## Open problems that concern computable sets $X \subseteq \mathbb{N}$ and cannot be stated formally as they refer to the current mathematical knowledge and require that the finiteness (infiniteness) of Xremains conjectured

Sławomir Kurpaska, Apoloniusz Tyszka

#### Abstract

Let  $\beta = ((24!)!)!$ , and let  $\mathcal{P}_{n^2+1}$  denote the set of all primes of the form  $n^2 + 1$ . Let  $\mathcal{M}$  denote the set of all positive multiples of elements of the set  $\mathcal{P}_{n^2+1} \cap (\beta, \infty)$ . The set  $X = \{0, \ldots, \beta\} \cup \mathcal{M}$ satisfies the following conditions: (1) card(X) is greater than a huge positive integer and it is conjectured that X is infinite, (2) we do not know any algorithm deciding the finiteness of X, (3) a known and short algorithm for every  $n \in \mathbb{N}$  decides whether or not  $n \in X$ , (4) a known and short algorithm returns an integer n such that X is infinite if and only if X contains an element greater than n. The following problem is open: simply define a set  $X \subseteq \mathbb{N}$  such that X satisfies conditions (1)-(4), and for every finite set  $\mathcal{T}$ , we do not know any definition of  $X \setminus \mathcal{T}$ simpler than the definition of X (5). Let f(3) = 4, and let f(n+1) = f(n)! for every integer  $n \ge 3$ . For an integer  $n \ge 3$ , let  $\Psi_n$  denote the following statement: if a system of equations  $S \subseteq \{x_i! = x_{i+1} : 1 \le i \le n-1\} \cup \{x_i \cdot x_j = x_{j+1} : 1 \le i \le j \le n-1\}$  has only finitely many solutions in positive integers  $x_1, \ldots, x_n$ , then each such solution  $(x_1, \ldots, x_n)$  satisfies  $x_1, \ldots, x_n \le f(n)$ . We prove that for every statement  $\Psi_n$  the bound f(n) cannot be decreased. The author's guess is that the statements  $\Psi_3, \ldots, \Psi_9$  are true. We prove that the statement  $\Psi_9$  implies that the set X of all non-negative integers k whose number of digits belongs to  $\mathcal{P}_{n^2+1}$  satisfies conditions (1)-(5).

**Key words and phrases:** Alexander Zenkin's super-induction method, arithmetical operations on huge integers cannot be performed by any physical process, computable set  $X \subseteq \mathbb{N}$  whose finiteness remains conjectured, computable set  $X \subseteq \mathbb{N}$  whose infiniteness remains conjectured.

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#### **1** Basic definitions and lemmas

**Definition 1.** *Let*  $\beta = ((24!)!)!$ .

**Lemma 1.**  $\beta \approx 10^{10} 10^{25.16114896940657}$ 

*Proof.* We ask Wolfram Alpha at http://wolframalpha.com.

**Definition 2.** We say that an integer  $m \ge -1$  is a threshold number of a set  $X \subseteq \mathbb{N}$ , if X is infinite if and only if X contains an element greater than m, cf. [10] and [11].

If a set  $X \subseteq \mathbb{N}$  is empty or infinite, then any integer  $m \ge -1$  is a threshold number of X. If a set  $X \subseteq \mathbb{N}$  is non-empty and finite, then the all threshold numbers of X form the set  $\{\max(X), \max(X) + 1, \max(X) + 2, \ldots\}$ .

**Definition 3.** We say that a non-negative integer *m* is a weak threshold number of a set  $X \subseteq \mathbb{N}$ , if X is infinite if and only if card(X) > m.

**Theorem 1.** For every  $X \subseteq \mathbb{N}$ , if an integer  $m \ge -1$  is a threshold number of X, then m + 1 is a weak threshold number of X.

*Proof.* For every  $X \subseteq \mathbb{N}$ , if  $m \in [-1, \infty) \cap \mathbb{Z}$  and  $\operatorname{card}(X) > m + 1$ , then  $X \cap [m + 1, \infty) \neq \emptyset$ .

We do not know any weak threshold number of the set of all primes of the form  $n^2 + 1$ . The same is true for the sets

$$\left\{n \in \mathbb{N} : 2^{2^n} + 1 \text{ is composite}\right\}$$

and

$${n \in \mathbb{N} : n! + 1 \text{ is a square}}$$

**Lemma 2.** For every positive integers x and y,  $x! \cdot y = y!$  if and only if

$$(x + 1 = y) \lor (x = y = 1)$$

**Lemma 3.** (Wilson's theorem, [1, p. 89]). For every integer  $x \ge 2$ , x is prime if and only if x divides (x-1)! + 1.

#### 2 Open Problems 1 and 2

The following two open problems cannot be stated formally as they refer to the current mathematical knowledge.

**Open Problem 1.** Simply define a set  $X \subseteq \mathbb{N}$  that satisfies the following conditions:

(1)  $\operatorname{card}(X)$  is greater than a huge positive integer and it is conjectured that X is infinite,

(2) we do not know any algorithm deciding the finiteness of X,

(3) a known and short algorithm for every  $n \in \mathbb{N}$  decides whether or not  $n \in X$ ,

(4•) a known and short algorithm returns an integer n such that X is infinite if and only if card(X) > n, (5) for every finite set T, we do not know any definition of  $X \setminus T$  simpler than the definition of X.

**Open Problem 2.** Simply define a set  $X \subseteq \mathbb{N}$  such that X satisfies conditions (1)-(3), (5), and a known and short algorithm returns an integer n such that X is infinite if and only if X contains an element greater than n (4).

**Theorem 2.** Open Problem 2 claims more than Open Problem 1.

*Proof.* By Theorem 1, condition (4) implies condition  $(4\bullet)$ .

#### **3** Two partial solutions to Open Problem 2

Edmund Landau's conjecture states that the set  $\mathcal{P}_{n^2+1}$  of all primes of the form  $n^2 + 1$  is infinite, see [4, pp. 37–38] and [7]. Let  $\mathcal{M}$  denote the set of all positive multiples of elements of the set  $\mathcal{P}_{n^2+1} \cap (\beta, \infty)$ .

**Theorem 3.** The set  $X = \{0, ..., \beta\} \cup \mathcal{M}$  satisfies conditions (1)-(4).

*Proof.* Condition (1) holds as  $\operatorname{card}(X) > \beta$  and the set  $\mathcal{P}_{n^2+1}$  is conjecturally infinite. By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of  $\mathcal{P}_{n^2+1}$  is greater than  $\beta$ . Thus condition (2) holds. Condition (3) holds trivially. Since the set  $\mathcal{M}$  is empty or infinite, the integer  $\beta$  is a threshold number of  $\mathcal{X}$ . Thus condition (4) holds.

Let  $[\cdot]$  denote the integer part function.

**Lemma 4.** For every non-negative integer n,  $\left[\frac{3n-3\beta+3}{3n-3\beta+2}\right]$  equals 0 or 1. The first case holds when  $n \leq \beta - 1$ . The second case holds when  $n \geq \beta$ .

Lemma 5. The function

$$\mathbb{N} \cap [\beta, \infty) \ni n \xrightarrow{\theta} \beta + n - \left[\sqrt{n}\right]^2 \in \mathbb{N} \cap [\beta, \infty)$$

takes every integer value  $k \ge \beta$  infinitely many times.

*Proof.* Let  $t = k - \beta$ . The equality  $\theta(n) = k$  holds for every

$$n \in \left\{ (t+0)^2 + t, \ (t+1)^2 + t, \ (t+2)^2 + t, \ldots) \right\} \cap [\beta, \infty)$$

**Theorem 4.** The set  $X = \left\{ n \in \mathbb{N} : 2 + \left[ \frac{3n - 3\beta + 3}{3n - 3\beta + 2} \right] \cdot \left( \left( \beta + n - \left[ \sqrt{n} \right]^2 \right)^2 - 1 \right) \text{ is prime} \right\}$  satisfies conditions (1)-(4).

*Proof.* Condition (3) holds trivially. By Lemma 4,  $X = \{0, ..., \beta - 1\} \cup \mathcal{H}$ , where

$$\mathcal{H} = \left\{ n \in \mathbb{N} \cap [\beta, \infty) : \left( \beta + n - \left[ \sqrt{n} \right]^2 \right)^2 + 1 \text{ is prime} \right\}$$

By Lemma 5, the set  $\mathcal{H}$  is empty or infinite. The second case holds when

$$\exists k \in \mathbb{N} \cap [\beta, \infty) \ k^2 + 1 \text{ is prime}$$
(6)

The equality  $X = \{0, ..., \beta - 1\} \cup \mathcal{H}$  and the last two sentences imply that  $\beta - 1$  is a threshold number of X and conditions (1) and (4) hold. Condition (2) holds as due to known physics we are not able to confirm statement (6) by a direct computation.

# 4 The statements $\Psi_n$ (n = 3, 4, 5, ...), which seem to be true for every $n \in \{3, ..., 9\}$

Let f(3) = 4, and let f(n + 1) = f(n)! for every integer  $n \ge 3$ . For an integer  $n \ge 3$ , let  $\mathcal{U}_n$  denote the following system of equations:

$$\begin{cases} \forall i \in \{1, \dots, n-1\} \setminus \{2\} \ x_i! = x_{i+1} \\ x_1 \cdot x_2 = x_3 \\ x_2 \cdot x_2 = x_3 \end{cases}$$

The diagram in Figure 1 illustrates the construction of the system  $\mathcal{U}_n$ .



**Fig. 1** Construction of the system  $\mathcal{U}_n$ 

**Lemma 6.** For every integer  $n \ge 3$ , the system  $\mathcal{U}_n$  has exactly two solutions in positive integers, namely  $(1, \ldots, 1)$  and  $(2, 2, f(3), \ldots, f(n))$ .

Let

$$B_n = \left\{ x_i! = x_{i+1} : 1 \le i \le n-1 \right\} \cup \left\{ x_i \cdot x_j = x_{j+1} : 1 \le i \le j \le n-1 \right\}$$

For an integer  $n \ge 3$ , let  $\Psi_n$  denote the following statement: *if a system of equations*  $S \subseteq B_n$  *has only finitely many solutions in positive integers*  $x_1, \ldots, x_n$ , *then each such solution*  $(x_1, \ldots, x_n)$  *satisfies*  $x_1, \ldots, x_n \le f(n)$ . The statement  $\Psi_n$  says that for subsystems of  $B_n$  with a finite number of solutions, the largest known solution is indeed the largest possible. The author's guess is that the statements  $\Psi_3, \ldots, \Psi_9$  are true.

**Theorem 5.** Every statement  $\Psi_n$  is true with an unknown integer bound that depends on n.

*Proof.* For every positive integer 
$$n$$
, the system  $B_n$  has a finite number of subsystems.  $\Box$ 

**Theorem 6.** For every statement  $\Psi_n$ , the bound f(n) cannot be decreased.

*Proof.* It follows from Lemma 6 because  $\mathcal{U}_n \subseteq B_n$ .

#### 5 The statement $\Psi_9$ solves Open Problem 2

Let  $\mathcal{A}$  denote the following system of equations:

Lemma 2 and the diagram in Figure 2 explain the construction of the system  $\mathcal{A}$ .

$$x_{1} \xrightarrow{\text{squaring}} x_{2} \xrightarrow{+1} x_{5}$$
or  $x_{2} = x_{5} = 1$ 

$$x_{3} \xrightarrow{x_{5} = x_{6}} x_{6} \xrightarrow{x_{5} \cdot x_{7} = x_{8}} \xrightarrow{+1} x_{8}$$

$$x_{3} \xrightarrow{+1} \text{or } x_{3} = x_{8} = 1$$

$$x_{4} \xrightarrow{x_{4} \cdot x_{8} = x_{9}} x_{9}$$

**Fig. 2** Construction of the system  $\mathcal{A}$ 

**Lemma 7.** For every integer  $x_1 \ge 2$ , the system  $\mathcal{A}$  is solvable in positive integers  $x_2, \ldots, x_9$  if and only if  $x_1^2 + 1$  is prime. In this case, the integers  $x_2, \ldots, x_9$  are uniquely determined by the following equalities:

$$\begin{aligned} x_2 &= x_1^2 \\ x_3 &= (x_1^2)! \\ x_4 &= ((x_1^2)!)! \\ x_5 &= x_1^2 + 1 \\ x_6 &= (x_1^2 + 1)! \\ x_7 &= \frac{(x_1^2)! + 1}{x_1^2 + 1} \\ x_8 &= (x_1^2)! + 1 \\ x_9 &= ((x_1^2)! + 1)! \end{aligned}$$

*Proof.* By Lemma 2, for every integer  $x_1 \ge 2$ , the system  $\mathcal{A}$  is solvable in positive integers  $x_2, \ldots, x_9$  if and only if  $x_1^2 + 1$  divides  $(x_1^2)! + 1$ . Hence, the claim of Lemma 7 follows from Lemma 3.

**Lemma 8.** There are only finitely many tuples  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$  which solve the system  $\mathcal{A}$  and satisfy  $x_1 = 1$ .

*Proof.* If a tuple  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$  solves the system  $\mathcal{A}$  and  $x_1 = 1$ , then  $x_1, \ldots, x_9 \leq 2$ . Indeed,  $x_1 = 1$  implies that  $x_2 = x_1^2 = 1$ . Hence, for example,  $x_3 = x_2! = 1$ . Therefore,  $x_8 = x_3 + 1 = 2$  or  $x_8 = 1$ . Consequently,  $x_9 = x_8! \leq 2$ .

**Theorem 7.** The statement  $\Psi_9$  proves the following implication: if there exists an integer  $x_1 \ge 2$  such that  $x_1^2 + 1$  is prime and greater than f(7), then the set  $\mathcal{P}_{n^2+1}$  is infinite.

*Proof.* Suppose that the antecedent holds. By Lemma 7, there exists a unique tuple  $(x_2, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^8$  such that the tuple  $(x_1, x_2, \ldots, x_9)$  solves the system  $\mathcal{A}$ . Since  $x_1^2 + 1 > f(7)$ , we obtain that  $x_1^2 \ge f(7)$ . Hence,  $(x_1^2)! \ge f(7)! = f(8)$ . Consequently,

$$x_9 = ((x_1^2)! + 1)! \ge (f(8) + 1)! > f(8)! = f(9)$$

Since  $\mathcal{A} \subseteq B_9$ , the statement  $\Psi_9$  and the inequality  $x_9 > f(9)$  imply that the system  $\mathcal{A}$  has infinitely many solutions  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$ . According to Lemmas 7 and 8 the set  $\mathcal{P}_{n^2+1}$  is infinite.  $\Box$ 

Let  $\mathcal{F}$  denote the set of all non-negative integers k whose number of digits belongs to  $\mathcal{P}_{n^2+1}$ .

**Lemma 9.**  $\operatorname{card}(\mathcal{F}) \ge 9 \cdot 10^9 \cdot 4^{747}$ .

Proof. The following PARI/GP ([6]) command

## (12:26) gp > isprime(1+9\*4^747,{flag=2}) %1 = 1

is shown together with its output. This command performs the APRCL primality test, the best deterministic primality test algorithm ([9, p. 226]). It rigorously shows that the number  $(3 \cdot 2^{747})^2 + 1$  is prime. Since  $9 \cdot 10^9 \cdot 4^{747}$  non-negative integers have  $1 + 9 \cdot 4^{747}$  digits, the desired inequality holds.

**Theorem 8.** The statement  $\Psi_9$  implies that  $X = \mathcal{F}$  satisfies conditions (1)-(5).

*Proof.* Suppose that the antecedent holds. Since the set  $\mathcal{P}_{n^2+1}$  is conjecturally infinite, Lemma 9 implies condition (1). Conditions (3) and (5) hold trivially. By Theorem 7,  $\underbrace{9...9}_{f(7) \text{ digits}}$  is a threshold number

of X. Thus condition (4) holds. By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of  $\mathcal{P}_{n^2+1}$  is greater than  $f(7) = ((24!)!)! = \beta$ . Thus condition (2) holds.

#### 6 Open Problems 3 and 4

**Definition 4.** Let (1 $\diamond$ ) denote the following condition: card(X) is greater than a huge positive integer and it is conjectured that  $X = \mathbb{N}$ .

**Definition 5.** Let (2 $\diamond$ ) denote the following condition: we do not know any algorithm deciding the equality  $X = \mathbb{N}$ .

The following two open problems cannot be stated formally as they refer to the current mathematical knowledge.

**Open Problem 3.** Simply define a set  $X \subseteq \mathbb{N}$  that satisfies conditions  $(1\diamond)-(2\diamond)$ , (2)-(3),  $(4\bullet)$ , and (5).

Open Problem 3 claims more than Open Problem 1 as condition (1) implies condition (1).

**Open Problem 4.** Simply define a set  $X \subseteq \mathbb{N}$  that satisfies conditions  $(1\diamond) - (2\diamond)$  and (2) - (5).

Open Problem 4 claims more than Open Problem 2 as condition (1) implies condition (1).

**Theorem 9.** Open Problem 4 claims more than Open Problem 3.

*Proof.* By Theorem 1, condition (4) implies condition (4•).

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#### 7 A partial solution to Open Problem 4

Let  $\mathcal V$  denote the set of all positive multiples of elements of the set

$${n \in \{\beta + 1, \beta + 2, \beta + 3, \ldots\} : 2^{2^n} + 1 \text{ is composite}\}}$$

**Theorem 10.** The set  $X = \{0, \ldots, \beta\} \cup \mathcal{V}$  satisfies conditions  $(1\diamond) - (2\diamond)$  and (2) - (4).

*Proof.* The inequality  $\operatorname{card}(X) > \beta$  holds trivially. Most mathematicians believe that  $2^{2^n} + 1$  is composite for every integer  $n \ge 5$ , see [2, p. 23]. These two facts imply conditions (1 $\diamond$ ) and (2 $\diamond$ ). Condition (3) holds trivially. Since the set  $\mathcal{V}$  is empty or infinite, the integer  $\beta$  is a threshold number of X. Thus condition (4) holds. The question of finiteness of the set { $n \in \mathbb{N} : 2^{2^n} + 1$  is composite} remains open, see [3, p. 159]. By this and Lemma 1, the question of emptiness of the set

$${n \in \{\beta + 1, \beta + 2, \beta + 3, \ldots\} : 2^{2^n} + 1 \text{ is composite}}$$

remains open. Therefore, the question of finiteness of the set  $\mathcal{V}$  remains open. Consequently, the question of finiteness of the set  $\mathcal{X}$  remains open and condition (2) holds.

#### 8 Open Problems 5 and 6

**Definition 6.** Let  $(1^*)$  denote the following condition: card(X) is greater than a huge positive integer and it is conjectured that X is finite.

The following two open problems cannot be stated formally as they refer to the current mathematical knowledge.

**Open Problem 5.** Simply define a set  $X \subseteq \mathbb{N}$  that satisfies conditions (1\*), (2)-(3), (4•), and (5).

**Open Problem 6.** Simply define a set  $X \subseteq \mathbb{N}$  that satisfies conditions (1\*) and (2)–(5).

**Theorem 11.** Open Problem 6 claims more than Open Problem 5.

*Proof.* By Theorem 1, condition (4) implies condition  $(4\bullet)$ .

#### 9 A partial solution to Open Problem 6

A weak form of Szpiro's conjecture implies that there are only finitely many solutions to the equation  $x! + 1 = y^2$ , see [5].

**Lemma 10.** ([8, p. 297]). It is conjectured that x! + 1 is a square only for  $x \in \{4, 5, 7\}$ .

Let W denote the set of all integers x greater than  $\beta$  such that x! + 1 is a square.

Theorem 12. The set

$$\mathcal{X} = \{0, \dots, \beta\} \cup \{k \cdot x : (k \in \mathbb{N} \setminus \{0\}) \land (x \in \mathcal{W})\}$$

satisfies conditions (1\*) and (2)-(4).

*Proof.* Condition (1\*) holds as  $card(X) > \beta$  and the set W is conjecturally empty by Lemma 10. Condition (3) holds trivially. We do not know any algorithm that decides the emptiness of W and the set

$$\mathcal{Y} = \{k \cdot x : (k \in \mathbb{N} \setminus \{0\}) \land (x \in \mathcal{W})\}$$

is empty or infinite. Thus condition (2) holds. Since the set  $\mathcal{Y}$  is empty or infinite, the integer  $\beta$  is a threshold number of  $\mathcal{X}$ . Thus condition (4) holds.

#### **10** The statement $\Psi_6$ solves Open Problem 6

Let *C* denote the following system of equations:

$$\begin{array}{rcl}
x_1! &=& x_2 \\
x_2! &=& x_3 \\
x_5! &=& x_6 \\
x_4 \cdot x_4 &=& x_5 \\
x_3 \cdot x_5 &=& x_6
\end{array}$$

Lemma 2 and the diagram in Figure 3 explain the construction of the system C.



Fig. 3 Construction of the system C

**Lemma 11.** For every  $x_1, x_4 \in \mathbb{N} \setminus \{0, 1\}$ , the system *C* is solvable in positive integers  $x_2, x_3, x_5, x_6$  if and only if  $x_1! + 1 = x_4^2$ . In this case, the integers  $x_2, x_3, x_5, x_6$  are uniquely determined by the following equalities:

$$x_{2} = x_{1}!$$
  

$$x_{3} = (x_{1}!)!$$
  

$$x_{5} = x_{1}! + 1$$
  

$$x_{6} = (x_{1}! + 1)!$$

*Proof.* It follows from Lemma 2.

**Theorem 13.** If the equation  $x_1! + 1 = x_4^2$  has only finitely many solutions in positive integers, then the statement  $\Psi_6$  guarantees that each such solution  $(x_1, x_4)$  belongs to the set  $\{(4, 5), (5, 11), (7, 71)\}$ .

*Proof.* Suppose that the antecedent holds. Let positive integers  $x_1$  and  $x_4$  satisfy  $x_1! + 1 = x_4^2$ . Then,  $x_1, x_4 \in \mathbb{N} \setminus \{0, 1\}$ . By Lemma 11, the system *C* is solvable in positive integers  $x_2, x_3, x_5, x_6$ . Since  $C \subseteq B_6$ , the statement  $\Psi_6$  implies that  $x_6 = (x_1! + 1)! \leq f(6) = f(5)!$ . Hence,  $x_1! + 1 \leq f(5) = f(4)!$ . Consequently,  $x_1 < f(4) = 24$ . If  $x_1 \in \{1, \dots, 23\}$ , then  $x_1! + 1$  is a square only for  $x_1 \in \{4, 5, 7\}$ .

**Theorem 14.** Let X denote the set of all non-negative integers n which have (((k!)!)!)! digits for some  $k \in \{m \in \mathbb{N} : m! + 1 \text{ is a square}\}$ . We claim that the statement  $\Psi_6$  implies that X satisfies conditions (1\*) and (2)-(5).

*Proof.* Let d = (((7!)!)!)!. Since  $7! + 1 = 71^2$ , we obtain that  $\{10^{d-1}, \dots, \underbrace{9 \dots 9}_{d \text{ digits}}\} \subseteq X$ . Hence,  $card(X) \ge d$ 

 $9 \cdot 10^{d-1}$ . By this and Lemma 10, condition (1\*) holds. Conditions (2)-(3) and (5) hold trivially. By Theorem 13, the statement  $\Psi_6$  implies that  $\underbrace{9...9}_{d \text{ digits}}$  is a threshold number of X. Thus condition (4)

holds.

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Sławomir Kurpaska Hugo Kołłątaj University Balicka 116B, 30-149 Kraków, Poland E-mail: rtkurpas@cyf-kr.edu.pl

Apoloniusz Tyszka Hugo Kołłątaj University Balicka 116B, 30-149 Kraków, Poland E-mail: rttyszka@cyf-kr.edu.pl

### On *ZFC*-formulae $\varphi(x)$ for which we know a non-negative integer *n* such that $\{x \in \mathbb{N} : \varphi(x)\} \subseteq \{x \in \mathbb{N} : x \leq n-1\}$ if the set $\{x \in \mathbb{N} : \varphi(x)\}$ is finite

Apoloniusz Tyszka

#### Abstract

Let  $\Gamma(k)$  denote (k-1)!, and let  $\Gamma_n(k)$  denote (k-1)!, where  $n \in \{3, ..., 16\}$  and  $k \in \{2\} \cup [2^{2^{n-3}} + 1, \infty) \cap \mathbb{N}$ . For an integer  $n \in \{3, ..., 16\}$ , let  $\Sigma_n$  denote the following statement: if a system of equations  $S \subseteq \{\Gamma_n(x_i) = x_k : i, k \in \{1, ..., n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, ..., n\}\}$  with  $\Gamma$  instead of  $\Gamma_n$  has only finitely many solutions in positive integers  $x_1, ..., x_n$ , then every tuple  $(x_1, ..., x_n) \in (\mathbb{N} \setminus \{0\})^n$  that solves the original system S satisfies  $x_1, ..., x_n \notin 2^{2^{n-2}}$ . Our hypothesis claims that the statements  $\Sigma_3, ..., \Sigma_{16}$  are true. The statement  $\Sigma_6$  proves the following implication: if the equation x(x + 1) = y! has only finitely many solutions in positive integers x and y, then each such solution (x, y) belongs to the set  $\{(1, 2), (2, 3)\}$ . The statement  $\Sigma_6$  proves the following implication: if the equation  $x! + 1 = y^2$  has only finitely many solutions in positive integers x and y, then each such solution (x, y) belongs to the set  $\{(4, 5), (5, 11), (7, 71)\}$ . The statement  $\Sigma_9$  implies the infinitude of primes of the form  $n^2 + 1$ . The statement  $\Sigma_9$  implies that any prime of the form n! + 1 with  $n \ge 2^{2^{9-3}}$  proves the infinitude of primes of the form n! + 1 implies the infinitude of twin primes. The statement  $\Sigma_{16}$  implies the infinitude of Sophie Germain primes.

**Key words and phrases:** Brocard's problem, Brocard-Ramanujan equation  $x! + 1 = y^2$ , composite Fermat numbers, decidability in the limit, Erdös' equation x(x + 1) = y!, finiteness of a set, infiniteness of a set, prime numbers of the form  $n^2 + 1$ , prime numbers of the form n! + 1, single query to an oracle for the halting problem, Sophie Germain primes, twin primes.

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#### 1 Introduction and basic lemmas

The phrase "we know a non-negative integer n" in the title means that we know an algorithm which returns *n*. The title of the article cannot be formalised in *ZFC* because the phrase "we know a non-negative integer n" refers to currently known non-negative integers *n* with some property. A formally stated title may look like this: On *ZFC*-formulae  $\varphi(x)$  for which there exists a non-negative integer *n* such that *ZFC* proves that

 $\operatorname{card}(\{x \in \mathbb{N} \colon \varphi(x)\}) < \infty \Longrightarrow \{x \in \mathbb{N} \colon \varphi(x)\} \subseteq \{x \in \mathbb{N} \colon x \le n-1\}$ 

Unfortunately, this formulation admits formulae  $\varphi(x)$  without any known non-negative integer *n* such that *ZFC* proves the above implication.

**Lemma 1.** For every non-negative integer n, card({ $x \in \mathbb{N}$ :  $x \leq n - 1$ }) = n.

**Corollary 1.** *The title altered to* "On *ZFC*-formulae  $\varphi(x)$  for which we know a non-negative integer *n* such that card( $\{x \in \mathbb{N} : \varphi(x)\}$ )  $\leq n$  if the set  $\{x \in \mathbb{N} : \varphi(x)\}$  is finite" *involves a weaker assumption on*  $\varphi(x)$ .

**Lemma 2.** For every positive integers x and y,  $x! \cdot y = y!$  if and only if

$$(x + 1 = y) \lor (x = y = 1)$$

Let  $\Gamma(k)$  denote (k - 1)!.

**Lemma 3.** For every positive integers x and y,  $x \cdot \Gamma(x) = \Gamma(y)$  if and only if

$$(x + 1 = y) \lor (x = y = 1)$$

**Lemma 4.** For every non-negative integers b and c, b + 1 = c if and only if  $2^{2^b} \cdot 2^{2^b} = 2^{2^c}$ .

**Lemma 5.** (Wilson's theorem, [8, p. 89]). For every positive integer x, x divides (x - 1)! + 1 if and only if x = 1 or x is prime.

#### **2** Subsets of $\mathbb{N}$ and their threshold numbers

We say that a non-negative integer *m* is a threshold number of a set  $X \subseteq \mathbb{N}$ , if X is infinite if and only if X contains an element greater than *m*, cf. [24] and [25]. If a set  $X \subseteq \mathbb{N}$  is empty or infinite, then any non-negative integer *m* is a threshold number of X. If a set  $X \subseteq \mathbb{N}$  is non-empty and finite, then the all threshold numbers of X form the set {max(X), max(X) + 1, max(X) + 2, ...}.

It is conjectured that the set of prime numbers of the form  $n^2 + 1$  is infinite, see [14, pp. 37–38]. It is conjectured that the set of prime numbers of the form n! + 1 is infinite, see [3, p. 443]. A twin prime is a prime number that differs from another prime number by 2. The twin prime conjecture states that the set of twin primes is infinite, see [14, p. 39]. It is conjectured that the set of composite numbers of the form  $2^{2^n} + 1$  is infinite, see [10, p. 23] and [11, pp. 158–159]. A prime p is said to be a Sophie Germain prime if both p and 2p + 1 are prime, see [22]. It is conjectured that the set of Sophie Germain primes is infinite, see [17, p. 330]. For each of these sets, we do not know any threshold number.

The following statement:

for every non-negative integer *n* there exist

prime numbers 
$$p$$
 and  $q$  such that  $p + 2 = q$  and  $p \in \left[10^n, 10^{n+1}\right]$  (1)

is a  $\Pi_1$  statement which strengthens the twin prime conjecture, see [4, p. 43]. C. H. Bennett claims that most mathematical conjectures can be settled indirectly by proving stronger  $\Pi_1$  statements, see [1]. Statement (1) is equivalent to the non-halting of a Turing machine. If a set  $X \subseteq \mathbb{N}$  is computable and we know a threshold number of X, then the infinity of X is equivalent to the halting of a Turing machine.

The height of a rational number  $\frac{p}{q}$  is denoted by  $H\left(\frac{p}{q}\right)$  and equals  $\max(|p|, |q|)$  provided  $\frac{p}{q}$  is written in lowest terms. The height of a rational tuple  $(x_1, \ldots, x_n)$  is denoted by  $H(x_1, \ldots, x_n)$  and equals  $\max(H(x_1), \ldots, H(x_n))$ .

**Lemma 6.** The equation  $x^5 - x = y^2 - y$  has only finitely many rational solutions, see [13, p. 212]. The known rational solutions are  $(x, y) = (-1, 0), (-1, 1), (0, 0), (0, 1), (1, 0), (1, 1), (2, -5), (2, 6), (3, -15), (3, 16), (30, -4929), (30, 4930), <math>(\frac{1}{4}, \frac{15}{32}), (\frac{1}{4}, \frac{17}{32}), (-\frac{15}{16}, -\frac{185}{1024}), (-\frac{15}{16}, \frac{1209}{1024})$ , and the existence of other solutions is an open question, see [18, pp. 223–224].

**Corollary 2.** The set  $\mathcal{T} = \{n \in \mathbb{N} : \text{the equation } x^5 - x = y^2 - y \text{ has a rational solution of height } n\}$  is finite. We know an algorithm which for every  $n \in \mathbb{N}$  decides whether or not  $n \in \mathcal{T}$ . We do not know any algorithm which returns a threshold number of  $\mathcal{T}$ .

Let  $\mathcal{L}$  denote the following system of equations:

$$\begin{cases} x^2 + y^2 &= s^2 \\ x^2 + z^2 &= t^2 \\ y^2 + z^2 &= u^2 \\ x^2 + y^2 + z^2 &= v^2 \end{cases}$$

Let

$$\mathcal{F} = \left\{ n \in \mathbb{N} \setminus \{0\} : \left( \text{the system } \mathcal{L} \text{ has no solutions in } \{1, \dots, n\}^7 \right) \land$$

$$\left( \text{the system } \mathcal{L} \text{ has a solution in } \{1, \dots, n+1\}^7 \right) \right\}$$

A perfect cuboid is a cuboid having integer side lengths, integer face diagonals, and an integer space diagonal.

Lemma 7. ([21]). No perfect cuboids are known.

**Corollary 3.** We know an algorithm which for every  $n \in \mathbb{N}$  decides whether or not  $n \in \mathcal{F}$ . ZFC proves that  $\operatorname{card}(\mathcal{F}) \in \{0, 1\}$ . We do not know any algorithm which returns  $\operatorname{card}(\mathcal{F})$ . We do not know any algorithm which returns a threshold number of  $\mathcal{F}$ .

Let

We do not know whether or not the set  $\mathcal{H}$  is finite.

**Proposition 1.** The number  $9^{99^{9^{7^{*}}}}$  is a threshold number of  $\mathcal{H}$ . We know an algorithm which decides the equality  $\mathcal{H} = \mathbb{N}$ . If  $\mathcal{H} \neq \mathbb{N}$ , then the set  $\mathcal{H}$  consists of all integers from 0 to a non-negative integer which can be computed by a known algorithm. We know an algorithm which for every  $n \in \mathbb{N}$  decides whether or not  $n \in \mathcal{H}$ .

Let

$$\mathcal{K} = \begin{cases} \{n\}, \text{ if } (n \in \mathbb{N}) \land \left(2^{\aleph_0} = \aleph_{n+1}\right) \\ \{0\}, \text{ if } 2^{\aleph_0} \ge \aleph_{\omega} \end{cases}$$

**Proposition 2.** *ZFC proves that*  $card(\mathcal{K}) = 1$ . *If ZFC is consistent, then for every*  $n \in \mathbb{N}$  *the sentences* "*n* is a threshold number of  $\mathcal{K}$ " *and* "*n* is not a threshold number of  $\mathcal{K}$ " *are not provable in ZFC.* 

*Proof.* It suffices to observe that  $2^{\aleph_0}$  can attain every value from the set  $\{\aleph_1, \aleph_2, \aleph_3, \ldots\}$ , see [7] and [9, p. 232].

# **3** A Diophantine equation whose non-solvability expresses the consistency of *ZFC*

Gödel's second incompleteness theorem and the Davis-Putnam-Robinson-Matiyasevich theorem imply the following theorem.

**Theorem 1.** ([5, p. 35]). There exists a polynomial  $D(x_1, ..., x_m)$  with integer coefficients such that if ZFC is arithmetically consistent, then the sentences "The equation  $D(x_1, ..., x_m) = 0$  is solvable in non-negative integers" and "The equation  $D(x_1, ..., x_m) = 0$  is not solvable in non-negative integers" are not provable in ZFC.

Let  $\mathcal{Y}$  denote the set of all non-negative integers k such that the equation  $D(x_1, \ldots, x_m) = 0$  has no solutions in  $\{0, \ldots, k\}^m$ . Since the set  $\{0, \ldots, k\}^m$  is finite, we know an algorithm which for every  $n \in \mathbb{N}$  decides whether or not  $n \in \mathcal{Y}$ . Theorem 1 implies the next theorem.

**Theorem 2.** For every  $n \in \mathbb{N}$ , ZFC proves that  $n \in \mathcal{Y}$ . If ZFC is arithmetically consistent, then the sentences " $\mathcal{Y}$  is finite" and " $\mathcal{Y}$  is infinite" are not provable in ZFC. If ZFC is arithmetically consistent, then for every  $n \in \mathbb{N}$  the sentences "n is a threshold number of  $\mathcal{Y}$ " and "n is not a threshold number of  $\mathcal{Y}$ " are not provable in ZFC.

Let  $\mathcal{E}$  denote the set of all non-negative integers k such that the equation  $D(x_1, \ldots, x_m) = 0$  has a solution in  $\{0, \ldots, k\}^m$ . Since the set  $\{0, \ldots, k\}^m$  is finite, we know an algorithm which for every  $n \in \mathbb{N}$  decides whether or not  $n \in \mathcal{E}$ . Theorem 1 implies the next theorem.

**Theorem 3.** The set  $\mathcal{E}$  is empty or infinite. In both cases, every non-negative integer n is a threshold number of  $\mathcal{E}$ . If ZFC is arithmetically consistent, then the sentences " $\mathcal{E}$  is empty", " $\mathcal{E}$  is not empty", " $\mathcal{E}$  is finite", and " $\mathcal{E}$  is infinite" are not provable in ZFC.

Let

 $\mathcal{V} = \{n \in \mathbb{N} : (\text{the polynomial } D(x_1, \dots, x_m) \text{ has no solutions in } \{0, \dots, n\}^m) \land$ 

(the polynomial  $D(x_1, \ldots, x_m)$  has a solution in  $\{0, \ldots, n+1\}^m$ )

Since the sets  $\{0, ..., n\}^m$  and  $\{0, ..., n + 1\}^m$  are finite, we know an algorithm which for every  $n \in \mathbb{N}$  decides whether or not  $n \in \mathcal{V}$ . Theorem 1 implies the next theorem.

**Theorem 4.** *ZFC* proves that  $card(\mathcal{V}) \in \{0, 1\}$ . For every  $n \in \mathbb{N}$ , *ZFC* proves that  $n \notin \mathcal{V}$ . *ZFC* does not prove the emptiness of  $\mathcal{V}$ , if *ZFC* is arithmetically consistent. For every  $n \in \mathbb{N}$ , the sentence "n is a threshold number of  $\mathcal{V}$ " is not provable in *ZFC*, if *ZFC* is arithmetically consistent.

#### **4** Hypothetical statements $\Psi_3, \ldots, \Psi_{16}$

For an integer  $n \ge 3$ , let  $\mathcal{U}_n$  denote the following system of equations:

$$\begin{cases} \forall i \in \{1, \dots, n-1\} \setminus \{2\} \ x_i! = x_{i+1} \\ x_1 \cdot x_2 = x_3 \\ x_2 \cdot x_2 = x_3 \end{cases}$$

The diagram in Figure 1 illustrates the construction of the system  $\mathcal{U}_n$ .



**Fig. 1** Construction of the system  $\mathcal{U}_n$ 

Let g(3) = 4, and let g(n + 1) = g(n)! for every integer  $n \ge 3$ .

**Lemma 8.** For every integer  $n \ge 3$ , the system  $\mathcal{U}_n$  has exactly two solutions in positive integers, namely  $(1, \ldots, 1)$  and  $(2, 2, g(3), \ldots, g(n))$ .

Let

 $B_n = \left\{ x_i ! = x_k : (i, k \in \{1, \dots, n\}) \land (i \neq k) \right\} \cup \left\{ x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\} \right\}$ 

For an integer  $n \ge 3$ , let  $\Psi_n$  denote the following statement: *if a system of equations*  $S \subseteq B_n$  *has only finitely many solutions in positive integers*  $x_1, \ldots, x_n$ , *then each such solution*  $(x_1, \ldots, x_n)$  *satisfies*  $x_1, \ldots, x_n \le g(n)$ . The statement  $\Psi_n$  says that for subsystems of  $B_n$  the largest known solution is indeed the largest possible.

**Hypothesis 1.** The statements  $\Psi_3, \ldots, \Psi_{16}$  are true.

**Proposition 3.** Every statement  $\Psi_n$  is true with an unknown integer bound that depends on n. *Proof.* For every positive integer n, the system  $B_n$  has a finite number of subsystems.  $\Box$  **Proposition 4.** For every statement  $\Psi_n$ , the bound g(n) cannot be decreased. *Proof.* It follows from Lemma 8 because  $\mathcal{U}_n \subseteq B_n$ .  $\Box$ 

#### **5** The Brocard-Ramanujan equation $x! + 1 = y^2$

Let  $\mathcal{A}$  denote the following system of equations:

$$\begin{array}{rcl}
x_1! &=& x_2 \\
x_2! &=& x_3 \\
x_5! &=& x_6 \\
x_4 \cdot x_4 &=& x_5 \\
x_3 \cdot x_5 &=& x_6
\end{array}$$

Lemma 2 and the diagram in Figure 2 explain the construction of the system  $\mathcal{A}$ .



**Fig. 2** Construction of the system  $\mathcal{A}$ 

**Lemma 9.** For every  $x_1, x_4 \in \mathbb{N} \setminus \{0, 1\}$ , the system  $\mathcal{A}$  is solvable in positive integers  $x_2, x_3, x_5, x_6$  if and only if  $x_1! + 1 = x_4^2$ . In this case, the integers  $x_2, x_3, x_5, x_6$  are uniquely determined by the following equalities:

$$x_{2} = x_{1}!$$
  

$$x_{3} = (x_{1}!)!$$
  

$$x_{5} = x_{1}! + 1$$
  

$$x_{6} = (x_{1}! + 1)!$$

*Proof.* It follows from Lemma 2.

It is conjectured that x! + 1 is a perfect square only for  $x \in \{4, 5, 7\}$ , see [20, p. 297]. A weak form of Szpiro's conjecture implies that there are only finitely many solutions to the equation  $x! + 1 = y^2$ , see [15].

**Theorem 5.** If the equation  $x_1! + 1 = x_4^2$  has only finitely many solutions in positive integers, then the statement  $\Psi_6$  guarantees that each such solution  $(x_1, x_4)$  belongs to the set {(4, 5), (5, 11), (7, 71)}.

*Proof.* Suppose that the antecedent holds. Let positive integers  $x_1$  and  $x_4$  satisfy  $x_1! + 1 = x_4^2$ . Then,  $x_1, x_4 \in \mathbb{N} \setminus \{0, 1\}$ . By Lemma 9, the system  $\mathcal{A}$  is solvable in positive integers  $x_2, x_3, x_5, x_6$ . Since  $\mathcal{A} \subseteq B_6$ , the statement  $\Psi_6$  implies that  $x_6 = (x_1! + 1)! \leq g(6) = g(5)!$ . Hence,  $x_1! + 1 \leq g(5) = g(4)!$ . Consequently,  $x_1 < g(4) = 24$ . If  $x_1 \in \{1, \dots, 23\}$ , then  $x_1! + 1$  is a perfect square only for  $x_1 \in \{4, 5, 7\}$ .

#### 6 Are there infinitely many prime numbers of the form $n^2 + 1$ ?

Edmund Landau's conjecture states that there are infinitely many primes of the form  $n^2 + 1$ , see [14, pp. 37–38]. Let  $\mathcal{B}$  denote the following system of equations:

$$\begin{cases} x_2! = x_3 \\ x_3! = x_4 \\ x_5! = x_6 \\ x_8! = x_9 \\ x_1 \cdot x_1 = x_2 \\ x_3 \cdot x_5 = x_6 \\ x_4 \cdot x_8 = x_9 \\ x_5 \cdot x_7 = x_8 \end{cases}$$

Lemma 2 and the diagram in Figure 3 explain the construction of the system  $\mathcal{B}$ .



**Fig. 3** Construction of the system  $\mathcal{B}$ 

**Lemma 10.** For every integer  $x_1 \ge 2$ , the system  $\mathcal{B}$  is solvable in positive integers  $x_2, \ldots, x_9$  if and only if  $x_1^2 + 1$  is prime. In this case, the integers  $x_2, \ldots, x_9$  are uniquely determined by the following equalities:

$$\begin{array}{rcl} x_2 &=& x_1^2 \\ x_3 &=& (x_1^2)! \\ x_4 &=& ((x_1^2)!)! \\ x_5 &=& x_1^2 + 1 \\ x_6 &=& (x_1^2 + 1)! \\ x_7 &=& \frac{(x_1^2)! + 1}{x_1^2 + 1} \\ x_8 &=& (x_1^2)! + 1 \\ x_9 &=& ((x_1^2)! + 1)! \end{array}$$

*Proof.* By Lemma 2, for every integer  $x_1 \ge 2$ , the system  $\mathcal{B}$  is solvable in positive integers  $x_2, \ldots, x_9$  if and only if  $x_1^2 + 1$  divides  $(x_1^2)! + 1$ . Hence, the claim of Lemma 10 follows from Lemma 5.

**Lemma 11.** There are only finitely many tuples  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$  which solve the system  $\mathcal{B}$  and satisfy  $x_1 = 1$ .

*Proof.* If a tuple  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$  solves the system  $\mathcal{B}$  and  $x_1 = 1$ , then  $x_1, \ldots, x_9 \leq 2$ . Indeed,  $x_1 = 1$  implies that  $x_2 = x_1^2 = 1$ . Hence, for example,  $x_3 = x_2! = 1$ . Therefore,  $x_8 = x_3 + 1 = 2$  or  $x_8 = 1$ . Consequently,  $x_9 = x_8! \leq 2$ .

**Theorem 6.** The statement  $\Psi_9$  proves the following implication: if there exists an integer  $x_1 \ge 2$  such that  $x_1^2 + 1$  is prime and greater than g(7), then there are infinitely many primes of the form  $n^2 + 1$ .

*Proof.* Suppose that the antecedent holds. By Lemma 10, there exists a unique tuple  $(x_2, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^8$  such that the tuple  $(x_1, x_2, \ldots, x_9)$  solves the system  $\mathcal{B}$ . Since  $x_1^2 + 1 > g(7)$ , we obtain that  $x_1^2 \ge g(7)$ . Hence,  $(x_1^2)! \ge g(7)! = g(8)$ . Consequently,

$$x_9 = ((x_1^2)! + 1)! \ge (g(8) + 1)! > g(8)! = g(9)$$

Since  $\mathcal{B} \subseteq B_9$ , the statement  $\Psi_9$  and the inequality  $x_9 > g(9)$  imply that the system  $\mathcal{B}$  has infinitely many solutions  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$ . According to Lemmas 10 and 11, there are infinitely many primes of the form  $n^2 + 1$ .

**Corollary 4.** Let  $X_9$  denote the set of primes of the form  $n^2 + 1$ . The statement  $\Psi_9$  implies that we know an algorithm such that it returns a threshold number of  $X_9$ , and this number equals  $\max(X_9)$ , if  $X_9$  is finite. Assuming the statement  $\Psi_9$ , a single query to an oracle for the halting problem decides the infinity of  $X_9$ . Assuming the statement  $\Psi_9$ , the infinity of  $X_9$  is decidable in the limit.

*Proof.* We consider an algorithm which computes  $\max(X_9 \cap [1, g(7)])$ .

#### 7 Are there infinitely many prime numbers of the form n! + 1?

It is conjectured that there are infinitely many primes of the form n! + 1, see [3, p. 443].

**Theorem 7.** (cf. Theorem 11). The statement  $\Psi_9$  proves the following implication: if there exists an integer  $x_1 \ge g(6)$  such that  $x_1! + 1$  is prime, then there are infinitely many primes of the form n! + 1.

*Proof.* We leave the analogous proof to the reader.

#### 8 The twin prime conjecture

A twin prime is a prime number that differs from another prime number by 2. The twin prime conjecture states that there are infinitely many twin primes, see [14, p. 39]. Let C denote the following system of equations:

$$x_{1}! = x_{2}$$

$$x_{2}! = x_{3}$$

$$x_{4}! = x_{5}$$

$$x_{6}! = x_{7}$$

$$x_{7}! = x_{8}$$

$$x_{9}! = x_{10}$$

$$x_{12}! = x_{13}$$

$$x_{15}! = x_{16}$$

$$x_{2} \cdot x_{4} = x_{5}$$

$$x_{5} \cdot x_{6} = x_{7}$$

$$x_{7} \cdot x_{9} = x_{10}$$

$$x_{4} \cdot x_{11} = x_{12}$$

$$x_{3} \cdot x_{12} = x_{13}$$

$$x_{9} \cdot x_{14} = x_{15}$$

$$x_{8} \cdot x_{15} = x_{16}$$

Lemma 2 and the diagram in Figure 4 explain the construction of the system C.



Fig. 4 Construction of the system C

**Lemma 12.** For every  $x_4, x_9 \in \mathbb{N} \setminus \{0, 1, 2\}$ , the system *C* is solvable in positive integers  $x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$  if and only if  $x_4$  and  $x_9$  are prime and  $x_4 + 2 = x_9$ . In this case, the integers  $x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$  are uniquely determined by the following equalities:

$$x_{1} = x_{4} - 1$$

$$x_{2} = (x_{4} - 1)!$$

$$x_{3} = ((x_{4} - 1)!)!$$

$$x_{5} = x_{4}!$$

$$x_{6} = x_{9} - 1$$

$$x_{7} = (x_{9} - 1)!$$

$$x_{8} = ((x_{9} - 1)!)!$$

$$x_{10} = x_{9}!$$

$$x_{11} = \frac{(x_{4} - 1)! + 1}{x_{4}}$$

$$x_{12} = (x_{4} - 1)! + 1$$

$$x_{13} = ((x_{4} - 1)! + 1)!$$

$$x_{14} = \frac{(x_{9} - 1)! + 1}{x_{9}}$$

$$x_{15} = (x_{9} - 1)! + 1$$

*Proof.* By Lemma 2, for every  $x_4, x_9 \in \mathbb{N} \setminus \{0, 1, 2\}$ , the system *C* is solvable in positive integers  $x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$  if and only if

$$(x_4 + 2 = x_9) \land (x_4 | (x_4 - 1)! + 1) \land (x_9 | (x_9 - 1)! + 1)$$

Hence, the claim of Lemma 12 follows from Lemma 5.

**Lemma 13.** There are only finitely many tuples  $(x_1, \ldots, x_{16}) \in (\mathbb{N} \setminus \{0\})^{16}$  which solve the system C and satisfy  $(x_4 \in \{1, 2\}) \lor (x_9 \in \{1, 2\})$ .

*Proof.* If a tuple  $(x_1, \ldots, x_{16}) \in (\mathbb{N} \setminus \{0\})^{16}$  solves the system *C* and  $(x_4 \in \{1, 2\}) \lor (x_9 \in \{1, 2\})$ , then  $x_1, \ldots, x_{16} \leq 7!$ . Indeed, for example, if  $x_4 = 2$  then  $x_6 = x_4 + 1 = 3$ . Hence,  $x_7 = x_6! = 6$ . Therefore,  $x_{15} = x_7 + 1 = 7$ . Consequently,  $x_{16} = x_{15}! = 7!$ .

**Theorem 8.** The statement  $\Psi_{16}$  proves the following implication: if there exists a twin prime greater than g(14), then there are infinitely many twin primes.

*Proof.* Suppose that the antecedent holds. Then, there exist prime numbers  $x_4$  and  $x_9$  such that  $x_9 = x_4 + 2 > g(14)$ . Hence,  $x_4, x_9 \in \mathbb{N} \setminus \{0, 1, 2\}$ . By Lemma 12, there exists a unique tuple  $(x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}) \in (\mathbb{N} \setminus \{0\})^{14}$  such that the tuple  $(x_1, \dots, x_{16})$  solves the system *C*. Since  $x_9 > g(14)$ , we obtain that  $x_9 - 1 \ge g(14)$ . Therefore,  $(x_9 - 1)! \ge g(14)! = g(15)$ . Hence,  $(x_9 - 1)! + 1 > g(15)$ . Consequently,

$$x_{16} = ((x_9 - 1)! + 1)! > g(15)! = g(16)$$

Since  $C \subseteq B_{16}$ , the statement  $\Psi_{16}$  and the inequality  $x_{16} > g(16)$  imply that the system *C* has infinitely many solutions in positive integers  $x_1, \ldots, x_{16}$ . According to Lemmas 12 and 13, there are infinitely many twin primes.

**Corollary 5.** (cf. [6]). Let  $X_{16}$  denote the set of twin primes. The statement  $\Psi_{16}$  implies that we know an algorithm such that it returns a threshold number of  $X_{16}$ , and this number equals  $\max(X_{16})$ , if  $X_{16}$ is finite. Assuming the statement  $\Psi_{16}$ , a single query to an oracle for the halting problem decides the infinity of  $X_{16}$ . Assuming the statement  $\Psi_{16}$ , the infinity of  $X_{16}$  is decidable in the limit.

*Proof.* We consider an algorithm which computes  $\max(X_{16} \cap [1, g(14)])$ .

#### **9** Hypothetical statements $\Delta_5, \ldots, \Delta_{14}$ and their consequences

Let  $\lambda(5) = \Gamma(25)$ , and let  $\lambda(n + 1) = \Gamma(\lambda(n))$  for every integer  $n \ge 5$ . For an integer  $n \ge 5$ , let  $\mathcal{J}_n$  denote the following system of equations:

$$\begin{cases} \forall i \in \{1, \dots, n-1\} \setminus \{3\} \ \Gamma(x_i) &= x_{i+1} \\ x_1 \cdot x_1 &= x_4 \\ x_2 \cdot x_3 &= x_5 \end{cases}$$

Lemma 3 and the diagram in Figure 5 explain the construction of the system  $\mathcal{J}_n$ .



**Fig. 5** Construction of the system  $\mathcal{J}_n$ 

For every integer  $n \ge 5$ , the system  $\mathcal{J}_n$  has exactly two solutions in positive integers, namely  $(1, \ldots, 1)$  and  $(5, 24, 23!, 25, \lambda(5), \ldots, \lambda(n))$ . For an integer  $n \ge 5$ , let  $\Delta_n$  denote the following statement: *if a system of equations*  $S \subseteq \{\Gamma(x_i) = x_k : i, k \in \{1, \ldots, n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \ldots, n\}\}$  has only finitely many solutions in positive integers  $x_1, \ldots, x_n$ , then each such solution  $(x_1, \ldots, x_n)$  satisfies  $x_1, \ldots, x_n \le \lambda(n)$ .

**Hypothesis 2.** The statements  $\Delta_5, \ldots, \Delta_{14}$  are true.

Lemmas 3 and 5 imply that the statements  $\Delta_n$  have similar consequences as the statements  $\Psi_n$ .

**Theorem 9.** The statement  $\Delta_6$  implies that any prime number  $p \ge 25$  proves the infinitude of primes.

Proof. It follows from Lemmas 3 and 5. We leave the details to the reader.

#### **10** Hypothetical statements $\Sigma_3, \ldots, \Sigma_{16}$ and their consequences

Let  $\Gamma_n(k)$  denote (k-1)!, where  $n \in \{3, ..., 16\}$  and  $k \in \{2\} \cup [2^{2^{n-3}} + 1, \infty) \cap \mathbb{N}$ . For an integer  $n \in \{3, ..., 16\}$ , let

$$Q_n = \{\Gamma_n(x_i) = x_k : i, k \in \{1, \dots, n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}\}$$

For an integer  $n \in \{3, ..., 16\}$ , let  $P_n$  denote the following system of equations:

$$\begin{cases} x_1 \cdot x_1 &= x_1 \\ \Gamma_n(x_2) &= x_1 \\ \forall i \in \{2, \dots, n-1\} x_i \cdot x_i &= x_{i+1} \end{cases}$$

**Lemma 14.** For every integer  $n \in \{3, ..., 16\}$ ,  $P_n \subseteq Q_n$  and the system  $P_n$  with  $\Gamma$  instead of  $\Gamma_n$  has exactly one solution in positive integers  $x_1, ..., x_n$ , namely  $(1, 2^{2^0}, 2^{2^1}, 2^{2^2}, ..., 2^{2^{n-2}})$ .

For an integer  $n \in \{3, ..., 16\}$ , let  $\Sigma_n$  denote the following statement: if a system of equations  $S \subseteq Q_n$  with  $\Gamma$  instead of  $\Gamma_n$  has only finitely many solutions in positive integers  $x_1, ..., x_n$ , then every tuple  $(x_1, ..., x_n) \in (\mathbb{N} \setminus \{0\})^n$  that solves the original system S satisfies  $x_1, ..., x_n \leq 2^{2^{n-2}}$ .

**Hypothesis 3.** The statements  $\Sigma_3, \ldots, \Sigma_{16}$  are true.

**Lemma 15.** (cf. Lemma 3). For every integer  $n \in \{4, ..., 16\}$  and for every positive integers x and y,  $x \cdot \Gamma_n(x) = \Gamma_n(y)$  if and only if  $(x + 1 = y) \land (x \ge 2^{2^{n-3}} + 1)$ .

Let  $Z_9 \subseteq Q_9$  be the system of equations in Figure 6.



**Fig. 6** Construction of the system  $Z_9$ 

**Lemma 16.** For every positive integer  $x_1$ , the system  $\mathbb{Z}_9$  is solvable in positive integers  $x_2, \ldots, x_9$  if and only if  $x_1 > 2^{2^{9-4}}$  and  $x_1^2 + 1$  is prime. In this case, positive integers  $x_2, \ldots, x_9$  are uniquely determined by  $x_1$ . For every positive integer n, at most finitely many tuples  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$  begin with n and solve the system  $\mathbb{Z}_9$  with  $\Gamma$  instead of  $\Gamma_9$ .

Proof. It follows from Lemmas 3, 5, and 15.

**Lemma 17.** ([19]). The number  $(13!)^2 + 1 = 38775788043632640001$  is prime.

Lemma 18. 
$$((13!)^2 \ge 2^{2^{9-3}} + 1 = 18446744073709551617) \land (\Gamma_9((13!)^2) > 2^{2^{9-2}})$$

**Theorem 10.** The statement  $\Sigma_9$  implies the infinitude of primes of the form  $n^2 + 1$ .

Proof. It follows from Lemmas 16-18.

**Theorem 11.** (cf. Theorem 7). The statement  $\Sigma_9$  implies that any prime of the form n! + 1 with  $n \ge 2^{2^{9-3}}$  proves the infinitude of primes of the form n! + 1.

*Proof.* We leave the proof to the reader.

**Corollary 6.** Let  $\mathcal{Y}_9$  denote the set of primes of the form n! + 1. The statement  $\Sigma_9$  implies that we know an algorithm such that it returns a threshold number of  $\mathcal{Y}_9$ , and this number equals  $\max(\mathcal{Y}_9)$ , if  $\mathcal{Y}_9$  is finite. Assuming the statement  $\Sigma_9$ , a single query to an oracle for the halting problem decides the infinity of  $\mathcal{Y}_9$ . Assuming the statement  $\Sigma_9$ , the infinity of  $\mathcal{Y}_9$  is decidable in the limit.

*Proof.* We consider an algorithm which computes 
$$\max(\mathcal{Y}_9 \cap [1, (2^{2^{9-3}} - 1)! + 1])$$
.  $\Box$ 

Let  $Z_{14} \subseteq Q_{14}$  be the system of equations in Figure 7.



**Fig. 7** Construction of the system  $Z_{14}$ 

**Lemma 19.** For every positive integer  $x_1$ , the system  $\mathbb{Z}_{14}$  is solvable in positive integers  $x_2, \ldots, x_{14}$  if and only if  $x_1$  and  $x_1 + 2$  are prime and  $x_1 \ge 2^{2^{14-3}} + 1$ . In this case, positive integers  $x_2, \ldots, x_{14}$  are uniquely determined by  $x_1$ . For every positive integer n, at most finitely many tuples  $(x_1, \ldots, x_{14}) \in (\mathbb{N} \setminus \{0\})^{14}$  begin with n and solve the system  $\mathbb{Z}_{14}$  with  $\Gamma$  instead of  $\Gamma_{14}$ .

Proof. It follows from Lemmas 3, 5, and 15.

**Lemma 20.** ([23, p. 87]). The numbers  $459 \cdot 2^{8529} - 1$  and  $459 \cdot 2^{8529} + 1$  are prime (Harvey Dubner).

**Lemma 21.**  $459 \cdot 2^{8529} - 1 > 2^{2^{14-2}} = 2^{4096}$ .

**Theorem 12.** The statement  $\Sigma_{14}$  implies the infinitude of twin primes.

Proof. It follows from Lemmas 19–21.

A prime p is said to be a Sophie Germain prime if both p and 2p + 1 are prime, see [22]. It is conjectured that there are infinitely many Sophie Germain primes, see [17, p. 330]. Let  $Z_{16} \subseteq Q_{16}$  be the system of equations in Figure 8.



**Fig. 8** Construction of the system  $Z_{16}$ 

**Lemma 22.** For every positive integer  $x_1$ , the system  $Z_{16}$  is solvable in positive integers  $x_2, \ldots, x_{16}$  if and only if  $x_1$  is a Sophie Germain prime and  $x_1 \ge 2^{2^{16-3}} + 1$ . In this case, positive integers  $x_2, \ldots, x_{16}$  are uniquely determined by  $x_1$ . For every positive integer n, at most finitely many tuples  $(x_1, \ldots, x_{16}) \in (\mathbb{N} \setminus \{0\})^{16}$  begin with n and solve the system  $Z_{16}$  with  $\Gamma$  instead of  $\Gamma_{16}$ .

Proof. It follows from Lemmas 3, 5, and 15.

**Lemma 23.** ([17, p. 330]). 8069496435 · 10<sup>5072</sup> – 1 is a Sophie Germain prime (Harvey Dubner).

**Lemma 24.**  $8069496435 \cdot 10^{5072} - 1 > 2^{2^{16-2}}$ 

**Theorem 13.** The statement  $\Sigma_{16}$  implies the infinitude of Sophie Germain primes.

Proof. It follows from Lemmas 22-24.

**Theorem 14.** The statement  $\Sigma_6$  proves the following implication: if the equation x(x + 1) = y! has only finitely many solutions in positive integers x and y, then each such solution (x, y) belongs to the set  $\{(1, 2), (2, 3)\}$ .

*Proof.* We leave the proof to the reader.

The question of solving the equation x(x + 1) = y! was posed by P. Erdös, see [2]. F. Luca proved that the *abc* conjecture implies that the equation x(x + 1) = y! has only finitely many solutions in positive integers, see [12].

**Theorem 15.** The statement  $\Sigma_6$  proves the following implication: if the equation  $x! + 1 = y^2$  has only finitely many solutions in positive integers x and y, then each such solution (x, y) belongs to the set  $\{(4, 5), (5, 11), (7, 71)\}$ .

*Proof.* We leave the proof to the reader.

#### 11 Hypothetical statements $\Omega_3, \ldots, \Omega_{16}$ and their consequences

For an integer  $n \in \{3, ..., 16\}$ , let  $\Omega_n$  denote the following statement: *if a system of equations*  $S \subseteq \{\Gamma(x_i) = x_k : i, k \in \{1, ..., n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, ..., n\}\}$  has a solution in integers  $x_1, ..., x_n$  greater than  $2^{2^{n-2}}$ , then S has infinitely many solutions in positive integers  $x_1, ..., x_n$ . For every  $n \in \{3, ..., 16\}$ , the statement  $\Sigma_n$  implies the statement  $\Omega_n$ .

**Lemma 25.** The number  $(65!)^2 + 1$  is prime and  $65! > 2^{2^{9-2}}$ .

Proof. The following PARI/GP ([16]) command

# (04:04) gp > isprime((65!)^2+1,{flag=2}) %1 = 1

is shown together with its output. This command performs the APRCL primality test, the best deterministic primality test algorithm ([23, p. 226]). It rigorously shows that the number  $(65!)^2 + 1$  is prime.

**Lemma 26.** If positive integers  $x_1, \ldots, x_9$  solve the system  $\mathbb{Z}_9$  and  $x_1 > 2^{2^{9-2}}$ , then  $x_1 = \min(x_1, \ldots, x_9)$ .

**Theorem 16.** The statement  $\Omega_9$  implies the infinitude of primes of the form  $n^2 + 1$ .

Proof. It follows from Lemmas 16 and 25-26.

**Lemma 27.** If positive integers  $x_1, ..., x_{14}$  solve the system  $Z_{14}$  and  $x_1 > 2^{2^{14-2}}$ , then  $x_1 = \min(x_1, ..., x_{14})$ .

**Theorem 17.** The statement  $\Omega_{14}$  implies the infinitude of twin primes.

*Proof.* It follows from Lemmas 19–21 and 27.

#### 12 Are there infinitely many composite Fermat numbers?

Integers of the form  $2^{2^n} + 1$  are called Fermat numbers. Primes of the form  $2^{2^n} + 1$  are called Fermat primes, as Fermat conjectured that every integer of the form  $2^{2^n} + 1$  is prime, see [11, p. 1]. Fermat correctly remarked that  $2^{2^0} + 1 = 3$ ,  $2^{2^1} + 1 = 5$ ,  $2^{2^2} + 1 = 17$ ,  $2^{2^3} + 1 = 257$ , and  $2^{2^4} + 1 = 65537$  are all prime, see [11, p. 1].

**Open Problem.** ([11, p. 159]). Are there infinitely many composite numbers of the form  $2^{2^n} + 1$ ? Most mathematicians believe that  $2^{2^n} + 1$  is composite for every integer  $n \ge 5$ , see [10, p. 23]. Let

$$H_n = \{x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}\} \cup \{2^{2^{X_i}} = x_k : i, k \in \{1, \dots, n\}\}$$

Let h(1) = 1, and let  $h(n + 1) = 2^{2h(n)}$  for every positive integer *n*.

**Lemma 28.** The following subsystem of  $H_n$ 

$$\begin{cases} x_1 \cdot x_1 = x_1 \\ \forall i \in \{1, \dots, n-1\} \ 2^{2^{X_i}} = x_{i+1} \end{cases}$$

has exactly one solution  $(x_1, \ldots, x_n) \in (\mathbb{N} \setminus \{0\})^n$ , namely  $(h(1), \ldots, h(n))$ .

For a positive integer *n*, let  $\xi_n$  denote the following statement: if a system of equations  $S \subseteq H_n$  has only finitely many solutions in positive integers  $x_1, \ldots, x_n$ , then each such solution  $(x_1, \ldots, x_n)$  satisfies  $x_1, \ldots, x_n \leq h(n)$ . The statement  $\xi_n$  says that for subsystems of  $H_n$  the largest known solution is indeed the largest possible.

**Hypothesis 4.** *The statements*  $\xi_1, \ldots, \xi_{13}$  *are true.* 

**Proposition 5.** Every statement  $\xi_n$  is true with an unknown integer bound that depends on n.

*Proof.* For every positive integer n, the system  $H_n$  has a finite number of subsystems.

**Theorem 18.** The statement  $\xi_{13}$  proves the following implication: if  $z \in \mathbb{N} \setminus \{0\}$  and  $2^{2^z} + 1$  is composite and greater than h(12), then  $2^{2^z} + 1$  is composite for infinitely many positive integers z.

Proof. Let us consider the equation

$$(x+1)(y+1) = 2^{2^{z}} + 1$$
<sup>(2)</sup>

in positive integers. By Lemma 4, we can transform equation (2) into an equivalent system of equations  $\mathcal{G}$  which has 13 variables (*x*, *y*, *z*, and 10 other variables) and which consists of equations of the forms  $\alpha \cdot \beta = \gamma$  and  $2^{2^{\alpha}} = \gamma$ , see the diagram in Figure 9.



Fig. 9 Construction of the system G

Since  $2^{2^{z}} + 1 > h(12)$ , we obtain that  $2^{2^{2^{z}}+1} > h(13)$ . By this, the statement  $\xi_{13}$  implies that the system  $\mathcal{G}$  has infinitely many solutions in positive integers. It means that there are infinitely many composite Fermat numbers.

**Corollary 7.** Let  $W_{13}$  denote the set of composite Fermat numbers. The statement  $\xi_{13}$  implies that we know an algorithm such that it returns a threshold number of  $W_{13}$ , and this number equals  $\max(W_{13})$ , if  $W_{13}$  is finite. Assuming the statement  $\xi_{13}$ , a single query to an oracle for the halting problem decides the infinity of  $W_{13}$ . Assuming the statement  $\xi_{13}$ , the infinity of  $W_{13}$  is decidable in the limit.

*Proof.* We consider an algorithm which computes  $\max(W_{13} \cap [1, h(12)])$ .

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Apoloniusz Tyszka Technical Faculty Hugo Kołłątaj University Balicka 116B, 30-149 Kraków, Poland E-mail: rttyszka@cyf-kr.edu.pl