The physical limits of computation inspire an open problem that concerns abstract computable sets $X\subseteq\mathbb{N}$ and cannot be formalized in the set theory ZFC as it refers to our current knowledge on X

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Abstract. Let f(1) = 2, f(2) = 4, and let f(n+1) = f(n)! for every integer $n \ge 2$. Edmund Landau's conjecture states that the set \mathcal{P}_{n^2+1} of primes of the form $n^2 + 1$ is infinite. Landau's conjecture implies the following unproven statement Φ : $\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq [2, f(7)]$. Let B denote the system of equations: $\{x_i! = x_k : i, k \in \mathbb{N}\}$ $\{1,\ldots,9\}$ \cup $\{x_i\cdot x_j=x_k:i,j,k\in\{1,\ldots,9\}\}$. We write down a system $\mathcal{U}\subseteq B$ of 9 equations which has exactly two solutions in positive integers, namely $(1, \ldots, 1)$ and $(f(1), \ldots, f(9))$. Let Ψ denote the statement: if a system $S \subseteq B$ has at most finitely many solutions in positive integers x_1, \ldots, x_9 , then each such solution (x_1, \ldots, x_9) satisfies $x_1, \ldots, x_9 \le f(9)$. We write down a system $\mathcal{A} \subseteq B$ of 8 equations. Theorem 1. The statement Ψ restricted to the system \mathcal{A} is equivalent to the statement Φ . Open Problem. Is there a set $X \subseteq \mathbb{N}$ that satisfies conditions (1)-(5)? (1) There are many elements of X and it is conjectured that Xis infinite. (2) No known algorithm decides the finiteness/infiniteness of X. (3) There is a known algorithm that for every $k \in \mathbb{N}$ decides whether or not $k \in X$. (4) There is a known algorithm that computes an integer n satisfying $card(X) < \omega \Rightarrow X \subseteq (-\infty, n]$. (5) There is a simple condition C, which can be formalized in ZFC, such that for almost all $k \in \mathbb{N}$, k satisfies the condition C if and only if $k \in X$. The simplest known such condition C defines in \mathbb{N} the set X. Condition (5) excludes non-naturally defined sets $X \subseteq \mathbb{N}$ as that in the statement (i). We prove: (i) the set $X = \{k \in \mathbb{N} : (f(7) < k) \Rightarrow (f(7), k) \cap \mathcal{P}_{n^2+1} \neq \emptyset \}$ satisfies conditions (1)-(4); (ii) the set $X = \{1\} \cup \mathcal{P}_{n^2+1}$ satisfies conditions (1)-(3) and (5). Proving Landau's conjecture will disprove statements (i) and (ii). Theorem 2. The statement Φ implies that the set $X = \{1\} \cup \mathcal{P}_{n^2+1}$ satisfies condition (4). Theorem 3. No set $X \subseteq \mathbb{N}$ will satisfy conditions (1)-(4) forever, if for every algorithm with no inputs that operates on integers, at some future day, a computer will be able to execute this algorithm in 1 second or less. The assumption of Theorem 3 contradicts the current paradigm of physics. Theorem 3 and the Open Problem explain the title of the article.

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1. Basic definitions and the philosophical goal of the article

Logicism is a programme in the philosophy of mathematics. It is mainly characterized by the contention that mathematics can be reduced to logic, provided that the latter includes set theory, see [3, p. 199].

Definition 1. Conditions (1)–(5) concern sets $X \subseteq \mathbb{N}$.

- (1) There are many elements of X and it is conjectured that X is infinite.
- (2) No known algorithm decides the finiteness/infiniteness of X.
- (3) There is a known algorithm that for every $k \in \mathbb{N}$ decides whether or not $k \in X$.
- (4) There is a known algorithm that computes an integer n satisfying $\operatorname{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$.
- (5) There is a simple condition C, which can be formalized in ZFC, such that for almost all $k \in \mathbb{N}$, k satisfies the condition C if and only if $k \in X$. The simplest known such condition C defines in \mathbb{N} the set X.

Condition (5) excludes non-naturally defined sets $X \subseteq \mathbb{N}$ as that in Statement 2.

Definition 2. We say that an integer n is a threshold number of a set $X \subseteq \mathbb{N}$, if $\operatorname{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$, cf. [7] and [8].

If a set $X \subseteq \mathbb{N}$ is empty or infinite, then any integer n is a threshold number of X. If a set $X \subseteq \mathbb{N}$ is non-empty and finite, then the all threshold numbers of X form the set $[\max(X), \infty) \cap \mathbb{N}$.

Edmund Landau's conjecture states that the set \mathcal{P}_{n^2+1} of primes of the form $n^2 + 1$ is infinite, see [4]–[6].

Definition 3. Let Φ denote the following unproven statement:

$$\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty, (((24!)!)!)!]$$

Landau's conjecture implies the statement Φ . In Section 4, we heuristically justify the statement Φ without invoking Landau's conjecture.

Statement 1. No known algorithm computes an integer k such that

$$\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty, k]$$

Proving the statement Φ will disprove Statement 1. Statement 1 cannot be formalized in ZFC because it refers to the current mathematical knowledge. The same is true for Statements 2–3 and Open Problem 1 in the next sections. It argues against logicism as Open Problem 1 concerns abstract computable sets $\mathcal{X} \subseteq \mathbb{N}$.

2. The physical limits of computation inspire Open Problem 1

Definition 4. *Let* $\beta = (((24!)!)!)!$.

Lemma 1. $\log_2(\log_2(\log_2(\log_2(\log_2(\log_2(\beta)))))) \approx 1.42298$.

Proof. We ask Wolfram Alpha at http://wolframalpha.com.

Statement 2. The set $X = \{k \in \mathbb{N} : (\beta < k) \Rightarrow (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$ satisfies conditions (1)-(4).

Proof. Condition (1) holds as $X \supseteq \{0, \dots, \beta\}$ and the set \mathcal{P}_{n^2+1} is conjecturally infinite. By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of \mathcal{P}_{n^2+1} is greater than β , see [2]. Thus condition (2) holds. Condition (3) holds trivially. Since the set

$$\{k \in \mathbb{N} : (\beta < k) \land (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$$

is empty or infinite, the integer β is a threshold number of X. Thus condition (4) holds.

Proving Landau's conjecture will disprove Statement 2.

Open Problem 1. *Is there a set* $X \subseteq \mathbb{N}$ *that satisfies conditions* (1)-(5)?

Theorem 1. No set $X \subseteq \mathbb{N}$ will satisfy conditions (1)–(4) forever, if for every algorithm with no inputs that operates on integers, at some future day, a computer will be able to execute this algorithm in 1 second or less.

Proof. The proof goes by contradiction. Since conditons (2)–(4) will hold forever, the algorithm in Figure 1 never terminates and sequentially prints the following sentences:

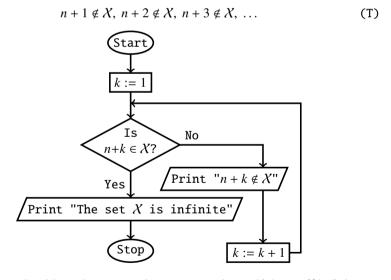


Fig. 1 An algorithm whose execution never terminates if the set X is finite

The sentences from the sequence (T) and our assumption imply that for every integer m > n computed by a known algorithm, at some future day, a computer will be able to confirm in 1 second or less that $(n, m] \cap \mathcal{X} = \emptyset$. Thus, at some future day, numerical evidence will support the conjecture that the set \mathcal{X} is finite, contrary to the conjecture in condition (1).

The assumption of Theorem 1 contradicts the current paradigm of physics. Theorem 1 and Open Problem 1 explain the title of the article.

3. Number-theoretic statements Ψ_n

Let f(1) = 2, f(2) = 4, and let f(n + 1) = f(n)! for every integer $n \ge 2$. Let \mathcal{U}_1 denote the system of equations which consists of the equation $x_1! = x_1$. For an integer $n \ge 2$, let \mathcal{U}_n denote the following system of equations:

$$\begin{cases} x_1! &= x_1 \\ x_1 \cdot x_1 &= x_2 \\ \forall i \in \{2, \dots, n-1\} x_i! &= x_{i+1} \end{cases}$$

The diagram in Figure 2 illustrates the construction of the system \mathcal{U}_n .

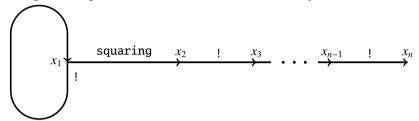


Fig. 2 Construction of the system \mathcal{U}_n

Lemma 2. For every positive integer n, the system \mathcal{U}_n has exactly two solutions in positive integers, namely (1, ..., 1) and (f(1), ..., f(n)).

Let B_n denote the following system of equations:

$${x_i! = x_k : i, k \in \{1, \dots, n\}} \cup {x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}}$$

For a positive integer n, let Ψ_n denote the following statement: if a system of equations $S \subseteq B_n$ has at most finitely many solutions in positive integers x_1, \ldots, x_n , then each such solution (x_1, \ldots, x_n) satisfies $x_1, \ldots, x_n \le f(n)$. The statement Ψ_n says that for subsystems of B_n with a finite number of solutions, the largest known solution is indeed the largest possible. The statements Ψ_1 and Ψ_2 hold trivially. The statement $\forall n \in \mathbb{N} \setminus \{0\} \ \Psi_n$ is very doubtful.

Theorem 2. For every statement Ψ_n , the bound f(n) cannot be decreased.

Proof. It follows from Lemma 2 because $\mathcal{U}_n \subseteq B_n$.

Theorem 3. For every integer $n \ge 2$, the statement Ψ_{n+1} implies the statement Ψ_n .

Proof. If a system $S \subseteq B_n$ has at most finitely many solutions in positive integers x_1, \ldots, x_n , then for every integer $i \in \{1, \ldots, n\}$ the system $S \cup \{x_i! = x_{n+1}\}$ has at most finitely many solutions in positive integers x_1, \ldots, x_{n+1} . The statement Ψ_{n+1} implies that $x_i! = x_{n+1} \le f(n+1) = f(n)!$. Hence, $x_i \le f(n)$.

Theorem 4. Every statement Ψ_n is true with an unknown integer bound that depends on n.

Proof. For every positive integer n, the system B_n has a finite number of subsystems.

4. A conjectural solution to Open Problem 1

Lemma 3. For every positive integers x and y, $x! \cdot y = y!$ if and only if

$$(x + 1 = y) \lor (x = y = 1)$$

Lemma 4. (Wilson's theorem, [1, p. 89]). For every integer $x \ge 2$, x is prime if and only if x divides (x - 1)! + 1.

Let \mathcal{A} denote the following system of equations:

$$\begin{cases} x_2! &= x_3 \\ x_3! &= x_4 \\ x_5! &= x_6 \\ x_8! &= x_9 \\ x_1 \cdot x_1 &= x_2 \\ x_3 \cdot x_5 &= x_6 \\ x_4 \cdot x_8 &= x_9 \\ x_5 \cdot x_7 &= x_8 \end{cases}$$

Lemma 3 and the diagram in Figure 3 explain the construction of the system \mathcal{A} .

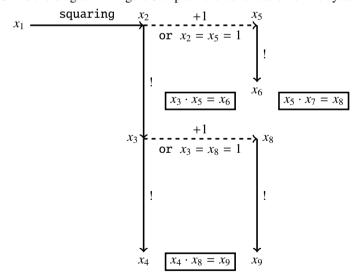


Fig. 3 Construction of the system \mathcal{A}

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Lemma 5. For every integer $x_1 \ge 2$, the system \mathcal{A} is solvable in positive integers x_2, \ldots, x_9 if and only if $x_1^2 + 1$ is prime. In this case, the integers x_2, \ldots, x_9 are uniquely determined by the following equalities:

$$x_{2} = x_{1}^{2}$$

$$x_{3} = (x_{1}^{2})!$$

$$x_{4} = ((x_{1}^{2})!)!$$

$$x_{5} = x_{1}^{2} + 1$$

$$x_{6} = (x_{1}^{2} + 1)!$$

$$x_{7} = \frac{(x_{1}^{2})! + 1}{x_{1}^{2} + 1}$$

$$x_{8} = (x_{1}^{2})! + 1$$

$$x_{9} = ((x_{1}^{2})! + 1)!$$

Proof. By Lemma 3, for every integer $x_1 \ge 2$, the system \mathcal{H} is solvable in positive integers x_2, \ldots, x_9 if and only if $x_1^2 + 1$ divides $(x_1^2)! + 1$. Hence, the claim of Lemma 5 follows from Lemma 4.

Lemma 6. There are only finitely many tuples $(x_1, ..., x_9) \in (\mathbb{N} \setminus \{0\})^9$, which solve the system \mathcal{A} and satisfy $x_1 = 1$. This is true as every such tuple $(x_1, ..., x_9)$ satisfies $x_1, ..., x_9 \in \{1, 2\}$.

Proof. The equality $x_1 = 1$ implies that $x_2 = x_1 \cdot x_1 = 1$. Hence, $x_3 = x_2! = 1$. Therefore, $x_4 = x_3! = 1$. The equalities $x_5! = x_6$ and $x_5 = 1 \cdot x_5 = x_3 \cdot x_5 = x_6$ imply that $x_5, x_6 \in \{1, 2\}$. The equalities $x_8! = x_9$ and $x_8 = 1 \cdot x_8 = x_4 \cdot x_8 = x_9$ imply that $x_8, x_9 \in \{1, 2\}$. The equality $x_5 \cdot x_7 = x_8$ implies that $x_7 = \frac{x_8}{x_5} \in \left\{\frac{1}{1}, \frac{1}{2}, \frac{2}{1}, \frac{2}{2}\right\} \cap \mathbb{N} = \{1, 2\}$. □

Conjecture 1. The statement Ψ_9 is true when is restricted to the system \mathcal{A} .

Theorem 5. Conjecture 1 proves the following implication: if there exists an integer $x_1 \ge 2$ such that $x_1^2 + 1$ is prime and greater than f(7), then the set \mathcal{P}_{n^2+1} is infinite.

Proof. Suppose that the antecedent holds. By Lemma 5, there exists a unique tuple $(x_2, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^8$ such that the tuple (x_1, x_2, \ldots, x_9) solves the system \mathcal{A} . Since $x_1^2 + 1 > f(7)$, we obtain that $x_1^2 \ge f(7)$. Hence, $(x_1^2)! \ge f(7)! = f(8)$. Consequently,

$$x_9 = ((x_1^2)! + 1)! \ge (f(8) + 1)! > f(8)! = f(9)$$

Conjecture 1 and the inequality $x_9 > f(9)$ imply that the system \mathcal{A} has infinitely many solutions $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$. According to Lemmas 5 and 6, the set \mathcal{P}_{n^2+1} is infinite.

Theorem 6. Conjecture 1 implies the statement Φ .

Proof. It follows from Theorem 5 and the equality f(7) = (((24!)!)!)!.

Theorem 7. The statement Φ implies Conjecture 1.

Proof. By Lemmas 5 and 6, if positive integers x_1, \ldots, x_9 solve the system \mathcal{A} , then

$$(x_1 \ge 2) \land (x_5 = x_1^2 + 1) \land (x_5 \text{ is prime})$$

or $x_1, \ldots, x_9 \in \{1, 2\}$. In the first case, Lemma 5 and the statement Φ imply that the inequality $x_5 \leq (((24!)!)!)! = f(7)$ holds when the system \mathcal{A} has at most finitely many solutions in positive integers x_1, \ldots, x_9 . Hence, $x_2 = x_5 - 1 < f(7)$ and $x_3 = x_2! < f(7)! = f(8)$. Continuing this reasoning in the same manner, we can show that every x_i does not exceed f(9).

Statement 3. The set $X = \{1\} \cup \mathcal{P}_{n^2+1}$ satisfies conditions (1)-(3) and (5).

Proof. The set \mathcal{P}_{n^2+1} is conjecturally infinite. There are 2199894223892 primes of the form n^2+1 in the interval $[2,10^{28})$, see [5]. These two facts imply condition (1). Condition (3) holds trivially. By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of $\{1\} \cup \mathcal{P}_{n^2+1}$ is greater than $f(7) = (((24!)!)!)! = \beta$, see [2]. Thus condition (2) holds. The following condition:

k-1 is a square and k has no divisors greater than 1 and smaller than k defines in \mathbb{N} the set $\{1\} \cup \mathcal{P}_{n^2+1}$. This proves condition (5).

Proving Landau's conjecture will disprove Statement 3.

Theorem 8. The statement Φ implies that the set $X = \{1\} \cup \mathcal{P}_{n^2+1}$ satisfies condition (4).

Proof. Suppose that the statement Φ is true. This means that β is a threshold number of $X = \{1\} \cup \mathcal{P}_{n^2+1}$. Thus condition (4) holds.

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