A true and falsifiable statement Ψ with the predicate $\mathcal{K}$ of the current mathematical knowledge, where Ψ is mathematically interesting and does not express what is proved or unproved in mathematics without $\mathcal{K}$

Apoloniusz Tyszka

ABSTRACT. The theorem of Royer and Case states that there exists a limit-computable function $\beta_1 : \mathbb{N} \to \mathbb{N}$ which eventually dominates every computable function $\delta_1 : \mathbb{N} \to \mathbb{N}$. We present an alternative proof of this theorem. $\mathcal{K}$ denotes both the knowledge predicate satisfied by every currently known theorem and the finite set of all currently known theorems. The set $\mathcal{K}$ is time-dependent and publicly available. Any theorem of any mathematician from past or present forever belongs to $\mathcal{K}$. We prove: `(1)` there exists a limit-computable function $f : \mathbb{N} \to \mathbb{N}$ of unknown computability which eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation. We present both constructive and non-constructive proof of `(1)`. Statement `(1)` claims that there exists a function $f : \mathbb{N} \to \mathbb{N}$ such that (f is computable in the limit) $\wedge$ ($\neg\mathcal{K}$(f is computable)) $\wedge$ ($\neg\mathcal{K}$(f is uncomputable)) $\wedge$ (f eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation). Since Martin Davis' conjecture on single-fold Diophantine representations disproves Statement `(1)`, Statement `(1)` justifies the title of the article.

Key words and phrases: eventual domination, limit-computable function, predicate $\mathcal{K}$ of the current mathematical knowledge, single-fold Diophantine representation, time-dependent truth in mathematics with the predicate $\mathcal{K}$.

2020 Mathematics Subject Classification: 03F65.

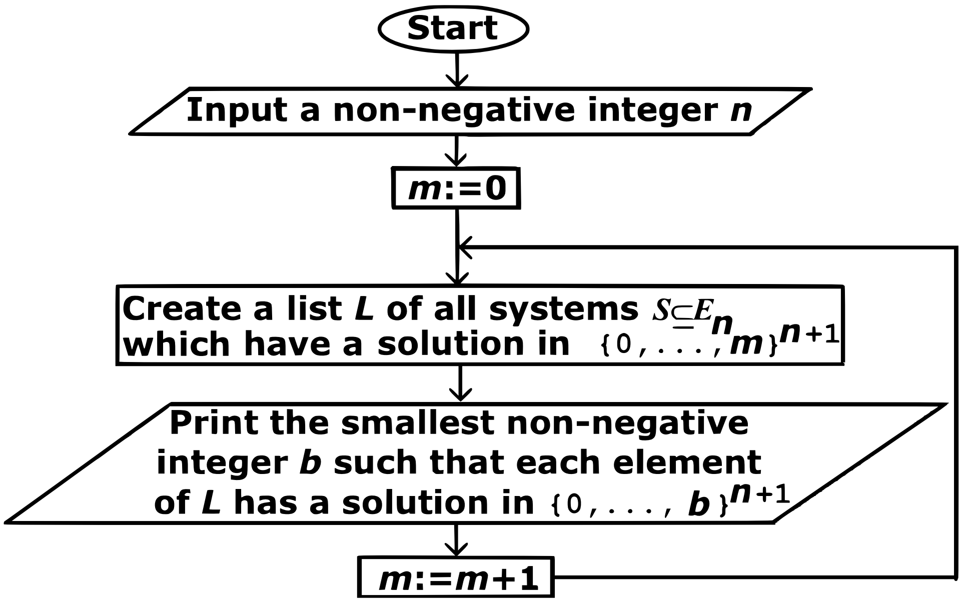
$\mathcal{K}$ denotes both the knowledge predicate satisfied by every currently known theorem and the finite set of all currently known theorems. The set $\mathcal{K}$ is time-dependent and publicly available. Any theorem of any mathematician from past or present forever belongs to $\mathcal{K}$. For $n \in \mathbb{N}$, let

$$E_n = \{1 = x_k,\ x_i + x_j = x_k,\ x_i \cdot x_j = x_k :\ i, j, k \in \{0, \ldots, n\}\}$$

1. CLASSICAL MATHEMATICS

Theorem 1. *(*[5, p. 118]*). There exists a limit-computable function $\beta_1 : \mathbb{N} \to \mathbb{N}$ which eventually dominates every computable function $\delta_1 : \mathbb{N} \to \mathbb{N}$.*

We present an alternative proof of Theorem 1. For every $n \in \mathbb{N}$, we define $\beta_1(n)$ as the smallest $b \in \mathbb{N}$ such that if a system of equations $S \subseteq E_n$ has a solution in $\mathbb{N}^{n+1}$, then this solution belongs to $\{0, \ldots, b\}^{n+1}$. The function $\beta_1 : \mathbb{N} \to \mathbb{N}$ is computable in the limit and eventually dominates every computable function $\delta_1 : \mathbb{N} \to \mathbb{N}$, see [6]. Flowchart 1 describes a semi-algorithm which computes $\beta_1(n)$ in the limit.



Flowchart 1

A semi-algorithm which computes $\beta_1(n)$ in the limit

Conjecture 1. *(*[1, pp. 341–342], [3, p. 42], [4, p. 745]*). Every listable set* $\mathcal{X} \subseteq \mathbb{N}^k$ $(k \in \mathbb{N} \setminus \{0\})$ *has a single-fold Diophantine representation.*

Let Φ denote the following statement: *the function* $\mathbb{N} \ni n \to 2^n \in \mathbb{N}$ *eventually dominates every function* $\delta : \mathbb{N} \to \mathbb{N}$ *with a single-fold Diophantine representation*. For $n \in \mathbb{N}$, let

$$g(n) = \begin{cases} 2^n, & \text{if } \Phi \text{ holds} \\ \beta_1(n), & \text{otherwise} \end{cases}$$

The function $g : \mathbb{N} \to \mathbb{N}$ is computable if and only if Φ holds. Currently,

$$(\neg\mathcal{K}(\Phi)) \wedge (\neg\mathcal{K}(\neg\Phi)) \wedge (\neg\mathcal{K}(g \ is \ computable)) \wedge (\neg\mathcal{K}(g \ is \ uncomputable))$$

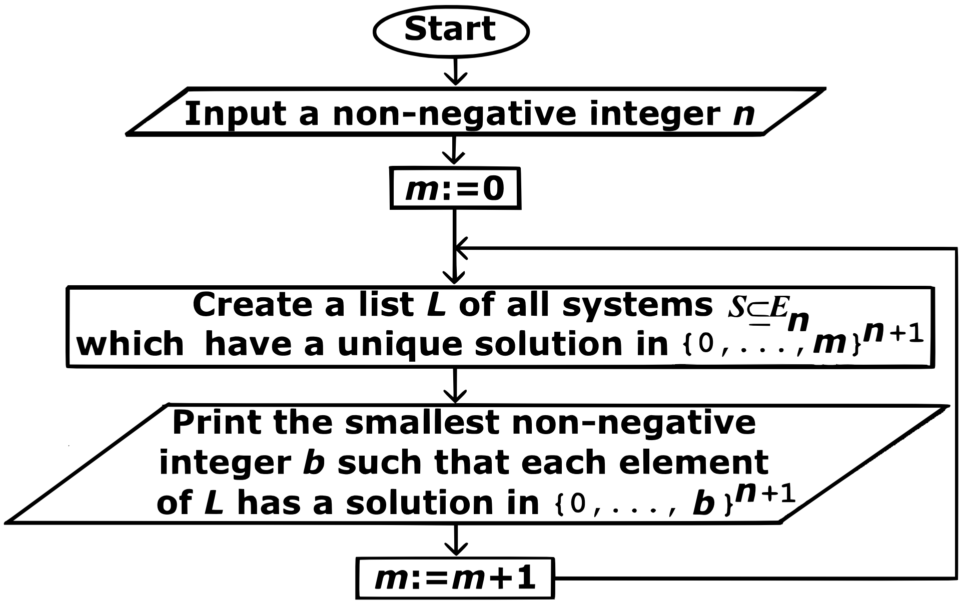
Lemma 1. *The function* g *is computable in the limit and eventually dominates every function* $\delta : \mathbb{N} \to \mathbb{N}$ *with a single-fold Diophantine representation.*

Proof. It follows from Theorem 1. $\square$

For every $n \in \mathbb{N}$, we define $\beta(n)$ as the smallest $b \in \mathbb{N}$ such that if a system of equations $S \subseteq E_n$ has a unique solution in $\mathbb{N}^{n+1}$, then this solution belongs to $\{0, \ldots, b\}^{n+1}$.

Theorem 2. *The function* $\beta : \mathbb{N} \to \mathbb{N}$ *is computable in the limit and eventually dominates every function* $\delta : \mathbb{N} \to \mathbb{N}$ *with a single-fold Diophantine representation.*

Proof. This is proved in [6]. The term *"dominated"* in the title of [6] means *"eventually dominated"*. Flowchart 2 describes a semi-algorithm which computes $\beta(n)$ in the limit.



Flowchart 2

A semi-algorithm which computes $\beta(n)$ in the limit

$\square$

2. Mathematics with the predicate $\mathcal{K}$ of the current mathematical knowledge

Let Γ denote the following true statement: *β is computable in the limit, the computability of β is unknown, and β eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation*.

Proposition 1. *The statement Γ does not justify the title of the article.*

Proof. The statement Γ expresses that

$$\mathcal{K}(\beta \text{ is computable in the limit}) \wedge (\neg\mathcal{K}(\beta \text{ is computable})) \wedge$$

$$(\neg\mathcal{K}(\beta \text{ is uncomputable})) \wedge \mathcal{K}(\beta \text{ eventually dominates every function } \delta : \mathbb{N} \to \mathbb{N}$$

$$\text{with a single-fold Diophantine representation})$$

$\square$

The analogous proposition holds for the function g instead of β.

Statement 1. *There exists a limit-computable function $f : \mathbb{N} \to \mathbb{N}$ of unknown computability which eventually dominates every function $\delta : \mathbb{N} \to \mathbb{N}$ with a single-fold Diophantine representation.*

Proof. Statement 1 follows constructively from Theorem 2 by taking $f = \beta$ and the following conjunction:

$$(\neg\mathcal{K}(\beta \text{ is computable})) \ \wedge \ (\neg\mathcal{K}(\beta \text{ is uncomputable}))$$

Statement 1 follows non-constructively from Lemma 1 by taking $f = g$ and the following conjunction:

$$(\neg\mathcal{K}(g \text{ is computable})) \ \wedge \ (\neg\mathcal{K}(g \text{ is uncomputable}))$$

$\square$

Proposition 2. *Statement 1 justifies the title of the article.*

Proof. Statement 1 claims that there exists a function $f : \mathbb{N} \to \mathbb{N}$ such that

$$(f \text{ is computable in the limit}) \wedge (\neg\mathcal{K}(f \text{ is computable})) \wedge$$

$$(\neg\mathcal{K}(f \text{ is uncomputable})) \wedge (f \text{ eventually dominates every function } \delta : \mathbb{N} \to \mathbb{N}$$

$$\text{with a single-fold Diophantine representation})$$

Conjecture 1 disproves Statement 1. $\square$

Since the function β_1 in Theorem 1 is not computable, Statement 1 does not follow from Theorem 1. Ignoring the epistemic condition in Statement 1, Statement 1 follows from Theorem 1 by taking $f = \beta_1$.

Statement 2. *There exists a positive integer b such that it is unknown whether or not the equation $x! + b = y^2$ has infinitely many solutions in positive integers x and y.*

Proof. It holds for $b = 1$, see [2]. $\square$

Proposition 3. *Statement 2 partially justifies the title of the article because it is widely known for $b = 1$.*

Proof. Statement 2 claims that there exists a positive integer b such that

$$\neg\mathcal{K}(\text{the equation } x! + b = y^2 \text{ has infinitely}$$

$$\text{many solutions in positive integers } x \text{ and } y) \wedge$$

$$\neg\mathcal{K}(\text{the equation } x! + b = y^2 \text{ has at most finitely}$$

$$\text{many solutions in positive integers } x \text{ and } y)$$

A weak form of Szpiro's conjecture implies that the equation $x! + b = y^2$ has at most finitely many solutions in positive integers x and y when $b \in \mathbb{N} \setminus \{0\}$, see [2]. $\square$

In [8], the author showed that the predicate $\mathcal{K}$ non-trivially extends constructive mathematics, see also [7], [9], [10].

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Apoloniusz Tyszka
Technical Faculty
Hugo Kołłątaj University
Balicka 116B, 30-149 Kraków, Poland
E-mail: *rttyszka@cyf-kr.edu.pl*