A true mathematical statement Ψ with the predicate \mathcal{K} of the current mathematical knowledge, where Ψ may be false in the future and does not express the current knowledge about the provability of any mathematical statements Ψ_1, \ldots, Ψ_n without \mathcal{K}

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ABSTRACT. The theorem of Royer and Case states that there exists a limit-computable function $\beta_1 : \mathbb{N} \to \mathbb{N}$ which eventually dominates every computable function $\delta_1 : \mathbb{N} \to \mathbb{N}$. We present an alternative proof of this theorem. \mathcal{K} denotes both the knowledge predicate satisfied by every currently known theorem and the finite set of all currently known theorems. The set \mathcal{K} is time-dependent and publicly available. Any theorem of any mathematician from past or present forever belongs to \mathcal{K} . We prove: (1) there exists a limit-computable function $f : \mathbb{N} \to \mathbb{N}$ of unknown computability which eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation. We present both constructive and non-constructive proof of (1). Statement (1) claims that there exists a function $f : \mathbb{N} \to \mathbb{N}$ such that (*f* is computable in the limit) $\land (\neg \mathcal{K}(f \text{ is computable})) \land (f$ eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation). Since Martin Davis' conjecture on single-fold Diophantine representations disproves Statement (1), Statement (1) justifies the title of the article.

Key words and phrases: eventual domination, limit-computable function, predicate \mathcal{K} of the current mathematical knowledge, single-fold Diophantine representation, time-dependent truth in mathematics with the predicate \mathcal{K} .

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Let

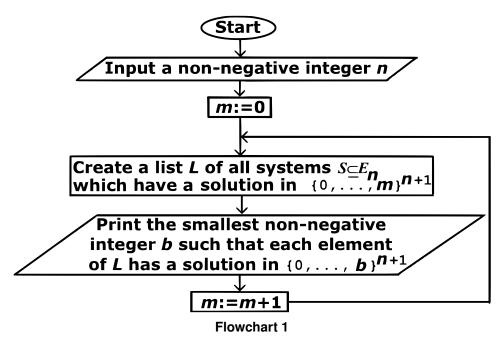
 $E_n = \{1 = x_k, \ x_i + x_j = x_k, \ x_i \cdot x_j = x_k : \ i, j, k \in \{0, \dots, n\}\}$

1. CLASSICAL MATHEMATICS

Theorem 1. ([4, p. 118]). There exists a limit-computable function $\beta_1 : \mathbb{N} \to \mathbb{N}$ which eventually dominates every computable function $\delta_1 : \mathbb{N} \to \mathbb{N}$.

We present an alternative proof of Theorem 1. For every $n \in \mathbb{N}$, we define $\beta_1(n)$ as the smallest $b \in \mathbb{N}$ such that if a system of equations $S \subseteq E_n$ has a solution in \mathbb{N}^{n+1} , then this solution belongs to $\{0, \ldots, b\}^{n+1}$. The function $\beta_1 : \mathbb{N} \to \mathbb{N}$ is computable in the limit and eventually dominates every computable function $\delta_1 : \mathbb{N} \to \mathbb{N}$, see [5]. Flowchart 1 describes a semi-algorithm which computes $\beta_1(n)$ in the limit.

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A semi-algorithm which computes $\beta_1(n)$ in the limit

Conjecture 1. ([1, pp. 341–342], [2, p. 42], [3]). Every listable set $\mathcal{X} \subseteq \mathbb{N}^k$ $(k \in \mathbb{N} \setminus \{0\})$ has a single-fold Diophantine representation.

Let Φ denote the following statement: the function $\mathbb{N} \ni n \to 2^n \in \mathbb{N}$ eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation. For $n \in \mathbb{N}$, let

$$g(n) = \begin{cases} 2^n, & \text{if } \Phi \text{ holds} \\ \beta_1(n), & \text{otherwise} \end{cases}$$

The function $g: \mathbb{N} \to \mathbb{N}$ is computable if and only if Φ holds. Currently,

 $(\neg \mathcal{K}(\Phi)) \land (\neg \mathcal{K}(\neg \Phi)) \land (\neg \mathcal{K}(g \ is \ computable)) \land (\neg \mathcal{K}(g \ is \ uncomputable))$

Lemma 1. The function g is computable in the limit and eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation.

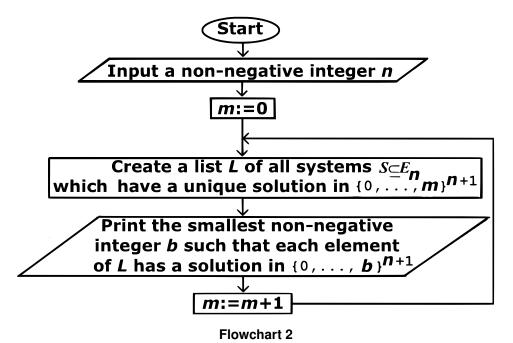
Proof. It follows from Theorem 1.

For every $n \in \mathbb{N}$, we define $\beta(n)$ as the smallest $b \in \mathbb{N}$ such that if a system of equations $S \subseteq E_n$ has a unique solution in \mathbb{N}^{n+1} , then this solution belongs to $\{0, \ldots, b\}^{n+1}$.

Theorem 2. The function $\beta : \mathbb{N} \to \mathbb{N}$ is computable in the limit and eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation.

Proof. This is proved in [5]. The term "dominated" in the title of [5] means "eventually dominated". Flowchart 2 describes a semi-algorithm which computes $\beta(n)$ in the limit.

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A semi-algorithm which computes $\beta(n)$ in the limit

2. Mathematics with the predicate ${\cal K}$ of the current mathematical knowledge

 \mathcal{K} denotes both the knowledge predicate satisfied by every currently known theorem and the finite set of all currently known theorems. The set \mathcal{K} is time-dependent and publicly available. Any theorem of any mathematician from past or present forever belongs to \mathcal{K} .

Let \mathcal{T} denote the set of twin primes. Let Γ denote the following true statement: *it is unknown whether or not there are infinitely many twin primes.*

Proposition 1. The statement Γ does not justify the title of the article.

Proof. The statement Γ expresses that

 $(\neg \mathcal{K}(\operatorname{card}(\mathcal{T}) = \omega)) \land (\neg \mathcal{K}(\operatorname{card}(\mathcal{T}) < \omega))$

This conjunction expresses the current knowledge about the provability of the statements $\operatorname{card}(\mathcal{T}) = \omega$ and $\operatorname{card}(\mathcal{T}) < \omega$.

Statement 1. There exists a limit-computable function $f : \mathbb{N} \to \mathbb{N}$ of unknown computability which eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation.

Proof. Statement 1 follows constructively from Theorem 2 by taking $f = \beta$ and the following conjunction:

 $(\neg \mathcal{K}(\beta \text{ is computable})) \land (\neg \mathcal{K}(\beta \text{ is uncomputable}))$

Statement 1 follows non-constructively from Lemma 1 by taking f = g and the following conjunction:

 $(\neg \mathcal{K}(g \text{ is computable})) \land (\neg \mathcal{K}(g \text{ is uncomputable}))$

Proposition 2. Statement 1 justifies the title of the article.

Proof. Statement 1 claims that there exists a function $f : \mathbb{N} \to \mathbb{N}$ such that

(*f* is computable in the limit) $\land (\neg \mathcal{K}(f \text{ is computable})) \land$

 $(\neg \mathcal{K}(f \text{ is uncomputable})) \land (f \text{ eventually dominates every function } \delta : \mathbb{N} \to \mathbb{N}$

with a single-fold Diophantine representation)

Conjecture 1 disproves Statement 1.

Since the function β_1 in Theorem 1 is not computable, Statement 1 does not follow from Theorem 1. Ignoring the epistemic condition in Statement 1, Statement 1 follows from Theorem 1 by taking $f = \beta_1$.

In [7], the author showed that the predicate \mathcal{K} non-trivially extends constructive mathematics, see also [6], [8], [9].

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