)$$

$$= \det \begin{pmatrix} A + X & -B + Y \\ B + Y & A - X \end{pmatrix} \quad (76)$$

Finally, we define the Clifford transpose:

Definition 24 (Clifford transpose (of a matrix of 2×2 matrix elements)). *The Clifford transpose is the geometric analogue to the conjugate transpose. Like the conjugate transpose can be interpreted as a transpose followed by an element-by-element application of the complex conjugate, here the Clifford transpose is a transpose, followed by an element-by-element application of the Clifford conjugate:*

$$\begin{pmatrix} \mathbf{u}_{00} & \cdots & \mathbf{u}_{0n} \\ \vdots & \ddots & \vdots \\ \mathbf{u}_{m0} & \cdots & \mathbf{u}_{mn} \end{pmatrix}^{\dagger} = \begin{pmatrix} \mathbf{u}_{00}^{\dagger} & \cdots & \mathbf{u}_{m0}^{\dagger} \\ \vdots & \ddots & \vdots \\ \mathbf{u}_{m0} & \cdots & \mathbf{u}_{nm}^{\dagger} \end{pmatrix} \quad (77)$$

If applied to a vector, then:

$$\begin{pmatrix} \mathbf{v}_1 \\ \vdots \\ \mathbf{v}_m \end{pmatrix}^{\dagger} = \begin{pmatrix} \mathbf{v}_1^{\dagger} & \cdots & \mathbf{v}_m^{\dagger} \end{pmatrix} \quad (78)$$

5.1.2 Axiomatic Definition of the Algebra, in 2D

Let \mathbb{V} be an m -dimensional vector space over $\mathbb{G}(2, \mathbb{R})$. A subset of vectors in \mathbb{V} forms an algebra of natural states $\mathcal{A}(\mathbb{V})$ iff the following holds:

1. $\forall \boldsymbol{\psi} \in \mathcal{A}(\mathbb{V})$, the bilinear map:

$$\begin{aligned} \langle \cdot, \cdot \rangle &: \mathbb{V} \times \mathbb{V} \longrightarrow \mathbb{G}(2, \mathbb{R}) \\ \langle \mathbf{u}, \mathbf{v} \rangle &\longmapsto \mathbf{u}^{\dagger} \mathbf{v} \end{aligned} \quad (79)$$

is positive-definite:

$$\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle \in \mathbb{R}_{>0} \quad (80)$$

2. $\forall \boldsymbol{\psi} \in \mathcal{A}(\mathbb{V})$, then for each element $\psi(q) \in \boldsymbol{\psi}$, the function:

$$\rho(\psi(q), \boldsymbol{\psi}) = \frac{1}{\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle} \psi(q)^{\dagger} \psi(q) \quad (81)$$

is positive-definite:

$$\rho(\psi(q), \boldsymbol{\psi}) \in \mathbb{R}_{>0} \quad (82)$$

We note the following comments and definitions:

- From (1) and (2) it follows that $\forall \boldsymbol{\psi} \in \mathcal{A}(\mathbb{V})$, the probabilities sum to unity:

$$\sum_{\psi(q) \in \boldsymbol{\psi}} \rho(\psi(q), \boldsymbol{\psi}) = 1 \quad (83)$$

- $\boldsymbol{\psi}$ is called a *natural* (or physical) state.
- $\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle$ is called the *partition function* of $\boldsymbol{\psi}$.
- $\rho(q, \boldsymbol{\psi})$ is called the *probability measure* (or generalized Born rule) of $\psi(q)$.
- The set of all matrices \mathbf{T} acting on $\boldsymbol{\psi}$, as $\mathbf{T}\boldsymbol{\psi} \rightarrow \boldsymbol{\psi}'$, which leaves the sum of probabilities normalized (invariant):

$$\sum_{\psi(q) \in \boldsymbol{\psi}} \rho(\psi(q), \mathbf{T}\boldsymbol{\psi}) = \sum_{\psi(q) \in \boldsymbol{\psi}} \rho(\psi(q), \boldsymbol{\psi}) = 1 \quad (84)$$

are the *natural* transformations of $\boldsymbol{\psi}$.

- A matrix \mathbf{O} such that $\forall \mathbf{u} \forall \mathbf{v} \in \mathcal{A}(\mathbb{V})$:

$$\langle \mathbf{O}\mathbf{u}, \mathbf{v} \rangle = \langle \mathbf{u}, \mathbf{O}\mathbf{v} \rangle \quad (85)$$

is called an observable.

- The expectation value of an observable \mathbf{O} is:

$$\langle \mathbf{O} \rangle = \frac{1}{\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle} \langle \mathbf{O}\boldsymbol{\psi}, \boldsymbol{\psi} \rangle \quad (86)$$

5.1.3 Reduction to Complex Hilbert Spaces

It is fairly easy to see that if we reduce the expression of our multi-vectors $(A + \mathbf{X} + \mathbf{B})|_{\mathbf{X} \rightarrow 0} = A + \mathbf{B}$ and further restrict $\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle \in \mathbb{R}_{>0}$ to $\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle = 1$, then we recover the unit vectors of the complex Hilbert spaces:

- Reduction to the conjugate transpose:

$$\left(\langle \mathbf{u}, \mathbf{v} \rangle = \mathbf{u}^\dagger \mathbf{v} \right) \Big|_{\mathbf{x} \rightarrow 0} \implies \langle \mathbf{u}, \mathbf{v} \rangle = \mathbf{u}^\dagger \mathbf{v} \quad (87)$$

- Reduction to the unitary transformations:

$$\left(\langle \mathbf{T}\mathbf{u}, \mathbf{T}\mathbf{v} \rangle = \langle \mathbf{u}, \mathbf{v} \rangle \right) \Big|_{\mathbf{x} \rightarrow 0} \implies \mathbf{T}^\dagger \mathbf{T} = I \quad (88)$$

- Reduction to the Born rule:

$$\left(\rho(q, \boldsymbol{\psi}) = \frac{1}{\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle} \psi(q)^\dagger \psi(q) \right) \Big|_{\mathbf{x} \rightarrow 0} \implies \rho(q, \boldsymbol{\psi}) = \frac{1}{\langle \boldsymbol{\psi}, \boldsymbol{\psi} \rangle} \psi(q)^\dagger \psi(q) \quad (89)$$

- Reduction of observables to Hermitian operators:

$$(\langle \mathbf{O}\mathbf{u}, \mathbf{v} \rangle = \langle \mathbf{u}, \mathbf{O}\mathbf{v} \rangle)|_{\mathbf{x} \rightarrow 0} \implies \mathbf{O}^\dagger = \mathbf{O} \quad (90)$$

Under this reduction, the formalism becomes equivalent to the Dirac-Von-Neumann formalism of quantum mechanics.

5.1.4 Observable, in 2D — Self-Adjoint Operator

Let us now investigate the general case of an observable in 2D. A matrix \mathbf{O} is an observable iff it is a self-adjoint operator; defined as:

$$\langle \mathbf{O}\phi, \psi \rangle = \langle \phi, \mathbf{O}\psi \rangle \quad (91)$$

$$\forall \mathbf{u} \forall \mathbf{v} \in \mathbb{V}.$$

Setup: Let $\mathbf{O} = \begin{pmatrix} O_{00} & O_{01} \\ O_{10} & O_{11} \end{pmatrix}$ be an observable. Let ϕ and ψ be 2 two-state vectors $\phi = \begin{pmatrix} \phi_1 \\ \phi_2 \end{pmatrix}$ and $\psi = \begin{pmatrix} \psi_1 \\ \psi_2 \end{pmatrix}$. Here, the components $\phi_1, \phi_2, \psi_1, \psi_2, O_{00}, O_{01}, O_{10}, O_{11}$ are multi-vectors of $\mathbb{G}(2, \mathbb{R})$.

Derivation: 1. Let us now calculate $\langle \mathbf{O}\phi, \psi \rangle$:

$$\begin{aligned} 2\langle \mathbf{O}\phi, \psi \rangle &= (O_{00}\phi_1 + O_{01}\phi_2)^\dagger \psi_1 + \psi_1^\dagger (O_{00}\phi_1 + O_{01}\phi_2) \\ &\quad + (O_{10}\phi_1 + O_{11}\phi_2)^\dagger \psi_2 + \psi_2^\dagger (O_{10}\phi_1 + O_{11}\phi_2) \end{aligned} \quad (92)$$

$$\begin{aligned} &= \phi_1^\dagger O_{00}^\dagger \psi_1 + \phi_2^\dagger O_{01}^\dagger \psi_1 + \psi_1^\dagger O_{00} \phi_1 + \psi_1^\dagger O_{01} \phi_2 \\ &\quad + \phi_1^\dagger O_{10}^\dagger \psi_2 + \phi_2^\dagger O_{11}^\dagger \psi_2 + \psi_2^\dagger O_{10} \phi_1 + \psi_2^\dagger O_{11} \phi_2 \end{aligned} \quad (93)$$

2. Now, $\langle \phi, \mathbf{O}\psi \rangle$:

$$\begin{aligned} 2\langle \phi, \mathbf{O}\psi \rangle &= \phi_1^\dagger (O_{00}\psi_1 + O_{01}\psi_2) + (O_{00}\psi_1 + O_{01}\psi_2)^\dagger \phi_1 \\ &\quad + \phi_2^\dagger (O_{10}\psi_1 + O_{11}\psi_2) + (O_{10}\psi_1 + O_{11}\psi_2)^\dagger \phi_2 \end{aligned} \quad (94)$$

$$\begin{aligned} &= \phi_1^\dagger O_{00} \psi_1 + \phi_1^\dagger O_{01} \psi_2 + \psi_1^\dagger O_{00}^\dagger \phi_1 + \psi_2^\dagger O_{01}^\dagger \phi_1 \\ &\quad + \phi_2^\dagger O_{10} \psi_1 + \phi_2^\dagger O_{11} \psi_2 + \psi_1^\dagger O_{10}^\dagger \phi_1 + \psi_2^\dagger O_{11}^\dagger \phi_2 \end{aligned} \quad (95)$$

For $\langle \mathbf{O}\phi, \psi \rangle = \langle \phi, \mathbf{O}\psi \rangle$ to be realized, it follows that these relations must hold:

$$O_{00}^\dagger = O_{00} \quad (96)$$

$$O_{01}^\dagger = O_{10} \quad (97)$$

$$O_{10}^\dagger = O_{01} \quad (98)$$

$$O_{11}^\dagger = O_{11} \quad (99)$$

Therefore, it follows that it must be the case that \mathbf{O} must be equal to its own Clifford transpose. Thus, \mathbf{O} is an observable iff:

$$\mathbf{O}^\dagger = \mathbf{O} \quad (100)$$

which is the equivalent of the self-adjoint operator $\mathbf{O}^\dagger = \mathbf{O}$ of complex Hilbert spaces.

5.1.5 Observable, in 2D — Eigenvalues / Spectral Theorem

Let us show how the spectral theorem applies to $\mathbf{O}^\dagger = \mathbf{O}$, such that its eigenvalues are real. Consider:

$$\mathbf{O} = \begin{pmatrix} a_{00} & a - xe_1 - ye_2 - be_{12} \\ a + xe_1 + ye_2 + be_{12} & a_{11} \end{pmatrix} \quad (101)$$

In this case, it follows that $\mathbf{O}^\dagger = \mathbf{O}$:

$$\mathbf{O}^\dagger = \begin{pmatrix} a_{00} & a - xe_1 - ye_2 - be_{12} \\ a + xe_1 + ye_2 + be_{12} & a_{11} \end{pmatrix} \quad (102)$$

This example is the most general 2×2 matrix \mathbf{O} such that $\mathbf{O}^\dagger = \mathbf{O}$. The eigenvalues are obtained as follows:

$$0 = \det(\mathbf{O} - \lambda I) = \det \begin{pmatrix} a_{00} - \lambda & a - xe_1 - ye_2 - be_{12} \\ a + xe_1 + ye_2 + be_{12} & a_{11} - \lambda \end{pmatrix} \quad (103)$$

implies:

$$0 = (a_{00} - \lambda)(a_{11} - \lambda) - (a - xe_1 - ye_2 - be_{12})(a + xe_1 + ye_2 + be_{12} + a_{11}) \quad (104)$$

$$0 = (a_{00} - \lambda)(a_{11} - \lambda) - (a^2 - x^2 - y^2 + b^2) \quad (105)$$

finally:

$$\lambda = \left\{ \frac{1}{2} \left(a_{00} + a_{11} - \sqrt{(a_{00} - a_{11})^2 + 4(a^2 - x^2 - y^2 + b^2)} \right) \right\}, \quad (106)$$

$$\frac{1}{2} \left(a_{00} + a_{11} + \sqrt{(a_{00} - a_{11})^2 + 4(a^2 - x^2 - y^2 + b^2)} \right) \} \quad (107)$$

We note that in the case where $a_{00} - a_{11} = 0$, the roots would be complex iff $a^2 - x^2 - y^2 + b^2 < 0$, but we already stated that the determinant of real matrices must be greater than zero because the exponential maps to the orientation-preserving general linear group—therefore it is the case that $a^2 - x^2 - y^2 + b^2 \geq 0$, as this expression is the determinant of the multi-vector. Consequently, $\mathbf{O}^\dagger = \mathbf{O}$ —implies, for orientation-preserving¹ transformations, that its roots are real-valued, and thus constitute a 'geometric' observable in the traditional sense of an observable whose eigenvalues are real-valued.

5.2 Algebra of Natural States, in 3D (brief)

The 3D case will be a stepping stone for the 4D case. A general multi-vector of $\mathbb{G}(3, \mathbb{R})$ can be written as follows:

$$\mathbf{u} = A + \mathbf{X} + \mathbf{V} + \mathbf{B} \quad (108)$$

where A is a scalar, \mathbf{X} is a vector, \mathbf{V} is a pseudo-vector and \mathbf{B} is a pseudo-scalar. Such multi-vectors form a complete representation of 2×2 complex matrices:

$$A + X\sigma_1 + Y\sigma_2 + Z\sigma_3 + V_1i\sigma_1 + V_2i\sigma_2 + V_3i\sigma_3 + B\sigma_1 \wedge \sigma_2 \wedge \sigma_3 \quad (109)$$

$$\cong \begin{pmatrix} A + iB + iV_2 + Z & V_1 + iV_3 + X - iY \\ -V_1 + iV_3 + X + iY & A + iB - iV_2 - Z \end{pmatrix} \quad (110)$$

and the determinant of this matrix connects to the determinant of the multi-vector as follows:

$$\begin{aligned} \det \cdot & : \mathbb{G}(3, \mathbb{R}) \longrightarrow \mathbb{C} \\ \mathbf{u} & \longmapsto \mathbf{u}^\dagger \mathbf{u} \end{aligned} \quad (111)$$

where \mathbf{u}^\dagger is the Clifford conjugate in 3D:

$$\mathbf{u}^\dagger := \langle \mathbf{u} \rangle_0 - \langle \mathbf{u} \rangle_1 - \langle \mathbf{u} \rangle_2 + \langle \mathbf{u} \rangle_3 \quad (112)$$

To produce a real number a further multiplication by its complex conjugate is required:

$$\begin{aligned} |\cdot| & : \mathbb{G}(3, \mathbb{R}) \longrightarrow \mathbb{R} \\ \mathbf{u} & \longmapsto (\mathbf{u}^\dagger \mathbf{u})^\dagger \mathbf{u}^\dagger \mathbf{u} \end{aligned} \quad (113)$$

¹We note the exception that a geometric observable may have real eigenvalues even in the case of a transformation that reverses the orientation if the elements $a_{00} - a_{11}$ are not zero and up to a certain magnitude, whereas transformations in the natural orientation are not bounded by a magnitude — thus creating an orientation-based asymmetry.

where \mathbf{u}^\dagger is defined as:

$$\mathbf{u}^\dagger := \langle \mathbf{u} \rangle_0 + \langle \mathbf{u} \rangle_1 + \langle \mathbf{u} \rangle_2 - \langle \mathbf{u} \rangle_3 \quad (114)$$

5.2.1 Axiomatic Definition of the Algebra, in 3D

Let \mathbb{V} be an m -dimensional vector space over $\mathbb{G}(3, \mathbb{R})$. A subset of vectors in \mathbb{V} forms an algebra of natural states $\mathcal{A}(\mathbb{V})$ iff the following holds:

1. $\forall \psi \in \mathcal{A}(\mathbb{V})$, the quadri-linear map:

$$\begin{aligned} \langle \cdot, \cdot, \cdot, \cdot \rangle &: \mathbb{V} \times \mathbb{V} \times \mathbb{V} \times \mathbb{V} \longrightarrow \mathbb{G}(3, \mathbb{R}) \\ \langle \mathbf{u}, \mathbf{v}, \mathbf{w}, \mathbf{x} \rangle &\longmapsto (\mathbf{u}^\dagger \mathbf{v})^\dagger \mathbf{w}^\dagger \mathbf{x} \end{aligned} \quad (115)$$

is positive-definite:

$$\langle \psi, \psi \rangle \in \mathbb{R}_{>0} \quad (116)$$

2. $\forall \psi \in \mathcal{A}(\mathbb{V})$, then for each element $\psi(q) \in \psi$, the function:

$$\rho(\psi(q), \psi) = \frac{1}{\langle \psi, \psi, \psi, \psi \rangle} (\psi(q)^\dagger \psi(q))^\dagger \psi(q)^\dagger \psi(q) \quad (117)$$

is positive-definite:

$$\rho(\psi(q), \psi) \in \mathbb{R}_{>0} \quad (118)$$

5.2.2 Reduction to Complex Hilbert Spaces

We now consider an algebra of natural states that comprises only those multi-vectors of the form $\mathbf{u}' \propto \mathbf{u}^\dagger \mathbf{u}$ (called a sub-algebra, sub-ring, or 'ideal' of the algebra). We also consider, as we obtain an exponential map from our entropy maximization procedure, only multi-vectors which are exponentiated. Then, the algebra reduces to the foundation of quantum mechanics on complex Hilbert spaces (with an extra internal geometric structure). For example, a wave-function would be of this form:

$$\psi = \begin{pmatrix} (\exp \frac{1}{2} \mathbf{u}_1)^\dagger \exp \frac{1}{2} \mathbf{u}_1 \\ \vdots \\ (\exp \frac{1}{2} \mathbf{u}_m)^\dagger \exp \frac{1}{2} \mathbf{u}_m \end{pmatrix} \quad (119)$$

Each element of ψ are of this form:

$$(\exp \frac{1}{2} \mathbf{u})^\dagger \exp \frac{1}{2} \mathbf{u} = \exp \frac{1}{2} (A - \mathbf{X} - \mathbf{V} + \mathbf{B}) \exp \frac{1}{2} (A + \mathbf{X} + \mathbf{V} + \mathbf{B}) \quad (120)$$

$$= \exp(A + \mathbf{B}) \exp \frac{1}{2} (-\mathbf{X} - \mathbf{V}) \exp \frac{1}{2} (\mathbf{X} + \mathbf{V}) \quad (121)$$

$$= \exp(A + \mathbf{B}) \quad (122)$$

Restricting the algebra to such states reduces the quadri-linear map to a bilinear form:

$$\begin{aligned} \langle \cdot, \cdot \rangle &: \mathcal{A}(\mathbb{V}) \times \mathcal{A}(\mathbb{V}) \longrightarrow \mathbb{C} \\ \langle \psi, \phi \rangle &\longmapsto \psi^\dagger \phi \end{aligned} \quad (123)$$

yielding, when applied to said reduced subset of vectors, the same theory as that of quantum mechanics on complex Hilbert space, but with an extra geometric structure for its observables. The 3D case is a stepping stone for the 4D case, where this extra geometric structure will be revealed to be (in the 4D case) the relativistic wave-function given in the form of a spinor field.

5.3 Algebra of Natural States, in 4D

We will now consider the general case for a vector space over 4×4 matrices.

5.3.1 Geometric Representation (in 4D)

The notation will be significantly improved if we use a geometric representation of matrices. Let $\mathbb{G}(4, \mathbb{R})$ be the two-dimensional geometric algebra over the reals. We can write a general multi-vector of $\mathbb{G}(4, \mathbb{R})$ as follows:

$$\mathbf{u} = A + \mathbf{X} + \mathbf{F} + \mathbf{V} + \mathbf{B} \quad (124)$$

where A is a scalar, \mathbf{X} is a vector, \mathbf{F} is a bivector, \mathbf{V} is a pseudo-vector, and \mathbf{B} is a pseudo-scalar. Each multi-vector has a structure-preserving (addition/multiplication) matrix representation. Explicitly, the multi-vectors of $\mathbb{G}(4, \mathbb{R})$ are represented as follows:

Definition 25 (Geometric representation of a matrix (4×4)).

$$\begin{aligned} &A + T\gamma_0 + X\gamma_1 + Y\gamma_2 + Z\gamma_3 \\ &+ F_{01}\gamma_0 \wedge \gamma_1 + F_{02}\gamma_0 \wedge \gamma_2 + F_{03}\gamma_0 \wedge \gamma_3 + F_{23}\gamma_2 \wedge \gamma_3 + F_{13}\gamma_1 \wedge \gamma_3 + F_{12}\gamma_1 \wedge \gamma_2 \\ &+ V_i\gamma_1 \wedge \gamma_2 \wedge \gamma_3 + V_x\gamma_0 \wedge \gamma_2 \wedge \gamma_3 + V_y\gamma_0 \wedge \gamma_1 \wedge \gamma_3 + V_z\gamma_0 \wedge \gamma_1 \wedge \gamma_2 \\ &+ B\gamma_0 \wedge \gamma_1 \wedge \gamma_2 \wedge \gamma_3 \\ &\cong \begin{pmatrix} A + X_0 - iF_{12} - iV_3 & F_{13} - iF_{23} + V_2 - iV_1 & -iB + X_3 + F_{03} - iV_0 & X_1 - iX_2 + F_{01} - iF_{02} \\ -F_{13} - iF_{23} - V_2 - iV_1 & A + X_0 + iF_{12} + iV_3 & X_1 + iX_2 + F_{01} + iF_{02} & -iB - X_3 - F_{03} - iV_0 \\ -iB - X_3 + F_{03} + iV_0 & -X_1 + iX_2 + F_{01} - iF_{02} & A - X_0 - iF_{12} + iV_3 & F_{13} - iF_{23} - V_2 + iV_1 \\ -X_1 - iX_2 + F_{01} + iF_{02} & -iB + X_3 - F_{03} + iV_0 & -F_{13} - iF_{23} + V_2 + iV_1 & A - X_0 + iF_{12} - iV_3 \end{pmatrix} \end{aligned} \quad (125)$$

And the converse is also true, each 4×4 real matrix is represented as a multi-vector of $\mathbb{G}(4, \mathbb{R})$.

We can define the determinant solely using constructs of geometric algebra[10].

Definition 26 (Clifford conjugate (of a $\mathbb{G}(4, \mathbb{R})$ multi-vector)).

$$\mathbf{u}^\dagger := \langle \mathbf{u} \rangle_0 - \langle \mathbf{u} \rangle_1 - \langle \mathbf{u} \rangle_2 + \langle \mathbf{u} \rangle_3 + \langle \mathbf{u} \rangle_4 \quad (126)$$

and $\lfloor \mathbf{m} \rfloor_{\{3,4\}}$ as the blade-conjugate of degree 3 and 4 (flipping the plus sign to a minus sign for blade 3 and blade 4):

$$\lfloor \mathbf{u} \rfloor_{\{3,4\}} := \langle \mathbf{u} \rangle_0 + \langle \mathbf{u} \rangle_1 + \langle \mathbf{u} \rangle_2 - \langle \mathbf{u} \rangle_3 - \langle \mathbf{u} \rangle_4 \quad (127)$$

The, the determinant of \mathbf{u} is:

Definition 27 (Geometric representation of the determinant (of a 4×4 matrix)).

$$\begin{aligned} \det &: \mathbb{G}(4, \mathbb{R}) \longrightarrow \mathbb{R} \\ \mathbf{u} &\longmapsto \lfloor \mathbf{u}^\dagger \mathbf{u} \rfloor_{3,4} \mathbf{u}^\dagger \mathbf{u} \end{aligned} \quad (128)$$

5.3.2 Axiomatic Definition of the Algebra, in 4D

Let \mathbb{V} be a m -dimensional vector space over the 4×4 real matrices. A subset of vectors in \mathbb{V} forms an algebra of natural states $\mathcal{A}(\mathbb{V})$ iff the following holds:

1. $\forall \psi \in \mathcal{A}(\mathbb{V})$, the quadri-linear form:

$$\begin{aligned} \langle \cdot, \cdot, \cdot, \cdot \rangle &: \mathbb{V} \times \mathbb{V} \times \mathbb{V} \times \mathbb{V} \longrightarrow \mathbb{G}(4, \mathbb{R}) \\ \langle \mathbf{u}, \mathbf{v}, \mathbf{w}, \mathbf{x} \rangle &\longmapsto \lfloor \mathbf{u}^\dagger \mathbf{v} \rfloor_{3,4} \mathbf{w}^\dagger \mathbf{x} \end{aligned} \quad (129)$$

is positive-definite:

$$\langle \psi, \psi, \psi, \psi \rangle \in \mathbb{R}_{>0} \quad (130)$$

2. $\forall \psi \in \mathcal{A}(\mathbb{V})$, then for each element $\psi(q) \in \psi$, the function:

$$\rho(\psi(q), \psi) = \frac{1}{\langle \psi, \psi, \psi, \psi \rangle} \lfloor \psi(q)^\dagger \psi(q) \rfloor_{3,4} \psi(q)^\dagger \psi(q) \quad (131)$$

is positive-definite:

$$\rho(\psi(q), \psi) \in \mathbb{R}_{>0} \quad (132)$$

We note the following properties, features and comments:

- ψ is called a *natural* (or physical) state.
- $\langle \psi, \psi, \psi, \psi \rangle$ is called the *partition function* of ψ .
- $\rho(\psi(q), \psi)$ is called the *probability measure* (or generalized Born rule) of $\psi(q)$.

- The set of all matrices \mathbf{T} acting on ψ such as $\mathbf{T}\psi \rightarrow \psi'$ which leaves the sum of probabilities normalized (invariant):

$$\sum_{\psi(q) \in \psi} \rho(\psi(q), \mathbf{T}\psi) = \sum_{\psi(q) \in \psi} \rho(\psi(q), \psi) = 1 \quad (133)$$

are the *natural* transformations of ψ .

- A matrix \mathbf{O} such that $\forall \mathbf{u} \forall \mathbf{v} \forall \mathbf{w} \forall \mathbf{x} \in \mathbb{V}$:

$$\langle \mathbf{O}\mathbf{u}, \mathbf{v}, \mathbf{w}, \mathbf{x} \rangle = \langle \mathbf{u}, \mathbf{O}\mathbf{v}, \mathbf{w}, \mathbf{x} \rangle = \langle \mathbf{u}, \mathbf{v}, \mathbf{O}\mathbf{w}, \mathbf{x} \rangle = \langle \mathbf{u}, \mathbf{v}, \mathbf{w}, \mathbf{O}\mathbf{x} \rangle \quad (134)$$

is called an observable.

- The expectation value of an observable \mathbf{O} is:

$$\langle \mathbf{O} \rangle = \frac{\langle \mathbf{O}\psi, \psi, \psi, \psi \rangle}{\langle \psi, \psi, \psi, \psi \rangle} \quad (135)$$

5.3.3 Reduction to Complex Hilbert Space

Let us select a subset of multi-vectors. The subset will contain all multi-vectors resulting from the multiplication of an even-multi-vector by its own Clifford conjugate. Consistent with our entropy maximization procedure, the elements will also be exponentiated.

$$\psi = \begin{pmatrix} (\exp \frac{1}{2} \mathbf{u}_1)^\dagger \exp \frac{1}{2} \mathbf{u}_1 \\ \vdots \\ (\exp \frac{1}{2} \mathbf{u}_m)^\dagger \exp \frac{1}{2} \mathbf{u}_m \end{pmatrix} \quad (136)$$

The form of the elements of ψ is:

$$\psi^\dagger \psi = \exp \frac{1}{2} (A - \mathbf{F} + \mathbf{B}) \exp \frac{1}{2} (A + \mathbf{F} + \mathbf{B}) \quad (137)$$

$$= \exp \frac{1}{2} A \exp -\frac{1}{2} \mathbf{F} \exp \frac{1}{2} \mathbf{B} \exp \frac{1}{2} A \exp \frac{1}{2} \mathbf{F} \exp \frac{1}{2} \mathbf{B} \quad (138)$$

$$= \exp A \exp \mathbf{B} \quad (139)$$

On such states, the quadri-linear map is reduced to the Born rule (a bilinear map):

$$\begin{aligned} \langle \cdot, \cdot \rangle &: \mathcal{A}(\mathbb{V}) \times \mathcal{A}(\mathbb{V}) \longrightarrow \mathbb{C} \\ \langle \psi, \phi \rangle &\longmapsto \psi^\dagger \phi \end{aligned} \quad (140)$$

In our example, and with this bilinear map, $\langle \psi^\dagger \psi, \psi^\dagger \psi \rangle = \exp 2A$.

We note the similarity of this sub-algebra to David Hestenes[11]'s geometric algebra formulation of the relativistic wave-function, given as $\psi = \sqrt{\rho} e^{iB/2} e^{\mathbf{F}/2}$. David Hestenes connects his wave-function to a complex number via the reverse $\tilde{\psi} := \sqrt{\rho} e^{iB/2} e^{-\mathbf{F}/2}$, such that $\psi \tilde{\psi} = \rho e^{iB}$.

5.4 Law of Motion

5.4.1 Probability-Preserving Evolution (Left Action in 2D)

A left action on a wave-function : $\mathbf{G}|\psi\rangle$, connects to the bilinear form as follows: $\langle\psi|\mathbf{G}^\dagger\mathbf{G}|\psi\rangle$. The invariance requirement on \mathbf{G} is as follows:

$$\langle\psi|\mathbf{G}^\dagger\mathbf{G}|\psi\rangle = \langle\psi|\psi\rangle \quad (141)$$

We are thus interested in the group of matrices such that:

$$\mathbf{G}^\dagger\mathbf{G} = I \quad (142)$$

Let us consider a two-state system. A general transformation is:

$$\mathbf{G} = \begin{pmatrix} u & v \\ w & x \end{pmatrix} \quad (143)$$

where u, v, w, x are multi-vectors of 2 dimensions. The expression $\mathbf{G}^\dagger\mathbf{G}$ is:

$$\mathbf{G}^\dagger\mathbf{G} = \begin{pmatrix} v^\dagger & u^\dagger \\ w^\dagger & x^\dagger \end{pmatrix} \begin{pmatrix} u & v \\ w & x \end{pmatrix} = \begin{pmatrix} v^\dagger v + u^\dagger u & v^\dagger w + u^\dagger x \\ w^\dagger v + x^\dagger u & w^\dagger w + x^\dagger x \end{pmatrix} \quad (144)$$

For the results to be the identity, it must be the case that:

$$v^\dagger v + u^\dagger u = 1 \quad (145)$$

$$v^\dagger w + u^\dagger x = 0 \quad (146)$$

$$w^\dagger v + x^\dagger u = 0 \quad (147)$$

$$w^\dagger w + x^\dagger x = 1 \quad (148)$$

This is the case if

$$\mathbf{G} = \frac{1}{\sqrt{v^\dagger v + u^\dagger u}} \begin{pmatrix} v & u \\ -e^\varphi u^\dagger & e^\varphi v^\dagger \end{pmatrix} \quad (149)$$

where u, v are multi-vectors of 2 dimensions, and where e^φ is a unit multi-vector. Comparatively, the unitary case is obtained with $\mathbf{X} \rightarrow 0$, and is:

$$\mathbf{U} = \frac{1}{\sqrt{|a|^2 + |b|^2}} \begin{pmatrix} a & b \\ -e^{i\theta} b^\dagger & e^{i\theta} a^\dagger \end{pmatrix} \quad (150)$$

We can show that $\mathbf{G}^\dagger \mathbf{G} = I$ as follows:

$$\Rightarrow \mathbf{G}^\dagger \mathbf{G} = \frac{1}{v^\dagger v + u^\dagger u} \begin{pmatrix} v^\dagger & -e^{-\varphi} u \\ u^\dagger & e^{-\varphi} v \end{pmatrix} \begin{pmatrix} v & u \\ -e^\varphi u^\dagger & e^\varphi v^\dagger \end{pmatrix} \quad (151)$$

$$= \frac{1}{v^\dagger v + u^\dagger u} \begin{pmatrix} v^\dagger v + u^\dagger u & v^\dagger u - v^\dagger u \\ u^\dagger v - u^\dagger v & u^\dagger u + v^\dagger v \end{pmatrix} \quad (152)$$

$$= I \quad (153)$$

In the case where \mathbf{G} and $|\psi\rangle$ are n -dimensional, we can find an expression for it starting from a diagonal matrix:

$$\mathbf{D} = \begin{pmatrix} e^{x_1 \hat{\mathbf{x}} + y_1 \hat{\mathbf{y}} + ib_1} & 0 \\ 0 & e^{x_2 \hat{\mathbf{x}} + y_2 \hat{\mathbf{y}} + ib_2} \end{pmatrix} \quad (154)$$

where $\mathbf{G} = \mathbf{P} \mathbf{D} \mathbf{P}^{-1}$. It follows quite easily that $\mathbf{D}^\dagger \mathbf{D} = I$, because each diagonal entry produces unity: $e^{-x_1 \hat{\mathbf{x}} - y_1 \hat{\mathbf{y}} - ib_1} e^{x_1 \hat{\mathbf{x}} + y_1 \hat{\mathbf{y}} + ib_1} = 1$.

5.4.2 Probability-Preserving Evolution (Adjoint Action in 2D)

Since the elements of $|\psi\rangle$ are matrices, in the general case, the transformation is given by adjoint action:

$$\mathbf{G} |\psi\rangle \mathbf{G}^{-1} \quad (155)$$

The bilinear form is:

$$(\mathbf{G} |\psi\rangle \mathbf{G}^{-1})^\dagger (\mathbf{G} |\psi\rangle \mathbf{G}^{-1}) = (\mathbf{G}^{-1})^\dagger \langle \psi | \mathbf{G}^\dagger \mathbf{G} |\psi\rangle \mathbf{G}^{-1} \quad (156)$$

and the invariance requirement on \mathbf{G} is as follows:

$$(\mathbf{G}^{-1})^\dagger \langle \psi | \mathbf{G}^\dagger \mathbf{G} |\psi\rangle \mathbf{G}^{-1} = \langle \psi | \psi \rangle \quad (157)$$

With a diagonal matrix, this occurs for general linear transformations:

$$\mathbf{D} = \begin{pmatrix} e^{a_1 + x_1 \hat{\mathbf{x}} + y_1 \hat{\mathbf{y}} + ib_1} & 0 & 0 \\ 0 & e^{a_2 + x_2 \hat{\mathbf{x}} + y_2 \hat{\mathbf{y}} + ib_2} & 0 \\ 0 & 0 & \ddots \end{pmatrix} \quad (158)$$

where $\mathbf{G} = \mathbf{P} \mathbf{D} \mathbf{P}^{-1}$.

Taking a single diagonal entry as an example, the reduction is:

$$e^{-a_1+x_1\hat{\mathbf{x}}+y_1\hat{\mathbf{y}}+ib_1}\psi_1^\dagger e^{a_1-x_1\hat{\mathbf{x}}-y_1\hat{\mathbf{y}}-ib_1}e^{a_1+x_1\hat{\mathbf{x}}+y_1\hat{\mathbf{y}}+ib_1}\psi_1 e^{-a_1-x_1\hat{\mathbf{x}}-y_1\hat{\mathbf{y}}-ib_1} \quad (159)$$

$$= e^{-a_1+x_1\hat{\mathbf{x}}+y_1\hat{\mathbf{y}}+ib_1}\psi_1^\dagger e^{2a_1}\psi_1 e^{-a_1-x_1\hat{\mathbf{x}}-y_1\hat{\mathbf{y}}-ib_1} \quad (160)$$

We note that $\psi^\dagger\psi$ is a scalar, therefore

$$= \psi_1^\dagger\psi_1 e^{2a_1} e^{-a_1+x_1\hat{\mathbf{x}}+y_1\hat{\mathbf{y}}+ib_1} e^{-a_1-x_1\hat{\mathbf{x}}-y_1\hat{\mathbf{y}}-ib_1} \quad (161)$$

$$= \psi_1^\dagger\psi_1 e^{2a_1} e^{-a_1} e^{-a_1} = \psi_1^\dagger\psi_1 \quad (162)$$

5.4.3 General Linear Schrödinger Equation (Left Action)

The standard Schrödinger equation can be derived as follows. First, assume $U(t) = e^{-itH}$, and its Taylor expansion to the first linear term: $U(\delta t) \approx 1 - i\delta tH$. Then:

$$|\psi(t + \delta t)\rangle = U(\delta t) |\psi(t)\rangle \approx (1 - i\delta tH) |\psi\rangle \quad (163)$$

$$\implies |\psi(t + \delta t)\rangle - |\psi\rangle \approx -i\delta tH |\psi\rangle \quad (164)$$

$$\implies i \frac{|\psi(t + \delta t)\rangle - |\psi\rangle}{\delta t} \approx H |\psi\rangle \quad (165)$$

$$\implies i \frac{\partial |\psi(t)\rangle}{\partial t} = H |\psi\rangle \quad (166)$$

Now, we wish to use the same derivation, but apply it to the 2D general linear version of the unitary group:

$$U^\dagger U = I \rightarrow (\mathbf{G}^{-1})^\dagger \langle\psi| \mathbf{G}^\dagger \mathbf{G} |\psi\rangle \mathbf{G}^{-1} = \langle\psi|\psi\rangle \quad (167)$$

In the general linear case, the imaginary number i is replaced with an arbitrary matrix \mathbf{M} , via the relation:

$$\mathbf{G} = e^{-\mathbf{M}\tau H} \quad (168)$$

where H is self-adjoint : $H^\dagger = H$.

Then, the general linear Schrödinger equation for the one-parameter group of the general linear group $\mathbf{G}(\tau) = e^{-\mathbf{M}\tau H}$, for the left action is:

$$|\psi(\tau + \delta\tau)\rangle = \mathbf{G}(\delta\tau) |\psi(\tau)\rangle \quad (169)$$

$$\approx (1 - \mathbf{M}\delta\tau H) |\psi(\tau)\rangle \quad (170)$$

$$= |\psi(\tau)\rangle - \mathbf{M}\delta\tau H |\psi(\tau)\rangle \quad (171)$$

$$\implies -\frac{\partial |\psi(\tau)\rangle}{\partial \tau} = \mathbf{M}H |\psi(\tau)\rangle \quad (172)$$

and iff $\exists \mathbf{M}^{-1}$, then

$$-\mathbf{M}^{-1} \frac{\partial |\psi(\tau)\rangle}{\partial \tau} = H |\psi(\tau)\rangle \quad (173)$$

5.4.4 General Linear Schrödinger Equation (Adjoint Action)

And for the adjoint action, it is:

$$|\psi(\tau + \delta\tau)\rangle = G(\delta\tau) |\psi(\tau)\rangle G(\delta\tau)^{-1} \quad (174)$$

$$\approx (1 - \mathbf{M}\delta\tau H) |\psi(\tau)\rangle (1 + \mathbf{M}\delta\tau H) \quad (175)$$

$$= |\psi(\tau)\rangle (1 + \mathbf{M}\delta\tau H) - \mathbf{M}\delta\tau H |\psi(\tau)\rangle (1 + \mathbf{M}\delta\tau H) \quad (176)$$

$$= |\psi(\tau)\rangle + |\psi(\tau)\rangle \mathbf{M}\delta\tau H - \mathbf{M}\delta\tau H |\psi(\tau)\rangle - \mathbf{M}\delta\tau H |\psi(\tau)\rangle \mathbf{M}\delta\tau H \quad (177)$$

$$\approx |\psi(\tau)\rangle + |\psi(\tau)\rangle \mathbf{M}\delta\tau H - \mathbf{M}\delta\tau H |\psi(\tau)\rangle \quad (178)$$

$$= |\psi(\tau)\rangle + \delta\tau [\mathbf{M}H, |\psi(\tau)\rangle] \quad (179)$$

$$\Rightarrow \frac{\partial |\psi(\tau)\rangle}{\partial \tau} = [\mathbf{M}H, |\psi(\tau)\rangle] \quad (180)$$

5.4.5 Conservation of Probability (Left Action in 2D)

For a parametrization of ψ , the probability must normalize. For instance, a x parametrization would yield:

$$\int \psi(\tau, x)^\dagger \psi(\tau, x) dx = N(\tau) \quad (181)$$

To lighten the notation we will not explicitly write the dependance of ψ in (τ, x) .

$$\frac{dN(\tau)}{d\tau} = 0 = \int \frac{\partial \psi^\dagger \psi}{\partial \tau} dx \quad (182)$$

$$= \int \frac{\partial \psi^\dagger}{\partial \tau} \psi dx + \int \psi^\dagger \frac{\partial \psi}{\partial \tau} dx \quad (183)$$

We now inject the following relation (derived from the general linear Schrödinger equation):

$$\frac{\partial \psi}{\partial \tau} = \mathbf{M}H\psi \quad (184)$$

$$\frac{\partial \psi^\dagger}{\partial \tau} = (\mathbf{M}H\psi)^\dagger = \psi^\dagger H^\dagger \mathbf{M}^\dagger \quad (185)$$

Injecting them, we get:

$$\frac{dN(\tau)}{d\tau} = \int \frac{\partial \psi^\dagger}{\partial \tau} \psi \, dx + \int \psi^\dagger \frac{\partial \psi}{\partial \tau} \, dx \quad (186)$$

$$= \int \psi^\dagger \mathbf{M}^\dagger H^\dagger \psi \, dx + \int \psi^\dagger \mathbf{M} H \psi \, dx \quad (187)$$

We note that $\mathbf{M}^\dagger = -\mathbf{M}$ (this requires that $\langle \mathbf{M} \rangle_0 = 0$, which is the case for the left action) and that $H^\dagger = H$ (which it is iff it is self-adjoint), therefore:

$$= - \int \psi^\dagger \mathbf{M} H \psi \, dx + \int \psi^\dagger \mathbf{M} H \psi \, dx \quad (188)$$

Finally, adding the requirement that $[\mathbf{M}, H] = 0$, we get the conservation of probability:

$$= 0 \quad (189)$$

The general linear form of the Schrödinger equation is a conservation of probability law of the general linear case.

5.4.6 Conservation of Probability (Adjoint Action in 2D)

For a parametrization of ψ , the probability must normalize. For instance, a x parametrization would yield:

$$\int \psi(\tau, x)^\dagger \psi(\tau, x) \, dx = N(\tau) \quad (190)$$

To lighten the notation we will not explicitly write the dependance of ψ in (τ, x) .

$$\frac{dN(\tau)}{d\tau} = 0 = \int \frac{\partial \psi^\dagger \psi}{\partial \tau} \, dx \quad (191)$$

$$= \int \frac{\partial \psi^\dagger}{\partial \tau} \psi \, dx + \int \psi^\dagger \frac{\partial \psi}{\partial \tau} \, dx \quad (192)$$

We now inject the following relation (derived from the general linear Schrödinger equation):

$$\frac{\partial \psi}{\partial \tau} = \mathbf{M} H \psi - \psi \mathbf{M} H \quad (193)$$

$$\frac{\partial \psi^\dagger}{\partial \tau} = (\mathbf{M} H \psi - \psi \mathbf{M} H)^\dagger = \psi^\dagger H^\dagger \mathbf{M}^\dagger - H^\dagger \mathbf{M}^\dagger \psi^\dagger \quad (194)$$

Injecting them, we get:

$$\frac{dN(\tau)}{d\tau} = \int \frac{\partial \psi^\dagger}{\partial \tau} \psi \, dx + \int \psi^\dagger \frac{\partial \psi}{\partial \tau} \, dx \quad (195)$$

$$= \int \psi^\dagger H^\dagger \mathbf{M}^\dagger \psi - H^\dagger \mathbf{M}^\dagger \psi^\dagger \psi \, dx + \int \psi^\dagger \mathbf{M} H \psi - \psi^\dagger \psi \mathbf{M} H \, dx \quad (196)$$

$$= \int \psi^\dagger H^\dagger \mathbf{M}^\dagger \psi - H^\dagger \mathbf{M}^\dagger \psi^\dagger \psi + \psi^\dagger \mathbf{M} H \psi - \psi^\dagger \psi \mathbf{M} H \, dx \quad (197)$$

We note that $\mathbf{M}^\dagger = A - \mathbf{m}$ and that $H^\dagger = H$, therefore (we also pose $\mathbf{M} = A + \mathbf{m}$):

$$= \int \psi^\dagger H(A - \mathbf{m})\psi - H(A - \mathbf{m})\psi^\dagger \psi + \psi^\dagger(A + \mathbf{m})H\psi - \psi^\dagger \psi(A + \mathbf{m})H \, dx \quad (198)$$

$$= \int -\psi^\dagger H\mathbf{m}\psi + H\mathbf{m}\psi^\dagger \psi + \psi^\dagger \mathbf{m}H\psi - \psi^\dagger \psi \mathbf{m}H \, dx \quad (199)$$

Finally, adding the requirement that $[\mathbf{m}, H] = 0$, we get the conservation of probability:

$$= 0 \quad (200)$$

The general linear form of the Schrödinger equation is a conservation of probability law of the general linear case.

The 4D case is omitted due to excessive verbosity.

5.5 Gravity as a Gauge Theory

5.5.1 Unitary Gauge (Recap)

The typical gauge theory in quantum electrodynamics is obtained by the production of a gauge covariant derivative over a $U(1)$ invariance associated with the use of the complex norm in any probability measure of quantum mechanics. Localizing the invariance group $\theta \rightarrow \theta(x)$ yields the corresponding covariant derivative:

$$D_\mu = \partial_\mu + iqA_\mu(x) \quad (201)$$

Where $A_\mu(x)$ is the gauge field. The $U(1)$ invariance results from the usage of the complex norm to construct a probability measure in a quantum theory, and the presence of the derivative is the result of constructing said probability measure as the Lagrangian of a Dirac field. If one then applies a gauge transformation to ψ and A_μ :

$$\psi \rightarrow e^{-iq\theta(x)}\psi \quad \text{and} \quad A_\mu \rightarrow A_\mu + \partial_\mu\theta(x) \quad (202)$$

Then, applies the covariant derivation, one gets:

$$D_\mu \psi = \partial_\mu \psi + iq A_\mu \psi \quad (203)$$

$$\rightarrow \partial_\mu (e^{-iq\theta(x)} \psi) + iq (A_\mu + \partial_\mu \theta(x)) (e^{-iq\theta(x)} \psi) \quad (204)$$

$$= e^{-iq\theta(x)} D_\mu \psi \quad (205)$$

Finally, the field is given as follows:

$$F_{\mu\nu} = [\mathcal{D}_\mu, \mathcal{D}_\nu] \quad (206)$$

where \mathcal{D}_μ is the covariant derivative with respect to the potential one-form $A_\mu = A_\mu^\alpha T_\alpha$, and where T_α are the generators of the lie algebra of $U(1)$.

5.5.2 General Linear Gauge

The fundamental invariance group of our measure is the orientation-preserving general linear group $GL^+(n, \mathbb{R})$, if the algebra is even, or the complex general linear group $GL(n, \mathbb{C})$ if the algebra is odd, rather than $U(1)$. Gauging the $GL^+(n, \mathbb{R})$ group is known to substantially connect to general relativity, as the resulting $GL(4, \mathbb{R})$ -valued field can be viewed as the Christoffel symbols Γ^μ , and the commutator of the covariant derivatives as the Riemann tensor..

A general linear transformation of ψ :

$$\psi'(x) \rightarrow g\psi(x)g^{-1} \quad (207)$$

leaves the probability measure invariant.

The gauge-covariant derivative is:

$$D_\mu \psi = \partial_\mu \psi - [iq A_\mu, \psi] \quad (208)$$

Finally, the field is given as follows:

$$R_{\mu\nu} = [D_\mu, D_\nu] \quad (209)$$

where $R_{\mu\nu}$ is the Riemann tensor.

Since this is the result of the GL gauge invariance, then gravity is fundamentally integrated throughout the present quantum mechanical framework because of GL invariance, for the same reason that electromagnetism is fundamentally integrated within quantum theory over the complex norm because of $U(1)$ invariance.

6 Testable Prediction

Certain linear transformations of the wave-function, under the general linear group and its subgroups, would produce richer interference patterns than what is possible merely with complex interference. The possibility of richer interference patterns has been proposed before; specifically, I note the work of B. I. Lev.[12] which suggests (theoretically) the possibility of an extended interference pattern associated with the David Hestenes form of the relativistic wave-function and for the subset of rotors.

We note that interference experiments have paid off substantial dividends in the history of physics and are somewhat easy to construct and more affordable than many alternative experiments.

6.1 Geometric Interference

Let us start by introducing a notation for a dot product, then we will list the various possible interference patterns.

6.1.1 Geometric Algebra Dot Product

Let us introduce a notation. We will define a bilinear form using the dot product notation, as follows:

$$\begin{aligned} \cdot & : \mathbb{G}(2n, \mathbb{R}) \times \mathbb{G}(2n, \mathbb{R}) \longrightarrow \mathbb{R} \\ \mathbf{u} \cdot \mathbf{v} & \longmapsto \frac{1}{2}(\det(\mathbf{u} + \mathbf{v}) - \det \mathbf{u} - \det \mathbf{v}) \end{aligned} \quad (210)$$

For example,

$$\mathbf{u} = A_1 + X_1 e_1 + Y_1 e_2 + B_1 e_{12} \quad (211)$$

$$\mathbf{v} = A_2 + X_2 e_1 + Y_2 e_2 + B_2 e_{12} \quad (212)$$

$$\implies \mathbf{u} \cdot \mathbf{v} = A_1 A_2 + B_1 B_2 - X_1 X_2 - Y_1 Y_2 \quad (213)$$

Iff $\det \mathbf{u} > 0$ and $\det \mathbf{v} > 0$ then $\mathbf{u} \cdot \mathbf{v}$ is always positive, and therefore qualifies as a positive inner product (over the positive det group), but no greater than either $\det \mathbf{u}$ or $\det \mathbf{v}$, whichever is larger. This definition of the dot product extends to multi-vectors of 4 dimensions.

2D: In 2D the dot product is equivalent to this form:

$$\frac{1}{2}(\det(\mathbf{u} + \mathbf{v}) - \det \mathbf{u} - \det \mathbf{v}) = \frac{1}{2} \left((\mathbf{u} + \mathbf{v})^\dagger (\mathbf{u} + \mathbf{v}) - \mathbf{u}^\dagger \mathbf{u} - \mathbf{v}^\dagger \mathbf{v} \right) \quad (214)$$

$$= \mathbf{u}^\dagger \mathbf{u} + \mathbf{u}^\dagger \mathbf{v} + \mathbf{v}^\dagger \mathbf{u} + \mathbf{v}^\dagger \mathbf{v} - \mathbf{u}^\dagger \mathbf{u} - \mathbf{v}^\dagger \mathbf{v} \quad (215)$$

$$= \mathbf{u}^\dagger \mathbf{v} + \mathbf{v}^\dagger \mathbf{u} \quad (216)$$

4D: In 4D it is substantially more verbose:

$$\frac{1}{2}(\det(\mathbf{u} + \mathbf{v}) - \det \mathbf{u} - \det \mathbf{v}) \quad (217)$$

$$= \frac{1}{2} \left([(\mathbf{u} + \mathbf{v})^\dagger(\mathbf{u} + \mathbf{v})]_{3,4}(\mathbf{u} + \mathbf{v})^\dagger(\mathbf{u} + \mathbf{v}) - [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{u} - [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{v} \right) \quad (218)$$

$$= \frac{1}{2} \left([\mathbf{u}^\dagger \mathbf{u} + \mathbf{u}^\dagger \mathbf{v} + \mathbf{v}^\dagger \mathbf{u} + \mathbf{v}^\dagger \mathbf{v}]_{3,4}(\mathbf{u}^\dagger \mathbf{u} + \mathbf{u}^\dagger \mathbf{v} + \mathbf{v}^\dagger \mathbf{u} + \mathbf{v}^\dagger \mathbf{v}) - \dots \right) \quad (219)$$

$$\begin{aligned} &= [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{u} + [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{u} + [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{v} \\ &\quad + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{u} + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{u} + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{v} \\ &\quad + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{u} + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{u} + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{v} \\ &\quad + [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{u} + [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{u} + [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{v} - \dots \end{aligned} \quad (220)$$

$$\begin{aligned} &= [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{u} + [\mathbf{u}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{v} \\ &\quad + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{u} + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{u} + [\mathbf{u}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{v} \\ &\quad + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{u} + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{u} + [\mathbf{v}^\dagger \mathbf{u}]_{3,4} \mathbf{v}^\dagger \mathbf{v} \\ &\quad + [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{u} + [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{u}^\dagger \mathbf{v} + [\mathbf{v}^\dagger \mathbf{v}]_{3,4} \mathbf{v}^\dagger \mathbf{u} \end{aligned} \quad (221)$$

6.1.2 Geometric Interference (General Form)

A multi-vector can be written as $\mathbf{u} = a + \mathbf{s}$, where a is a scalar and \mathbf{s} is the multi-vectorial part. In general, the exponential $\exp \mathbf{u}$ equals $\exp a \exp \mathbf{s}$ because a commutes with \mathbf{s} .

One can thus write a general two-state system as follows:

$$\psi = \psi_1 + \psi_2 = e^{\frac{2}{n}A_1} e^{\frac{2}{n}\mathbf{S}_1} + e^{\frac{2}{n}A_2} e^{\frac{2}{n}\mathbf{S}_2} \quad (222)$$

$$(223)$$

The general interference pattern will be of the following form:

$$\det \psi_1 + \psi_2 = \det \psi_1 + \det \psi_2 + \psi_1 \cdot \psi_2 \quad (224)$$

$$= e^{2A_1} + e^{2A_2} + \psi_1 \cdot \psi_2 \quad (225)$$

where $\det \psi_1 + \det \psi_2$ is a sum of probabilities and where $\psi_1 \cdot \psi_2$ is the interference pattern.

6.1.3 Complex Interference (Recall)

Consider a two-state wave-function:

$$\psi = \psi_1 + \psi_2 = e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{\mathbf{B}_2} \quad (226)$$

The interference pattern familiar to quantum mechanics is the result of the complex norm:

$$\psi^\dagger \psi = \psi_1^\dagger \psi_1 + \psi_2^\dagger \psi_2 + \psi_1^\dagger \psi_2 + \psi_2^\dagger \psi_1 \quad (227)$$

$$= e^{A_1} e^{-\mathbf{B}_1} e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{-\mathbf{B}_2} e^{A_2} e^{\mathbf{B}_2} + e^{A_1} e^{-\mathbf{B}_1} e^{A_2} e^{\mathbf{B}_2} + e^{A_2} e^{-\mathbf{B}_2} e^{A_1} e^{\mathbf{B}_1} \quad (228)$$

$$= e^{2A_1} + e^{2A_2} + e^{A_1+A_2} (e^{-\mathbf{B}_1+\mathbf{B}_2} + e^{-(\mathbf{B}_1+\mathbf{B}_2)}) \quad (229)$$

$$= \underbrace{e^{2A_1} + e^{2A_2}}_{\text{sum}} + \underbrace{2e^{A_1+A_2} \cos(B_1 - B_2)}_{\text{interference}} \quad (230)$$

6.1.4 Geometric Interference in 2D

Consider a two-state wave-function:

$$\psi = \psi_1 + \psi_2 = e^{A_1} e^{\mathbf{X}_1+\mathbf{B}_1} + e^{A_2} e^{\mathbf{X}_2+\mathbf{B}_2} \quad (231)$$

To lighten the notation we will write it as follows:

$$\psi = \psi_1 + \psi_2 = e^{A_1} e^{\mathbf{S}_1} + e^{A_2} e^{\mathbf{S}_2} \quad (232)$$

where

$$\mathbf{S} = \mathbf{X} + \mathbf{B} \quad (233)$$

The interference pattern for a full general linear transformation on a two-state wave-function in 2D is:

$$\psi^\dagger \psi = \psi_1^\dagger \psi_1 + \psi_2^\dagger \psi_2 + \psi_1^\dagger \psi_2 + \psi_2^\dagger \psi_1 \quad (234)$$

$$= e^{A_1} (e^{\mathbf{S}_1})^\dagger e^{A_1} e^{\mathbf{S}_1} + e^{A_2} (e^{\mathbf{S}_2})^\dagger e^{A_2} e^{\mathbf{S}_2} + e^{A_1} (e^{\mathbf{S}_1})^\dagger e^{A_2} e^{\mathbf{S}_2} + e^{A_2} (e^{\mathbf{S}_2})^\dagger e^{A_1} e^{\mathbf{S}_1} \quad (235)$$

$$= e^{2A_1} + e^{2A_2} + e^{A_1+A_2} ((e^{\mathbf{S}_1})^\dagger e^{\mathbf{S}_2} + (e^{\mathbf{S}_2})^\dagger e^{\mathbf{S}_1}) \quad (236)$$

$$= \underbrace{e^{2A_1} + e^{2A_2}}_{\text{sum}} + \underbrace{e^{A_1+A_2} (e^{-\mathbf{X}_1-\mathbf{B}_1} e^{\mathbf{X}_2+\mathbf{B}_2} + e^{-\mathbf{X}_2-\mathbf{B}_2} e^{\mathbf{X}_1+\mathbf{B}_1})}_{\text{interference}} \quad (237)$$

6.1.5 Geometric Interference in 4D

Consider a two-state wave-function:

$$\psi = \psi_1 + \psi_2 = e^{\frac{1}{2}A_1} e^{\frac{1}{2}(\mathbf{X}_1+\mathbf{F}_1+\mathbf{V}_1+\mathbf{B}_1)} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}(\mathbf{X}_2+\mathbf{F}_2+\mathbf{V}_2+\mathbf{B}_2)} \quad (238)$$

To lighten the notation we will write it as follows:

$$\psi = \psi_1 + \psi_2 = e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{S}_1} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{S}_2} \quad (239)$$

where

$$\mathbf{S} = \mathbf{X} + \mathbf{F} + \mathbf{V} + \mathbf{B} \quad (240)$$

The geometric interference patterns for a full general linear transformation in 4D is given by the product:

$$[\psi^\dagger \psi]_{3,4} \psi^\dagger \psi = [\psi_1^\dagger \psi_1]_{3,4} \psi_1^\dagger \psi_1 + [\psi_2^\dagger \psi_2]_{3,4} \psi_2^\dagger \psi_2 + \psi_1 \cdot \psi_2 \quad (241)$$

$$= e^{2A_1} + e^{2A_2} + \left(e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{S}_1} \right) \cdot \left(e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{S}_2} \right) \quad (242)$$

In many cases of interest, the pattern simplifies:

6.1.6 Geometric Interference in 4D (Shallow Phase Rotation)

If we consider a sub-algebra in 4D comprised of even-multi-vector products $\psi^\dagger \psi$, then a two-state system is given as:

$$\psi = \psi_1 + \psi_2 \quad (243)$$

where

$$\psi_1 = (e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{F}_1} e^{\frac{1}{2}\mathbf{B}_1})^\dagger (e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{F}_1} e^{\frac{1}{2}\mathbf{B}_1}) = e^{A_1} e^{\mathbf{B}_1} \quad (244)$$

$$\psi_2 = (e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{F}_1} e^{\frac{1}{2}\mathbf{B}_1})^\dagger (e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{F}_1} e^{\frac{1}{2}\mathbf{B}_1}) = e^{A_2} e^{\mathbf{B}_2} \quad (245)$$

Thus

$$\psi = e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{\mathbf{B}_2} \quad (246)$$

The quadri-linear map becomes a bilinear map:

$$\psi^\dagger \psi = (e^{A_1} e^{-\mathbf{B}_1} + e^{A_2} e^{-\mathbf{B}_2})(e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{\mathbf{B}_2}) \quad (247)$$

$$= e^{A_1} e^{-\mathbf{B}_1} e^{A_1} e^{\mathbf{B}_1} + e^{A_1} e^{-\mathbf{B}_1} e^{A_2} e^{\mathbf{B}_2} + e^{A_2} e^{-\mathbf{B}_2} e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{-\mathbf{B}_2} e^{A_2} e^{\mathbf{B}_2} \quad (248)$$

$$= \underbrace{e^{2A_1} + e^{2A_2}}_{\text{sum}} + \underbrace{2e^{A_1+A_2} \cos(B_1 - B_2)}_{\text{complex interference}} \quad (249)$$

6.1.7 Geometric Interference in 4D (Deep Phase Rotation)

A phase rotation on the base algebra (rather than the sub-algebra) produces a difference interference pattern. Consider a two-state wave-function:

$$\psi = \psi_1 + \psi_2 = e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{B}_2} \quad (250)$$

The sub-product part is:

$$\psi^\dagger \psi = (e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{B}_2})(e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{B}_2}) \quad (251)$$

$$= e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{B}_1} e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{B}_1} e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{B}_2} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{B}_2} e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{B}_2} e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{B}_2} \quad (252)$$

$$= e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{\mathbf{B}_2} + 2e^{\frac{1}{2}(A_1+A_2)} e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} \quad (253)$$

The final product is:

$$[\psi^\dagger \psi]_{3,4} \psi^\dagger \psi = (e^{A_1} e^{-\mathbf{B}_1} + e^{A_2} e^{-\mathbf{B}_2} + 2e^{\frac{1}{2}(A_1+A_2)} e^{-\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}) \times (e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{\mathbf{B}_2} + 2e^{\frac{1}{2}(A_1+A_2)} e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}) \quad (254)$$

$$= e^{A_1} e^{-\mathbf{B}_1} e^{A_1} e^{\mathbf{B}_1} + e^{A_1} e^{-\mathbf{B}_1} e^{A_2} e^{\mathbf{B}_2} + e^{A_1} e^{-\mathbf{B}_1} 2e^{\frac{1}{2}(A_1+A_2)} e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} + e^{A_2} e^{-\mathbf{B}_2} e^{A_1} e^{\mathbf{B}_1} + e^{A_2} e^{-\mathbf{B}_2} e^{A_2} e^{\mathbf{B}_2} + e^{A_2} e^{-\mathbf{B}_2} 2e^{\frac{1}{2}(A_1+A_2)} e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} + 2e^{\frac{1}{2}(A_1+A_2)} e^{-\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} e^{A_1} e^{\mathbf{B}_1} + 2e^{\frac{1}{2}(A_1+A_2)} e^{-\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} e^{A_2} e^{\mathbf{B}_2} + 2e^{\frac{1}{2}(A_1+A_2)} e^{-\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} 2e^{\frac{1}{2}(A_1+A_2)} e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} \quad (255)$$

$$= e^{2A_1} + e^{2A_2} + 2e^{A_1+A_2} \cos(B_1 - B_2) + e^{A_1} e^{-\mathbf{B}_1} 2e^{\frac{1}{2}(A_1+A_2)} e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} + e^{A_2} e^{-\mathbf{B}_2} 2e^{\frac{1}{2}(A_1+A_2)} e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} + 2e^{\frac{1}{2}(A_1+A_2)} e^{-\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} e^{A_1} e^{\mathbf{B}_1} + 2e^{\frac{1}{2}(A_1+A_2)} e^{-\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)} e^{A_2} e^{\mathbf{B}_2} + 4e^{A_1+A_2} \quad (256)$$

$$= \underbrace{e^{2A_1} + e^{2A_2}}_{\text{sum}} + \underbrace{2e^{A_1+A_2} \cos(B_1 - B_2)}_{\text{complex interference}} + \underbrace{2e^{\frac{1}{2}(A_1+A_2)} (e^{A_1} + e^{A_2}) \cos \frac{1}{2}(B_1 - B_2) + 4e^{A_1+A_2}}_{\text{deep phase interference}} \quad (257)$$

6.1.8 Geometric Interference in 4D (Deep Spinor Rotation)

Consider a two-state wave-function (we note that $[\mathbf{F}, \mathbf{B}] = 0$):

$$\psi = \psi_1 + \psi_2 = e^{\frac{1}{2}A_1} e^{\frac{1}{2}\mathbf{F}_1} e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2} e^{\frac{1}{2}\mathbf{F}_2} e^{\frac{1}{2}\mathbf{B}_2} \quad (258)$$

The geometric interference pattern for a full general linear transformation in 4D is given by the product:

$$[\psi^\dagger\psi]_{3,4}\psi^\dagger\psi \quad (259)$$

Let us start with the sub-product:

$$\psi^\dagger\psi = (e^{\frac{1}{2}A_1}e^{-\frac{1}{2}\mathbf{F}_1}e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2}e^{-\frac{1}{2}\mathbf{F}_2}e^{\frac{1}{2}\mathbf{B}_2})(e^{\frac{1}{2}A_1}e^{\frac{1}{2}\mathbf{F}_1}e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2}e^{\frac{1}{2}\mathbf{F}_2}e^{\frac{1}{2}\mathbf{B}_2}) \quad (260)$$

$$\begin{aligned} &= e^{\frac{1}{2}A_1}e^{-\frac{1}{2}\mathbf{F}_1}e^{\frac{1}{2}\mathbf{B}_1}e^{\frac{1}{2}A_1}e^{\frac{1}{2}\mathbf{F}_1}e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_1}e^{-\frac{1}{2}\mathbf{F}_1}e^{\frac{1}{2}\mathbf{B}_1}e^{\frac{1}{2}A_2}e^{\frac{1}{2}\mathbf{F}_2}e^{\frac{1}{2}\mathbf{B}_2} \\ &\quad + e^{\frac{1}{2}A_2}e^{-\frac{1}{2}\mathbf{F}_2}e^{\frac{1}{2}\mathbf{B}_2}e^{\frac{1}{2}A_1}e^{\frac{1}{2}\mathbf{F}_1}e^{\frac{1}{2}\mathbf{B}_1} + e^{\frac{1}{2}A_2}e^{-\frac{1}{2}\mathbf{F}_2}e^{\frac{1}{2}\mathbf{B}_2}e^{\frac{1}{2}A_2}e^{\frac{1}{2}\mathbf{F}_2}e^{\frac{1}{2}\mathbf{B}_2} \end{aligned} \quad (261)$$

$$= e^{A_1}e^{\mathbf{B}_1} + e^{A_2}e^{\mathbf{B}_2} + e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}(e^{-\frac{1}{2}\mathbf{F}_1}e^{\frac{1}{2}\mathbf{F}_2}) + e^{-\frac{1}{2}\mathbf{F}_2}e^{\frac{1}{2}\mathbf{F}_1}) \quad (262)$$

$$= e^{A_1}e^{\mathbf{B}_1} + e^{A_2}e^{\mathbf{B}_2} + e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1) \quad (263)$$

where $R = e^{\frac{1}{2}\mathbf{F}}$, and where $\tilde{R} = e^{-\frac{1}{2}\mathbf{F}}$.

The full product is:

$$\begin{aligned} [\psi^\dagger\psi]_{3,4}\psi^\dagger\psi &= \left(e^{A_1}e^{-\mathbf{B}_1} + e^{A_2}e^{-\mathbf{B}_2} + e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(-\mathbf{B}_1-\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1) \right) \\ &\quad \times \left(e^{A_1}e^{\mathbf{B}_1} + e^{A_2}e^{\mathbf{B}_2} + e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1) \right) \quad (264) \\ &= e^{A_1}e^{-\mathbf{B}_1}e^{A_1}e^{\mathbf{B}_1} + e^{A_1}e^{-\mathbf{B}_1}e^{A_2}e^{\mathbf{B}_2} + e^{A_1}e^{-\mathbf{B}_1}e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1) \\ &\quad + e^{A_2}e^{-\mathbf{B}_2}e^{A_1}e^{\mathbf{B}_1} + e^{A_2}e^{-\mathbf{B}_2}e^{A_2}e^{\mathbf{B}_2} + e^{A_2}e^{-\mathbf{B}_2}e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1) \\ &\quad + e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(-\mathbf{B}_1-\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1)e^{A_1}e^{\mathbf{B}_1} \\ &\quad + e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(-\mathbf{B}_1-\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1)e^{A_2}e^{\mathbf{B}_2} \\ &\quad + e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(-\mathbf{B}_1-\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1)e^{\frac{1}{2}(A_1+A_2)}e^{\frac{1}{2}(\mathbf{B}_1+\mathbf{B}_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1) \end{aligned} \quad (265)$$

$$= e^{2A_1} + e^{2A_2} + 2e^{A_1+A_2}\cos(B_1 - B_2) \quad (266)$$

$$+ e^{\frac{1}{2}(A_1+A_2)}(\tilde{R}_1R_2 + \tilde{R}_2R_1)(\quad (267)$$

$$e^{A_1}(e^{\frac{1}{2}(-\mathbf{B}_1+\mathbf{B}_2)} + e^{\frac{1}{2}(\mathbf{B}_1-\mathbf{B}_2)}) \quad (268)$$

$$+ e^{A_2}(e^{\frac{1}{2}(\mathbf{B}_1-\mathbf{B}_2)} + e^{\frac{1}{2}(-\mathbf{B}_1+\mathbf{B}_2)})) \quad (269)$$

$$+ e^{A_1+A_2}(\tilde{R}_1R_2 + \tilde{R}_2R_1)^2 \quad (270)$$

$$= \underbrace{e^{2A_1} + e^{2A_2}}_{\text{sum}} + \underbrace{2e^{A_1+A_2}\cos(B_1 - B_2)}_{\text{complex interference}} \\ + \underbrace{2e^{\frac{1}{2}(A_1+A_2)}(e^{A_1} + e^{A_2})(\tilde{R}_1R_2 + \tilde{R}_2R_1)(\cos\frac{1}{2}(B_1 - B_2)) + e^{A_1+A_2}(\tilde{R}_1R_2 + \tilde{R}_2R_1)^2}_{\text{deep spinor interference}} \quad (271)$$

6.1.9 Geometric Interference Experiment (Sketch)

In the case of the general linear group, the interference pattern is much more complicated than the simple cosine of the standard Born rule, but that is to be expected as it comprises the full general linear group and not just the unitary group. It accounts for the group of all geometric transformations which preserves the probability distribution ρ for a two-state general linear system.

General linear interference can be understood as a generalization of complex interference, which is recovered under a "shallow" phase rotation in 4D and under just a plain normal phase rotation in 2D. Furthermore, when all elements of the odd-sub-algebra are eliminated (by posing $\mathbf{X} \rightarrow 0$, $\mathbf{V} \rightarrow 0$), then the wave-function reduces to the geometric algebra form of the relativistic wave-function identified by David Hestenes, in terms of a spinor field.

Such reductions produce a series of interference patterns of decreasing complexity, and as such they provide a method to experimentally identify which group of geometric transformations the world obeys, using interference experiments as the identification tool. Identification of the full general linear interference pattern (with all the elements $A, \mathbf{X}, \mathbf{F}, \mathbf{V}, \mathbf{B}$) in a lab experiment would suggest a gauge-theoretical theory of gravity, whereas identification of a reduced interference pattern (produced by $A, \mathbf{F}, \mathbf{B}$) and subsequently showing a failure to observe the full general linear interference ($\mathbf{X} \rightarrow 0, \mathbf{V} \rightarrow 0$) would suggest at most spinor-level interference.

In any such case, a general experimental setup would send a particle into two distinct paths. Then, either: a) one of the paths undergoes a general linear transformation, while the other doesn't or b) both paths undergo a different general linear transformation. Then, the paths are recombined to produce an interference pattern on a screen. Depending on the nature of the transformation, a deformation of the interference pattern based on the geometry of the setup should be observed.

One can further utilize the non-commutativity of the general linear transformations to identify only the difference between complex-interference and general linear interference. One would apply the same general linear transformations to each path, but would reverse the order in which the transformations are applied. The resulting interference pattern would then be compared to a case where both paths are transformed in the same order. Then, complex-interference, as it is fully commutative, would predict the same interference pattern irrespective of the order the transformations are applied in — whereas, with general linear interference, as it is non-commutative, would predict different interference patterns.

To achieve this it may be necessary to use a three-dimensional detector, whose idealized construction is a homogeneous bath of impurities (allowing photons to 'click' anywhere within the volume of the detector), instead of a two-dimensional screen, since the opportunity for non-commutative behaviour often kicks in at three dimensions or higher.

7 Discussion

The final piece of the puzzle is to provide a mechanism for the wave-function collapse and an interpretation of quantum mechanics. Previously, we have derived from first principles the probability measure (Born rule) as the fundamental object and the wave-function as a convenient construction to work with linear transformations. In fact, we have found the Born rule as a special case of the Gibbs measure and occurs when the entropy is maximized under the constraint of a phase-invariance instrument. Let us see in more details.

In statistical physics, constraints on the entropy are interpreted as instruments acting on the system. For instance, an energy constraint on the entropy:

$$\overline{E} = \sum_{q \in \mathbb{Q}} \rho(q) E(q) \quad (272)$$

is interpreted, physically, as a energy-meter measuring the system and producing a series of energy measurement E_1, E_2, \dots converging to a average value \overline{E} . Another typical constraint is that of the volume:

$$\overline{V} = \sum_{q \in \mathbb{Q}} \rho(q) V(q) \quad (273)$$

associated to a volume-meter acting on the system and converging towards an average volume value \overline{V} , also by producing a sequence of measurements of the volume V_1, V_2, \dots . With these two constraints, the typical system of statistical physics is obtained, and its Gibbs measure is:

$$\rho(q, \beta, p) = \frac{1}{Z} \exp(-\beta(E(q) + pV(q))) \quad (274)$$

Comparatively, in the present recovery of QM, the statistical physics interpretation is also adopted. Instead of an energy-meter or a volume-meter, we have a phase-invariant meter, and the constraint is given as follows:

$$\text{tr} \begin{pmatrix} 0 & -\overline{b} \\ \overline{b} & 0 \end{pmatrix} = \sum_{q \in \mathbb{Q}} \rho(q) \text{tr} \begin{pmatrix} 0 & -b(q) \\ b(q) & 0 \end{pmatrix} \quad (275)$$

The usage of the trace enforces the phase-invariance of the instrument. Yet, and quite simply, maximizing the entropy under this constraints produces the probability measure of the wave-function including the Born rule. This is the true origin of the Born rule, here reported for the first time. The interpretation simply becomes that of an instrument performing a sequence of measurement on the system such that an average value is obtained, but instead of the simpler scalar instruments used in statistical physics, here we have a phase-invariant instrument. This instrument is responsible for the quantum mechanical behaviour

associated with the wave-function. What is an example of such a detector; quite simply a photo-counter would be one. Such an instrument produces a sequence of incidences ('clicks') as photons are detected and "advanced features" such as an interference pattern is a consequence of this phase-invariance.

Now, the correct interpretation of quantum mechanics is simply the consequence of the existence of phase-invariant photon counters producing series of incidence counts which constrains the entropy of a physical system. Finally, maximizing said entropy using the typical tools of statistical physics produces the Born rule as a special case of the Gibbs ensemble. The interpretation of QM and statistical physics are now unified as the same interpretation.

It is not a collapse that occurs but instead a derivation of the wave-function from a sequence of incidences recorded by phase-invariant instruments. This is the natural interpretation; as a system of statistical physics whose entropy is maximized under proper phase-invariant constraint. Consequently, the interpretation does not admit a collapse problem, anymore than a coin landing on head creates a probability collapse problem (from 50-head/50-tail to 100-head after landing) for classical statistics.

8 Conclusion

We believe the formal system of science here-in presented to be a more powerful formulation of the laws of physics, because it can provide mathematical definitions for experimental data, the observer, it can derive the laws of physics from the scientific method, provide an account for the wave-function collapse, identify the wave-function as a special form of the Gibbs ensemble, account for the origin of the Born rule and finally provide an interpretation of quantum physics matching that of statistical physics.

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A Notation

S will denote the entropy, \mathcal{A} the action, L the Lagrangian, and \mathcal{L} the Lagrangian density. Sets, unless a prior convention assigns it another symbol, will be written using the blackboard bold typography (ex: $\mathbb{L}, \mathbb{W}, \mathbb{Q}$, etc.). Matrices will be in bold upper case (ex: \mathbf{A}, \mathbf{B}), whereas vectors and multi-vectors will be in bold lower case (ex: $\mathbf{u}, \mathbf{v}, \mathbf{g}$) and most other constructions (ex.: scalars, functions) will have plain typography (ex. a, A). The identity matrix is I , the unit pseudo-scalar (of geometric algebra) is \mathbf{I} and the imaginary number is i . The Dirac gamma matrices are $\gamma_0, \gamma_1, \gamma_2, \gamma_3$ and the Pauli matrices are $\sigma_x, \sigma_y, \sigma_z$. The basis elements of an arbitrary curvilinear geometric basis will be denoted $\mathbf{e}_0, \mathbf{e}_1, \mathbf{e}_2, \dots, \mathbf{e}_n$ (such that $\mathbf{e}_\nu \cdot \mathbf{e}_\mu = g_{\mu\nu}$) and if they are orthonormal as $\hat{\mathbf{x}}_0, \hat{\mathbf{x}}_1, \hat{\mathbf{x}}_2, \dots, \hat{\mathbf{x}}_n$ (such that $\hat{\mathbf{x}}_\mu \cdot \hat{\mathbf{x}}_\nu = \eta_{\mu\nu}$). The asterisk z^* denotes the complex conjugate of z , and the dagger \mathbf{A}^\dagger denotes the conjugate transpose of \mathbf{A} . A geometric algebra of m dimensions over a field \mathbb{F} is noted as $\mathbb{G}(m, \mathbb{F})$. The grades of a multi-vector will be denoted as $\langle \mathbf{v} \rangle_k$. Specifically, $\langle \mathbf{v} \rangle_0$ is a scalar, $\langle \mathbf{v} \rangle_1$ is a vector, $\langle \mathbf{v} \rangle_2$ is a bi-vector, $\langle \mathbf{v} \rangle_{n-1}$ is a pseudo-vector and $\langle \mathbf{v} \rangle_n$ is a pseudo-scalar. Furthermore, a scalar and a vector $\langle \mathbf{v} \rangle_0 + \langle \mathbf{v} \rangle_1$ is a para-vector, and a combination of even grades ($\langle \mathbf{v} \rangle_0 + \langle \mathbf{v} \rangle_2 + \langle \mathbf{v} \rangle_4 + \dots$) or odd grades ($\langle \mathbf{v} \rangle_1 + \langle \mathbf{v} \rangle_3 + \dots$) are even-multi-vectors or odd-multi-vectors, respectively. The commutator is defined as $[\mathbf{A}, \mathbf{B}] := \mathbf{AB} - \mathbf{BA}$ and the anti-commutator as $\{\mathbf{A}, \mathbf{B}\} := \mathbf{AB} + \mathbf{BA}$. We use the symbol \cong to relate two sets that are related by a group isomorphism. We denote the Hadamard product, or element-wise multiplication, of two matrices using \odot , and is written for instance as $\mathbf{M} \odot \mathbf{P}$, and for a multivector as

$\mathbf{u} \odot \mathbf{v}$; for instance: $(a_0 + x_0 \hat{\mathbf{x}} + y_0 \hat{\mathbf{y}} + b_0 \hat{\mathbf{x}} \wedge \hat{\mathbf{y}}) \odot (a_1 + x_1 \hat{\mathbf{x}} + y_1 \hat{\mathbf{y}} + b_1 \hat{\mathbf{x}} \wedge \hat{\mathbf{y}})$ would equal $a_0 a_1 + x_0 x_1 \hat{\mathbf{x}} + y_0 y_1 \hat{\mathbf{y}} + b_0 b_1 \hat{\mathbf{x}} \wedge \hat{\mathbf{y}}$.

B Lagrange equation

The Lagrangian equation to maximize is:

$$\mathcal{L}(\rho, \alpha, \tau) = -k_B \sum_{q \in \mathbb{Q}} \rho(q) \ln \rho(q) + \alpha \left(1 - \sum_{q \in \mathbb{Q}} \rho(q) \right) + \tau \operatorname{tr} \left(\bar{\mathbf{M}} - \sum_{q \in \mathbb{Q}} \rho(q) \mathbf{M}(q) \right) \quad (276)$$

where α and τ are the Lagrange multipliers. We note the usage of the trace operator for the geometric constraint such that a scalar-valued equation is maximized. Maximizing this equation for ρ by posing $\frac{\partial \mathcal{L}}{\partial \rho(p)} = 0$, where $p \in \mathbb{Q}$, we obtain:

$$\frac{\partial \mathcal{L}}{\partial \rho(p)} = -k_B \ln \rho(p) - k_B - \alpha - \tau \operatorname{tr} \mathbf{M}(p) \quad (277)$$

$$0 = k_B \ln \rho(p) + k_B + \alpha + \tau \operatorname{tr} \mathbf{M}(p) \quad (278)$$

$$\implies \ln \rho(p) = \frac{1}{k_B} (-k_B - \alpha - \tau \operatorname{tr} \mathbf{M}(p)) \quad (279)$$

$$\implies \rho(p) = \exp \left(\frac{-k_B - \alpha}{k_B} \right) \exp \left(-\frac{\tau}{k_B} \operatorname{tr} \mathbf{M}(p) \right) \quad (280)$$

$$= \frac{1}{Z} \det \exp \left(-\frac{\tau}{k_B} \mathbf{M}(p) \right) \quad (281)$$

where Z is obtained as follows:

$$1 = \sum_{q \in \mathbb{Q}} \exp \left(\frac{-k_B - \alpha}{k_B} \right) \exp \left(-\frac{\tau}{k_B} \operatorname{tr} \mathbf{M}(q) \right) \quad (282)$$

$$\implies \left(\exp \left(\frac{-k_B - \alpha}{k_B} \right) \right)^{-1} = \sum_{q \in \mathbb{Q}} \exp \left(-\frac{\tau}{k_B} \operatorname{tr} \mathbf{M}(q) \right) \quad (283)$$

$$Z := \sum_{q \in \mathbb{Q}} \det \exp \left(-\frac{\tau}{k_B} \mathbf{M}(q) \right) \quad (284)$$

We note that the Trace in the exponential drops down to a determinant, via the relation $\det \exp A \equiv \exp \operatorname{tr} A$.

B.1 Multiple constraints

Consider a set of constraints:

$$\overline{\mathbf{M}}_1 = \sum_{q \in \mathbb{Q}} \rho(q) \mathbf{M}_1(q) \quad (285)$$

$$\vdots \quad (286)$$

$$\overline{\mathbf{M}}_n = \sum_{q \in \mathbb{Q}} \rho(q) \mathbf{M}_n(q) \quad (287)$$

Then the Lagrange equation becomes:

$$\begin{aligned} \mathcal{L} = -k_B \sum_{q \in \mathbb{Q}} \rho(q) \ln \rho(q) + \alpha \left(1 - \sum_{q \in \mathbb{Q}} \rho(q) \right) + \tau_1 \operatorname{tr} \left(\overline{\mathbf{M}}_1 - \sum_{q \in \mathbb{Q}} \rho(q) \mathbf{M}_1(q) \right) + \dots \\ + \tau_n \operatorname{tr} \left(\overline{\mathbf{M}}_n - \sum_{q \in \mathbb{Q}} \rho(q) \mathbf{M}_n(q) \right) \end{aligned} \quad (288)$$

and the measure references all n constraints:

$$\rho(q) = \frac{1}{Z} \det \exp \left(-\frac{\tau_1}{k_B} \mathbf{M}_1(q) - \dots - \frac{\tau_n}{k_B} \mathbf{M}_n(q) \right) \quad (289)$$

B.2 Multiple constraints - General Case

In the general case of a multi-constraint system, each entry of the matrix corresponds to a constraint:

$$\overline{M}_{00} \begin{pmatrix} 1 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 0 \end{pmatrix} = \sum_{q \in \mathbb{Q}} \rho(q) M_{00}(q) \begin{pmatrix} 1 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 0 \end{pmatrix} \quad (290)$$

$$\vdots \quad (291)$$

$$\overline{M}_{01} \begin{pmatrix} 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 0 \end{pmatrix} = \sum_{q \in \mathbb{Q}} \rho(q) M_{01}(q) \begin{pmatrix} 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 0 \end{pmatrix} \quad (292)$$

$$\vdots \quad (293)$$

$$\overline{M}_{nn} \begin{pmatrix} 0 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 1 \end{pmatrix} = \sum_{q \in \mathbb{Q}} \rho(q) M_{nn}(q) \begin{pmatrix} 0 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 1 \end{pmatrix} \quad (294)$$

For a $n \times n$ matrix, there are n^2 constraints.

The probability measure which maximizes the entropy is as follows:

$$\rho(q) = \frac{1}{Z} \det \exp \left(-\frac{1}{k_B} \boldsymbol{\tau} \odot \mathbf{M}(q) \right) \quad (295)$$

where $\boldsymbol{\tau}$ is a matrix of Lagrange multipliers, and \odot , the element-wise multiplication, assigns the corresponding Lagrange multiplier to each constraint.