# On sets $\mathcal{W} \subseteq \mathbb{N}$ whose infinity follows from the existence in $\mathcal{W}$ of an element which is greater than a threshold number computed for $\mathcal{W}$ 

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#### Abstract

Let $f(1)=2, f(2)=4$, and let $f(n+1)=f(n)$ ! for every integer $n \geqslant 2$. For a positive integer $n$, let $\Theta_{n}$ denote the statement: if a system $\mathcal{S} \subseteq\left\{x_{i}!=x_{k}: i, k \in\{1, \ldots, n\}\right\} \cup$ $\left\{x_{i} \cdot x_{j}=x_{k}: i, j, k \in\{1, \ldots, n\}\right\}$ has only finitely many solutions in integers $x_{1}, \ldots, x_{n}$ greater than 1 , then each such solution $\left(x_{1}, \ldots, x_{n}\right)$ satisfies $\min \left(x_{1}, \ldots, x_{n}\right) \leqslant f(n)$. The statement $\Theta_{9}$ proves that if there exists an integer $x>f(9)$ such that $x^{2}+1$ (alternatively, $x!+1$ ) is prime, then there are infinitely many primes of the form $n^{2}+1$ (respectively, $n!+1$ ). The statement $\Theta_{16}$ proves that if there exists a twin prime greater than $f(16)+3$, then there are infinitely many twin primes. We formulate the statements $\Phi_{n}$ and prove: $\Phi_{4}$ equivalently expresses that there are infinitely many primes of the form $n!+1, \Phi_{6}$ implies that for infinitely many primes $p$ the number $p!+1$ is prime, $\Phi_{6}$ implies that there are infinitely many primes of the form $n!-1, \Phi_{7}$ implies that there are infinitely many twin primes.

Index Terms-composite Fermat numbers, prime numbers of the form $n!+1$, prime numbers of the form $n!-1$, prime numbers of the form $n^{2}+1$, prime numbers $p$ such that $p!+1$ is prime, single query to an oracle for the halting problem, twin prime conjecture.


## I. Spectra of sentences and their threshold numbers

TIHE following observation concerns the theme described in the title of the article.

Observation 1. If $\mathcal{W}$ is a subset of $\{0, \ldots, n\}$ where $n$ is a non-negative integer, then we take any integer $m \geqslant n$ as a threshold number for $\mathcal{W}$. If $\mathcal{W}$ is an infinite subset of $\mathbb{N}$, then we take any non-negative integer $m$ as a threshold number for $\mathcal{W}$.

We define the set $\mathcal{U} \subseteq \mathbb{N}$ by declaring that a non-negative integer $n$ belongs to $\mathcal{U}$ if and only if $\sin \left(10^{10^{10^{10}}}\right)>0$. This inequality is practically undecidable, see [5].
Corollary 1. The set $\mathcal{U}$ equals $\emptyset$ or $\mathbb{N}$. The statement " $\mathcal{U}=\emptyset$ " remains unproven and the statement " $\mathcal{U}=\mathbb{N}$ " remains unproven. Every non-negative integer $m$ is a threshold number for $\mathcal{U}$. For every non-negative integer $k$, the sentence " $k \in \mathcal{U}$ " is only theoretically decidable.

The first-order language of graph theory contains two relation symbols of arity $2: \sim$ and $=$, respectively for adjacency and equality of vertices. The term first-order imposes the condition that the variables represent vertices and hence the quantifiers apply to vertices only. For a first-order sentence $\Lambda$ about graphs, let $\operatorname{Spectrum}(\Lambda)$ denote the set of all positive integers $n$ such that there is a graph on $n$ vertices satisfying $\Lambda$. By a graph on $n$ vertices we understand a set of $n$ elements with a binary relation which is symmetric and irreflexive.

Theorem 1. ([]12 p. 171]). If a sentence $\Lambda$ in the language of graph theory has the form $\exists x_{1} \ldots x_{k} \forall y_{1} \ldots y_{l}$ $\Upsilon\left(x_{1}, \ldots, x_{k}, y_{1}, \ldots, y_{l}\right)$, where $\Upsilon\left(x_{1}, \ldots, x_{k}, y_{1}, \ldots, y_{l}\right)$ is quantifier-free, then either $\operatorname{Spectrum}(\Lambda) \subseteq\left[1,\left(2^{k} \cdot 4^{l}\right)-1\right]$ or $\operatorname{Spectrum}(\Lambda) \supseteq[k+l, \infty) \cap \mathbb{N}$.
Corollary 2. The number $\left(2^{k} \cdot 4^{l}\right)-1$ is a threshold number for $\operatorname{Spectrum}(\Lambda)$.

The classes of the infinite recursively enumerable sets and of the infinite recursive sets are not recursively enumerable, see [10, p. 234].
Corollary 3. If an algorithm $\mathrm{Alg}_{1}$ for every recursive set $\mathcal{W} \subseteq \mathbb{N}$ finds a non-negative integer $\operatorname{Alg}_{1}(\mathcal{W})$, then there exists a finite set $\mathcal{M} \subseteq \mathbb{N}$ such that $\mathcal{M} \cap\left[\operatorname{Alg}_{1}(\mathcal{M})+1, \infty\right) \neq \emptyset$.

Corollary 4. If an algorithm $\mathrm{Alg}_{2}$ for every recursively enumerable set $\mathcal{W} \subseteq \mathbb{N}$ finds a non-negative integer $\operatorname{Alg}_{2}(\mathcal{W})$, then there exists a finite set $\mathcal{M} \subseteq \mathbb{N}$ such that $\mathcal{M} \cap\left[\mathrm{Alg}_{2}(\mathcal{M})+1, \infty\right) \neq \emptyset$.

$$
\text { Let } K=\left\{j \in \mathbb{N}: 2^{\boldsymbol{\aleph}_{j}}=\boldsymbol{\aleph}_{j+1}\right\}
$$

Theorem 2. If ZFC is consistent, then for every non-negative integer $n$ the sentence
$" n$ is a threshold number for $K "$
is not provable in $Z F C$.

Proof. There exists a model $\mathcal{E}$ of ZFC such that

$$
\forall i \in\{0, \ldots, n+1\} \mathcal{E} \vDash 2^{\boldsymbol{\aleph}_{i}}=\boldsymbol{\aleph}_{i+1}
$$

and

$$
\forall i \in\{n+2, n+3, n+4, \ldots\} \mathcal{E} \vDash 2^{\boldsymbol{\aleph}_{i}}=\boldsymbol{\aleph}_{i+2}
$$

see [3] and [6, p. 232]. In the model $\mathcal{E}, K=\{0, \ldots, n+1\}$ and $n$ is not a threshold number for $K$.

Theorem 3. If ZFC is consistent, then for every non-negative integer $n$ the sentence
$" n$ is not a threshold number for $K "$
is not provable in ZFC.
Proof. The Generalized Continuum Hypothesis (GCH) is consistent with ZFC, see [6, p. 188] and [6, p. 190]. GCH implies that $K=\mathbb{N}$. Consequently, GCH implies that every non-negative integer $n$ is a threshold number for $K$.

## II. Basic lemmas

Let $f(1)=2, f(2)=4$, and let $f(n+1)=f(n)$ ! for every integer $n \geqslant 2$. Let $\mathcal{V}_{1}$ denote the system of equations $\left\{x_{1}!=x_{1}\right\}$, and let $\mathcal{V}_{2}$ denote the system of equations $\left\{x_{1}!=x_{1}, x_{1} \cdot x_{1}=x_{2}\right\}$. For an integer $n \geqslant 3$, let $\mathcal{V}_{n}$ denote the following system of equations:

$$
\left\{\begin{aligned}
x_{1}! & =x_{1} \\
x_{1} \cdot x_{1} & =x_{2} \\
\forall i \in\{2, \ldots, n-1\} x_{i}! & =x_{i+1}
\end{aligned}\right.
$$

The diagram in Figure 1 illustrates the construction of the system $\mathcal{V}_{n}$.


Fig. 1 Construction of the system $\mathcal{V}_{n}$
Lemma 1. For every positive integer $n$, the system $\mathcal{V}_{n}$ has exactly one solution in integers greater than 1, namely $(f(1), \ldots, f(n))$.

Let

$$
\begin{array}{r}
H_{n}=\left\{x_{i}!=x_{k}: \quad i, k \in\{1, \ldots, n\}\right\} \cup \\
\left\{x_{i} \cdot x_{j}=x_{k}: \quad i, j, k \in\{1, \ldots, n\}\right\}
\end{array}
$$

For a positive integer $n$, let $\Theta_{n}$ denote the following statement: if a system $\mathcal{S} \subseteq H_{n}$ has at most finitely many solutions in integers $x_{1}, \ldots, x_{n}$ greater than 1 , then each such solution $\left(x_{1}, \ldots, x_{n}\right)$ satisfies $\min \left(x_{1}, \ldots, x_{n}\right) \leqslant f(n)$. The assumption $\min \left(x_{1}, \ldots, x_{n}\right) \leqslant f(n)$ is weaker than the assumption $\max \left(x_{1}, \ldots, x_{n}\right) \leqslant f(n)$ suggested by Lemma 1
Lemma 2. For every positive integer $n$, the system $H_{n}$ has a finite number of subsystems.

Theorem 4. Every statement $\Theta_{n}$ is true with an unknown integer bound that depends on $n$.
Proof. It follows from Lemma 2.
Lemma 3. For every integers $x$ and $y$ greater than 1 , $x!\cdot y=y!$ if and only if $x+1=y$.
Lemma 4. If $x \geqslant 4$, then $\frac{(x-1)!+1}{x}>1$.
Lemma 5. (Wilson's theorem, [4] p. 89]). For every integer $x \geqslant 2, x$ is prime if and only if $x$ divides $(x-1)!+1$.

## III. Brocard's problem

A weak form of Szpiro's conjecture implies that there are only finitely many solutions to the Brocard-Ramanujan equation $x!+1=y^{2}$, see [11]. It is conjectured that $x!+1$ is a square only for $x \in\{4,5,7\}$, see [18, p. 297].

Let $\mathcal{A}$ denote the following system of equations:

$$
\left\{\begin{aligned}
x_{1}! & =x_{2} \\
x_{2}! & =x_{3} \\
x_{5}! & =x_{6} \\
x_{4} \cdot x_{4} & =x_{5} \\
x_{3} \cdot x_{5} & =x_{6}
\end{aligned}\right.
$$

Lemma 3 and the diagram in Figure 2 explain the construction of the system $\mathcal{A}$.


Fig. 2 Construction of the system $\mathcal{A}$
Lemma 6. For every integers $x_{1}$ and $x_{4}$ greater than 1 , the system $\mathcal{A}$ is solvable in integers $x_{2}, x_{3}, x_{5}, x_{6}$ greater than 1 if and only if $x_{1}!+1=x_{4}^{2}$. In this case, the integers $x_{2}, x_{3}, x_{5}, x_{6}$ are uniquely determined by the following equalities:

$$
\begin{aligned}
& x_{2}=x_{1}! \\
& x_{3}=\left(x_{1}!\right)! \\
& x_{5}=x_{1}!+1 \\
& x_{6}=\left(x_{1}!+1\right)!
\end{aligned}
$$

and $x_{1}=\min \left(x_{1}, \ldots, x_{6}\right)$.
Proof. It follows from Lemma 3.
Theorem 5. The statement $\Theta_{6}$ proves the following implication: if the equation $x_{1}!+1=x_{4}^{2}$ has only finitely many solutions in positive integers, then each such solution $\left(x_{1}, x_{4}\right)$ satisfies $x_{1} \leqslant f(6)$.
Proof. Let positive integers $x_{1}$ and $x_{4}$ satisfy $x_{1}!+1=x_{4}^{2}$. Then, $x_{1}, x_{4} \in \mathbb{N} \backslash\{0,1\}$. By Lemma 6, there exists a
unique tuple $\left(x_{2}, x_{3}, x_{5}, x_{6}\right) \in(\mathbb{N} \backslash\{0,1\})^{4}$ such that the tuple $\left(x_{1}, \ldots, x_{6}\right)$ solves the system $\mathcal{A}$. Lemma 6 guarantees that $x_{1}=\min \left(x_{1}, \ldots, x_{6}\right)$. By the antecedent and Lemma 6 , the system $\mathcal{A}$ has only finitely many solutions in integers $x_{1}, \ldots, x_{6}$ greater than 1 . Therefore, the statement $\Theta_{6}$ implies that $x_{1}=\min \left(x_{1}, \ldots, x_{6}\right) \leqslant f(6)$.

Hypothesis 1. The implication in Theorem 5 is true.
Corollary 5. Assuming Hypothesis [1] a single query to an oracle for the halting problem decides the problem of the infinitude of the solutions of the equation $x!+1=y^{2}$.
IV. Are there infinitely many prime numbers of the form

$$
n^{2}+1 ?
$$

Landau's conjecture states that there are infinitely many primes of the form $n^{2}+1$, see [9, pp. 37-38].

Let $\mathcal{B}$ denote the following system of equations:

$$
\left\{\begin{aligned}
x_{2}! & =x_{3} \\
x_{3}! & =x_{4} \\
x_{5}! & =x_{6} \\
x_{8}! & =x_{9} \\
x_{1} \cdot x_{1} & =x_{2} \\
x_{3} \cdot x_{5} & =x_{6} \\
x_{4} \cdot x_{8} & =x_{9} \\
x_{5} \cdot x_{7} & =x_{8}
\end{aligned}\right.
$$

Lemma 3 and the diagram in Figure 3 explain the construction of the system $\mathcal{B}$.


Fig. 3 Construction of the system $\mathcal{B}$
Lemma 7. For every integer $x_{1} \geqslant 2$, the system $\mathcal{B}$ is solvable in integers $x_{2}, \ldots, x_{9}$ greater than 1 if and only if $x_{1}^{2}+1$
is prime. In this case, the integers $x_{2}, \ldots, x_{9}$ are uniquely determined by the following equalities:

$$
\begin{aligned}
x_{2} & =x_{1}^{2} \\
x_{3} & =\left(x_{1}^{2}\right)! \\
x_{4} & =\left(\left(x_{1}^{2}\right)!\right)! \\
x_{5} & =x_{1}^{2}+1 \\
x_{6} & =\left(x_{1}^{2}+1\right)! \\
x_{7} & =\frac{\left(x_{1}^{2}\right)!+1}{x_{1}^{2}+1} \\
x_{8} & =\left(x_{1}^{2}\right)!+1 \\
x_{9} & =\left(\left(x_{1}^{2}\right)!+1\right)!
\end{aligned}
$$

and $\min \left(x_{1}, \ldots, x_{9}\right)=x_{1}$.
Proof. By Lemmas 3 and 4 for every integer $x_{1} \geqslant 2$, the system $\mathcal{B}$ is solvable in integers $x_{2}, \ldots, x_{9}$ greater than 1 if and only if $x_{1}^{2}+1$ divides $\left(x_{1}^{2}\right)!+1$. Hence, the claim of Lemma 7 follows from Lemma 5

Theorem 6. The statement $\Theta_{9}$ proves the following implication: if there exists an integer $x_{1}>f(9)$ such that $x_{1}^{2}+1$ is prime, then there are infinitely many primes of the form $n^{2}+1$.

Proof. Assume that an integer $x_{1}$ is greater than $f(9)$ and $x_{1}^{2}+1$ is prime. By Lemma 7, there exists a unique tuple $\left(x_{2}, \ldots, x_{9}\right) \in(\mathbb{N} \backslash\{0,1\})^{8}$ such that the tuple $\left(x_{1}, x_{2}, \ldots, x_{9}\right)$ solves the system $\mathcal{B}$. Lemma 7 guarantees that $\min \left(x_{1}, \ldots, x_{9}\right)=x_{1}$. Since $\mathcal{B} \subseteq H_{9}$, the statement $\Theta_{9}$ and the inequality $\min \left(x_{1}, \ldots, x_{9}\right)=x_{1}>f(9)$ imply that the system $\mathcal{B}$ has infinitely many solutions $\left(x_{1}, \ldots, x_{9}\right) \in(\mathbb{N} \backslash\{0,1\})^{9}$. According to Lemma 7, there are infinitely many primes of the form $n^{2}+1$.

Hypothesis 2. The implication in Theorem 6 is true.
Corollary 6. Assuming Hypothesis 2, a single query to an oracle for the halting problem decides the problem of the infinitude of primes of the form $n^{2}+1$.
V. Are there infinitely many prime numbers of the form

$$
n!+1 ?
$$

It is conjectured that there are infinitely many primes of the form $n!+1$, see [1, p. 443] and [14]. Let $\mathcal{G}$ denote the following system of equations:

$$
\left\{\begin{aligned}
x_{1}! & =x_{2} \\
x_{2}! & =x_{3} \\
x_{3}! & =x_{4} \\
x_{5}! & =x_{6} \\
x_{8}! & =x_{9} \\
x_{3} \cdot x_{5} & =x_{6} \\
x_{4} \cdot x_{8} & =x_{9} \\
x_{5} \cdot x_{7} & =x_{8}
\end{aligned}\right.
$$

Lemma 3 and the diagram in Figure 4 explain the construction of the system $\mathcal{G}$.


Fig. 4 Construction of the system $\mathcal{G}$
Lemma 8. For every integer $x_{1} \geqslant 2$, the system $\mathcal{G}$ is solvable in integers $x_{2}, \ldots, x_{9}$ greater than 1 if and only if $x_{1}!+1$ is prime. In this case, the integers $x_{2}, \ldots, x_{9}$ are uniquely determined by the following equalities:

$$
\begin{aligned}
& x_{2}=x_{1}! \\
& x_{3}=\left(x_{1}!\right)! \\
& x_{4}=\left(\left(x_{1}!\right)!\right)! \\
& x_{5}=x_{1}^{!}+1 \\
& x_{6}=\left(x_{1}!+1\right)! \\
& x_{7}=\frac{\left(x_{1}!\right)!+1}{x_{1}!+1} \\
& x_{8}=\left(x_{1}!\right)!+1 \\
& x_{9}=\left(\left(x_{1}!\right)!+1\right)!
\end{aligned}
$$

and $\min \left(x_{1}, \ldots, x_{9}\right)=x_{1}$.
Proof. By Lemmas 3 and 4, for every integer $x_{1} \geqslant 2$, the system $\mathcal{G}$ is solvable in integers $x_{2}, \ldots, x_{9}$ greater than 1 if and only if $x_{1}!+1$ divides $\left(x_{1}!\right)!+1$. Hence, the claim of Lemma 8 follows from Lemma 5

Theorem 7. The statement $\Theta_{9}$ proves the following implication: if there exists an integer $x_{1}>f(9)$ such that $x_{1}!+1$ is prime, then there are infinitely many primes of the form $n!+1$.

Proof. Assume that an integer $x_{1}$ is greater than $f(9)$ and $x_{1}!+1$ is prime. By Lemma 8, there exists a unique tuple $\left(x_{2}, \ldots, x_{9}\right) \in(\mathbb{N} \backslash\{0,1\})^{8}$ such that the tuple $\left(x_{1}, x_{2}, \ldots, x_{9}\right)$ solves the system $\mathcal{G}$. Lemma 8 guarantees that $\min \left(x_{1}, \ldots, x_{9}\right)=x_{1}$. Since $\mathcal{G} \subseteq H_{9}$, the statement $\Theta_{9}$ and the inequality $\min \left(x_{1}, \ldots, x_{9}\right)=x_{1}>f(9)$ imply that the system $\mathcal{G}$ has infinitely many solutions $\left(x_{1}, \ldots, x_{9}\right) \in(\mathbb{N} \backslash\{0,1\})^{9}$. According to Lemma 8, there are infinitely many primes of the form $n!+1$.

Hypothesis 3. The implication in Theorem 7 is true.
Corollary 7. Assuming Hypothesis 3 a single query to an oracle for the halting problem decides the problem of the infinitude of primes of the form $n!+1$.

## VI. The twin prime conjecture

A twin prime is a prime number that is either 2 less or 2 more than another prime number. The twin prime conjecture states that there are infinitely many twin primes, see [9, p. 39].

Let $C$ denote the following system of equations:

$$
\left\{\begin{aligned}
x_{1}! & =x_{2} \\
x_{2}! & =x_{3} \\
x_{4}! & =x_{5} \\
x_{6}! & =x_{7} \\
x_{7}! & =x_{8} \\
x_{9}! & =x_{10} \\
x_{12}! & =x_{13} \\
x_{15}! & =x_{16} \\
x_{2} \cdot x_{4} & =x_{5} \\
x_{5} \cdot x_{6} & =x_{7} \\
x_{7} \cdot x_{9} & =x_{10} \\
x_{4} \cdot x_{11} & =x_{12} \\
x_{3} \cdot x_{12} & =x_{13} \\
x_{9} \cdot x_{14} & =x_{15} \\
x_{8} \cdot x_{15} & =x_{16}
\end{aligned}\right.
$$

Lemma 3 and the diagram in Figure 5 explain the construction of the system $C$.


Fig. 5 Construction of the system $C$
Lemma 9. If $x_{4}=2$, then the system $C$ has no solutions in integers $x_{1}, \ldots, x_{16}$ greater than 1.

Proof. The equality $x_{2} \cdot x_{4}=x_{5}=x_{4}$ ! and the equality $x_{4}=2$ imply that $x_{2}=1$.

Lemma 10. If $x_{4}=3$, then the system $C$ has no solutions in integers $x_{1}, \ldots, x_{16}$ greater than 1 .

Proof. The equality $x_{4} \cdot x_{11}=x_{12}=\left(x_{4}-1\right)!+1$ and the equality $x_{4}=3$ imply that $x_{11}=1$.

Lemma 11. For every $x_{4} \in \mathbb{N} \backslash\{0,1,2,3\}$ and for every $x_{9} \in \mathbb{N} \backslash\{0,1\}$, the system $C$ is solvable in integers $x_{1}, x_{2}, x_{3}$, $x_{5}, x_{6}, x_{7}, x_{8}, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$ greater than 1 if and only if $x_{4}$ and $x_{9}$ are prime and $x_{4}+2=x_{9}$. In this case, the integers $x_{1}, x_{2}, x_{3}, x_{5}, x_{6}, x_{7}, x_{8}, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}$, $x_{15}, x_{16}$ are uniquely determined by the following equalities:

$$
\begin{aligned}
x_{1} & =x_{4}-1 \\
x_{2} & =\left(x_{4}-1\right)! \\
x_{3} & =\left(\left(x_{4}-1\right)!\right)! \\
x_{5} & =x_{4}! \\
x_{6} & =x_{9}-1 \\
x_{7} & =\left(x_{9}-1\right)! \\
x_{8} & =\left(\left(x_{9}-1\right)!\right)! \\
x_{10} & =x_{9}! \\
x_{11} & =\frac{\left(x_{4}-1\right)!+1}{x_{4}} \\
x_{12} & =\left(x_{4}-1\right)!+1 \\
x_{13} & =\left(\left(x_{4}-1\right)!+1\right)! \\
x_{14} & =\frac{\left(x_{9}-1\right)!+1}{x_{9}} \\
x_{15} & =\left(x_{9}-1\right)!+1 \\
x_{16} & =\left(\left(x_{9}-1\right)!+1\right)!
\end{aligned}
$$

and $\min \left(x_{1}, \ldots, x_{16}\right)=x_{1}=x_{9}-3$.
Proof. By Lemmas 3 and 4 for every $x_{4} \in \mathbb{N} \backslash\{0,1,2,3\}$ and for every $x_{9} \in \mathbb{N} \backslash\{0,1\}$, the system $C$ is solvable in integers $x_{1}, x_{2}, x_{3}, x_{5}, x_{6}, x_{7}, x_{8}, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$ greater than 1 if and only if

$$
\left(x_{4}+2=x_{9}\right) \wedge\left(x_{4} \mid\left(x_{4}-1\right)!+1\right) \wedge\left(x_{9} \mid\left(x_{9}-1\right)!+1\right)
$$

Hence, the claim of Lemma 11 follows from Lemma 5
Theorem 8. The statement $\Theta_{16}$ proves the following implication: if there exists a twin prime greater than $f(16)+3$, then there are infinitely many twin primes.

Proof. Assume the antecedent holds. Then, there exist prime numbers $x_{4}$ and $x_{9}$ such that $x_{9}=x_{4}+2>f(16)+3$. Hence, $x_{4} \in \mathbb{N} \backslash\{0,1,2,3\}$. By Lemma 11 , there exists a unique tuple $\left(x_{1}, x_{2}, x_{3}, x_{5}, x_{6}, x_{7}, x_{8}, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}\right) \in$ $(\mathbb{N} \backslash\{0,1\})^{14}$ such that the tuple $\left(x_{1}, \ldots, x_{16}\right)$ solves the system $C$. Lemma 11 guarantees that $\min \left(x_{1}, \ldots, x_{16}\right)=x_{1}=$ $x_{9}-3>f(16)$. Since $C \subseteq H_{16}$, the statement $\Theta_{16}$ and the inequality $\min \left(x_{1}, \ldots, x_{16}\right)>f(16)$ imply that the system $C$ has infinitely many solutions in integers $x_{1}, \ldots, x_{16}$ greater than 1. According to Lemmas 9,11 , there are infinitely many twin primes.

Hypothesis 4. The implication in Theorem 8 is true.
Corollary 8. (cf. [2]]). Assuming Hypothesis 4. a single query to an oracle for the halting problem decides the twin prime problem.
VII. Are there infinitely many composite Fermat numbers?

Primes of the form $2^{2^{n}}+1$ are called Fermat primes, as Fermat conjectured that every integer of the form $2^{2^{n}}+1$ is prime, see [8, p. 1]. Fermat correctly remarked that
$2^{2^{0}}+1=3,2^{2^{1}}+1=5,2^{2^{2}}+1=17,2^{2^{3}}+1=257$, and $2^{2^{4}}+1=65537$ are all prime, see [8, p. 1].
Open Problem. ([8, p. 159]). Are there infinitely many composite numbers of the form $2^{2^{n}}+1$ ?
Most mathematicians believe that $2^{2^{n}}+1$ is composite for every integer $n \geqslant 5$, see [7, p. 23].

Lemma 12. ([8 p. 38]). For every positive integer $n$, if a prime number $p$ divides $2^{2^{n}}+1$, then there exists a positive integer $k$ such that $p=k \cdot 2^{n+1}+1$.
Corollary 9. Since $k \cdot 2^{n+1}+1 \geqslant 2^{n+1}+1 \geqslant n+3$, for every positive integers $x, y$, and $n$, the equality $(x+1)(y+1)=2^{2^{n}}+1$ implies that $\min (n, x, x+1, y, y+1)=n$.

Let $g(1)=1$, and let $g(n+1)=2^{2^{g(n)}}$ for every positive integer $n$. Let

$$
\begin{aligned}
G_{n}= & \left\{x_{i} \cdot x_{j}=x_{k}: i, j, k \in\{1, \ldots, n\}\right\} \cup \\
& \left\{2^{2^{x_{i}}}=x_{k}: i, k \in\{1, \ldots, n\}\right\}
\end{aligned}
$$

Lemma 13. The following subsystem of $G_{n}$

$$
\left\{\begin{aligned}
x_{1} \cdot x_{1} & =x_{1} \\
\forall i \in\{1, \ldots, n-1\} 2^{2^{x_{i}}} & =x_{i+1}
\end{aligned}\right.
$$

has exactly one solution $\left(x_{1}, \ldots, x_{n}\right) \in(\mathbb{N} \backslash\{0\})^{n}$, namely $(g(1), \ldots, g(n))$.

For a positive integer $n$, let $\Psi_{n}$ denote the following statement: if a system $S \subseteq G_{n}$ has at most finitely many solutions in positive integers $x_{1}, \ldots, x_{n}$, then each such solution $\left(x_{1}, \ldots, x_{n}\right)$ satisfies $\min \left(x_{1}, \ldots, x_{n}\right) \leqslant g(n)$. The assumption $\min \left(x_{1}, \ldots, x_{n}\right) \leqslant g(n)$ is weaker than the assumption $\max \left(x_{1}, \ldots, x_{n}\right) \leqslant g(n)$ suggested by Lemma 13

Lemma 14. For every positive integer $n$, the system $G_{n}$ has a finite number of subsystems.

Theorem 9. Every statement $\Psi_{n}$ is true with an unknown integer bound that depends on $n$.

Proof. It follows from Lemma 14
Lemma 15. For every non-negative integers $b$ and $c, b+1=c$ if and only if $2^{2^{b}} \cdot 2^{2^{b}}=2^{2^{c}}$.

Theorem 10. The statement $\Psi_{13}$ proves the following implication: if $2^{2^{n}}+1$ is composite for some integer $n>g(13)$, then $2^{2^{n}}+1$ is composite for infinitely many positive integers $n$.
Proof. Let us consider the equation

$$
\begin{equation*}
(x+1)(y+1)=2^{2^{z}}+1 \tag{1}
\end{equation*}
$$

in positive integers. By Lemma 15 , we can transform equation (1) into an equivalent system $\mathcal{F}$ which has 13 variables
( $x, y, z$, and 10 other variables) and which consists of equations of the forms $\alpha \cdot \beta=\gamma$ and $2^{2^{\alpha}}=\gamma$, see the diagram in Figure 6.


Fig. 6 Construction of the system $\mathcal{F}$
Assume that $2^{2^{n}}+1$ is composite for some integer $n>g(13)$. By this and Corollary 9 , equation (1) has a solution $(x, y, z) \in(\mathbb{N} \backslash\{0\})^{3}$ such that $z=n$ and $z=\min (z, x, x+1, y, y+1)$. Hence, the system $\mathcal{F}$ has a solution in positive integers such that $z=n$ and $n$ is the smallest number in the solution sequence. Since $n>g(13)$, the statement $\Psi_{13}$ implies that the system $\mathcal{F}$ has infinitely many solutions in positive integers. Therefore, there are infinitely many positive integers $n$ such that $2^{2^{n}}+1$ is composite.

Hypothesis 5. The implication in Theorem 10 is true.
Corollary 10. Assuming Hypothesis 5 a single query to an oracle for the halting problem decides whether or not the set of composite Fermat numbers is infinite.
VIII. Computations of length $n$ and the statements $\Phi_{n}$

For a positive integer $x$, let $\Gamma(x)$ denote $(x-1)$ !. Let fact $^{-1}:\{1,2,6,24, \ldots\} \rightarrow \mathbb{N} \backslash\{0\}$ denote the inverse function to the factorial function. For positive integers $x$ and $y$, let $\operatorname{rem}(x, y)$ denote the remainder from dividing $x$ by $y$.
Definition. For a positive integer $n$, by a computation of length $n$ we understand any sequence of terms $x_{1}, \ldots, x_{n}$ such that $x_{1}$ is defined as the variable $x$, and for every integer $i \in\{2, \ldots, n\}, x_{i}$ is defined as $\Gamma\left(x_{i-1}\right)$, or fact $^{-1}\left(x_{i-1}\right)$, or $\operatorname{rem}\left(x_{i-1}, x_{i-2}\right)$ (only if $i \geqslant 3$ and $x_{i-1}$ is defined as $\Gamma\left(x_{i-2}\right)$ ).

For a positive integer $n$, let $c(n)$ denote the number of computations of length $n$. Then, $c(1)=1, c(2)=2$, and $c(n)=c(n-2)+2 \cdot c(n-1)$ for every integer $n \geqslant 3$. Hence, $c(3)=5, c(4)=12, c(5)=29, c(6)=70$, and $c(7)=169$.

Let $\mathcal{P}$ denote the set of prime numbers.
Lemma 16. ([]3] pp. 214-215]). For every positive integer $x$, $\operatorname{rem}(\Gamma(x), x) \in \mathbb{N} \backslash\{0\}$ if and only if $x \in\{4\} \cup \mathcal{P}$.

Let $h(4)=3$, and let $h(n+1)=h(n)$ ! for every integer $n \geqslant 4$.
Theorem 11. For every integer $n \geqslant 4$ and for every positive integer $x$, the following computation $\mathcal{H}_{n}$

$$
\left\{\begin{aligned}
x_{1} & :=x \\
\forall i \in\{2, \ldots, n-3\} x_{i} & :=\operatorname{fact}^{-1}\left(x_{i-1}\right) \\
x_{n-2} & :=\Gamma\left(x_{n-3}\right) \\
x_{n-1} & :=\Gamma\left(x_{n-2}\right) \\
x_{n} & :=\operatorname{rem}\left(x_{n-1}, x_{n-2}\right)
\end{aligned}\right.
$$

returns positive integers $x_{1}, \ldots, x_{n}$ if and only if $x=h(n)$.
Proof. We make three observations.
Observation 2. If $x_{n-3}=3$, then $x_{1}, \ldots, x_{n-3} \in \mathbb{N} \backslash\{0\}$ and $x=x_{1}=h(n)$. If $x=h(n)$, then $x_{1}, \ldots, x_{n-3} \in \mathbb{N} \backslash\{0\}$ and $x_{n-3}=3$. Hence, $x_{n-2}=\Gamma\left(x_{n-3}\right)=2$ and $x_{n-1}=\Gamma\left(x_{n-2}\right)=1$. Therefore, $x_{n}=\operatorname{rem}\left(x_{n-1}, x_{n-2}\right)=1$.

Observation 3. If $x_{n-3}=2$, then $x=x_{1}=\ldots=x_{n-3}=2$. If $x=2$, then $x_{1}=\ldots=x_{n-3}=2$. Hence, $x_{n-2}=\Gamma\left(x_{n-3}\right)=1$ and $x_{n-1}=\Gamma\left(x_{n-2}\right)=1$. Therefore, $x_{n}=\operatorname{rem}\left(x_{n-1}, x_{n-2}\right)=0 \notin$ $\mathbb{N} \backslash\{0\}$.

Observation 4. If $x_{n-3}=1$, then $x_{n-2}=\Gamma\left(x_{n-3}\right)=1$. Hence, $x_{n-1}=\Gamma\left(x_{n-2}\right)=1$. Therefore, $x_{n}=\operatorname{rem}\left(x_{n-1}, x_{n-2}\right)=0 \notin$ $\mathbb{N} \backslash\{0\}$.

Observations 2 4 cover the case when $x_{n-3} \in\{1,2,3\}$. If $x_{n-3} \geqslant 4$, then $x_{n-2}=\Gamma\left(x_{n-3}\right)$ is greater than 4 and composite. By Lemma 16, $x_{n}=\operatorname{rem}\left(x_{n-1}, x_{n-2}\right)=\operatorname{rem}\left(\Gamma\left(x_{n-2}\right), x_{n-2}\right)=$ $0 \notin \mathbb{N} \backslash\{0\}$.

For an integer $n \geqslant 4$, let $\Phi_{n}$ denote the following statement: if a computation of length $n$ returns positive integers $x_{1}, \ldots, x_{n}$ for at most finitely many positive integers $x$, then every such $x$ does not exceed $h(n)$.

Theorem 12．For every integer $n \geqslant 4$ ，the bound $h(n)$ in the statement $\Phi_{n}$ cannot be decreased．

Proof．It follows from Theorem 11
Lemma 17．For every positive integer $n$ ，there are only finitely many computations of length $n$ ．

Theorem 13．For every integer $n \geqslant 4$ ，the statement $\Phi_{n}$ is true with an unknown integer bound that depends on $n$ ．
Proof．It follows from Lemma 17.
IX．Consequences of the statements $\Phi_{4}, \ldots, \Phi_{7}$
Lemma 18．If $x \in \mathcal{P}$ ，then $\operatorname{rem}(\Gamma(x), x)=x-1$ ．
Proof．It follows from Lemma 5 ．
Lemma 19．For every positive integer $x$ ，the following com－ putation $\mathcal{T}$

$$
\left\{\begin{aligned}
x_{1} & :=x \\
x_{2} & :=\Gamma\left(x_{1}\right) \\
x_{3} & :=\operatorname{rem}^{\left(x_{2}, x_{1}\right)} \\
x_{4} & :=\operatorname{fact}^{-1}\left(x_{3}\right)
\end{aligned}\right.
$$

returns positive integers $x_{1}, \ldots, x_{4}$ if and only if $x=4$ or $x$ is a prime number of the form $n!+1$ ．

Proof．For an integer $i \in\{1, \ldots, 4\}$ ，let $T_{i}$ denote the set of positive integers $x$ such that the first $i$ instructions of the computation $\mathcal{T}$ returns positive integers $x_{1}, \ldots, x_{i}$ ．We show that

$$
\begin{equation*}
T_{4}=\{4\} \cup(\{n!+1: n \in \mathbb{N} \backslash\{0\}\} \cap \mathcal{P}) \tag{2}
\end{equation*}
$$

For every positive integer $x$ ，the terms $x_{1}$ and $x_{2}$ belong to $\mathbb{N} \backslash\{0\}$ ．By Lemma 16，the term $x_{3}$（which equals $\operatorname{rem}(\Gamma(x), x)$ ）belongs to $\overline{\mathbb{N}} \backslash\{0\}$ if and only if $x \in\{4\} \cup \mathcal{P}$ ． Hence，$T_{3}=\{4\} \cup \mathcal{P}$ ．If $x=4$ ，then $x_{1}, \ldots, x_{4} \in \mathbb{N} \backslash\{0\}$ ． Hence， $4 \in T_{4}$ ．If $x \in \mathcal{P}$ ，then Lemma 18 implies that $x_{3}=\operatorname{rem}(\Gamma(x), x)=x-1 \in \mathbb{N} \backslash\{0\}$ ．Therefore，for every $x \in \mathcal{P}$ ，the term $x_{4}=\operatorname{fact}^{-1}\left(x_{3}\right)$ belongs to $\mathbb{N} \backslash\{0\}$ if and only if $x \in\{n!+1: n \in \mathbb{N} \backslash\{0\}\}$ ．This proves equality（2）．

Theorem 14．The statement $\Phi_{4}$ implies that the set of primes of the form $n!+1$ is infinite．
Proof．The number $3!+1=7$ is prime．By Lemma 19，for $x=7$ the computation $\mathcal{T}$ returns positive integers $x_{1}, \ldots, x_{4}$ ． Since $x=7>3=h(4)$ ，the statement $\Phi_{4}$ guarantees that the computation $\mathcal{T}$ returns positive integers $x_{1}, \ldots, x_{4}$ for infinitely many positive integers $x$ ．By Lemma 19 there are infinitely many primes of the form $n!+1$ ．

Lemma 20．If $x \in \mathbb{N} \backslash\{0,1\}$ ，then fact $^{-1}(\Gamma(x))=x-1$ ．
Theorem 15．If the set of primes of the form $n!+1$ is infinite， then the statement $\Phi_{4}$ is true．

Proof．There exist exactly 10 computations of length 4 that differ from $\mathcal{H}_{4}$ and $\mathcal{T}$ ，see Table 1．For every such computa－ tion $\mathcal{F}_{i}$ ，we determine the set $S_{i}$ of all positive integers $x$ such that the computation $\mathcal{F}_{i}$ outputs positive integers $x_{1}, \ldots, x_{4}$ on
input $x$ ．We omit 10 easy proofs which use Lemmas 16 and 20 The sets $S_{i}$ are infinite，see Table 1.

| $0_{0}^{7}$ | ${ }_{6}{ }^{4}$ | $\infty^{4}$ | $\mathrm{c}^{4}$ | $0^{4}$ | 4 | U4 | ${ }_{+}^{+7}$ | $\omega^{4}$ | $\pm$ | $\mathrm{N}^{4}$ | － |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \dddot{x} \\ & \text { ii } \\ & \vdots \end{aligned}$ | $\begin{aligned} & \dddot{y} \\ & \text { ii } \\ & \underset{y}{2} \end{aligned}$ | $\begin{aligned} & x \\ & \text { ii } \\ & x \\ & x \end{aligned}$ | $\begin{aligned} & \ddot{x} \\ & \text { ii } \\ & \ddot{x} \end{aligned}$ | $\begin{aligned} & \Varangle \\ & \vdots \\ & i i \\ & \vdots \end{aligned}$ | $\begin{aligned} & \ddot{z} \\ & i i \\ & z \\ & \hline \end{aligned}$ | $\begin{aligned} & \underset{y}{x} \\ & \text { ii } \\ & \ddot{x} \end{aligned}$ | \＃ |  | $\begin{aligned} & \vdots \\ & \vdots \\ & \vdots \\ & \hdashline \end{aligned}$ | $\begin{aligned} & \varkappa \\ & \text { ii } \\ & \vdots \end{aligned}$ | $\begin{aligned} & \varkappa \\ & \text { ii } \\ & \underset{y}{n} \end{aligned}$ |
|  |  |  |  |  |  | $$ |  | $$ | $\begin{gathered} \text { N } \\ \text { ii } \\ \underset{y}{\Xi} \\ \hline \end{gathered}$ | $$ | $\begin{aligned} & \text { ふ } \\ & \text { ii } \\ & \underset{~}{¿} \end{aligned}$ |
|  | $\begin{aligned} & \dot{\omega} \\ & \text { ii } \\ & \stackrel{\rightharpoonup}{\circ} \\ & \underset{\sim}{\prime} \\ & \underset{\sim}{心} \end{aligned}$ | $\begin{aligned} & \dot{心} \\ & \text { ii } \\ & \underset{\sim}{\mathcal{N}} \\ & \hline \end{aligned}$ |  | $\begin{aligned} & \dot{\omega} \\ & \text { ii } \\ & \underset{\widetilde{\prime}}{心} \end{aligned}$ |  |  |  |  | $\begin{gathered} \stackrel{y}{\omega} \\ \text { ii } \\ \underset{\sim}{心} \\ \stackrel{y}{c} \end{gathered}$ | $$ | $\begin{aligned} & \text { w } \\ & \text { ii } \\ & \underset{\sim}{心} \\ & \text { in } \end{aligned}$ |
|  |  |  |  | $\begin{aligned} & \stackrel{y}{\stackrel{1}{2}} \\ & \text { ii } \\ & \underset{\sim}{\underset{心}{2}} \end{aligned}$ |  |  |  |  |  |  | $\begin{aligned} & \ddagger \\ & \text { ※i } \\ & \text { ii } \\ & \stackrel{\rightharpoonup}{\omega} \end{aligned}$ |
|  |  |  |  |  |  |  |  |  |  |  |  |

Tab． 112 computations of length $4, x \in \mathbb{N} \backslash\{0\}$
This completes the proof．
Hypothesis 6．The statements $\Phi_{4}, \ldots, \Phi_{7}$ are true．
Lemma 21．For every positive integer $x$ ，the following com－ putation $y$

$$
\left\{\begin{aligned}
x_{1} & :=x \\
x_{2} & :=\Gamma\left(x_{1}\right) \\
x_{3} & :=\operatorname{rem}^{2}\left(x_{2}, x_{1}\right) \\
x_{4} & :=\operatorname{fact}^{-1}\left(x_{3}\right) \\
x_{5} & :=\Gamma\left(x_{4}\right) \\
x_{6} & :=\operatorname{rem}\left(x_{5}, x_{4}\right)
\end{aligned}\right.
$$

returns positive integers $x_{1}, \ldots, x_{6}$ if and only if $x \in\{4\} \cup$ $\{p!+1: p \in \mathcal{P}\} \cap \mathcal{P}$.

Proof. For an integer $i \in\{1, \ldots, 6\}$, let $Y_{i}$ denote the set of positive integers $x$ such that the first $i$ instructions of the computation $y$ returns positive integers $x_{1}, \ldots, x_{i}$. Since the computations $\mathcal{T}$ and $\mathcal{Y}$ have the same first four instructions, the equality $Y_{i}=T_{i}$ holds for every $i \in\{1, \ldots, 4\}$. In particular,

$$
Y_{4}=\{4\} \cup(\{n!+1: n \in \mathbb{N} \backslash\{0\}\} \cap \mathcal{P})
$$

We show that

$$
\begin{equation*}
Y_{6}=\{4\} \cup(\{p!+1: p \in \mathcal{P}\} \cap \mathcal{P}) \tag{3}
\end{equation*}
$$

If $x=4$, then $x_{1}, \ldots, x_{6} \in \mathbb{N} \backslash\{0\}$. Hence, $4 \in Y_{6}$. Let $x \in \mathcal{P}$, and let $x=n!+1$, where $n \in \mathbb{N} \backslash\{0\}$. Hence, $n \neq 4$. Lemma 18 implies that $x_{3}=\operatorname{rem}(\Gamma(x), x)=x-1=n!$. Hence, $x_{4}=$ fact $^{-1}\left(x_{3}\right)=n$ and $x_{5}=\Gamma\left(x_{4}\right)=\Gamma(n) \in \mathbb{N} \backslash\{0\}$. By Lemma 16, the term $x_{6}$ (which equals rem $(\Gamma(n), n)$ ) belongs to $\mathbb{N} \backslash\{0\}$ if and only if $n \in\{4\} \cup \mathcal{P}$. This proves equality (3) as $n \neq 4$.

Theorem 16. The statement $\Phi_{6}$ implies that for infinitely many primes $p$ the number $p!+1$ is prime.

Proof. The numbers 11 and $11!+1$ are prime, see [1] p. 441] and [16]. By Lemma [21, for $x=11!+1$ the computation $y$ returns positive integers $x_{1}, \ldots, x_{6}$. Since $x=11!+1>6!=h(6)$, the statement $\Phi_{6}$ guarantees that the computation $\mathcal{Y}$ returns positive integers $x_{1}, \ldots, x_{6}$ for infinitely many positive integers $x$. By Lemma 21, for infinitely many primes $p$ the number $p!+1$ is prime.

Lemma 22. For every positive integer $x$, the following computation $\mathcal{L}$

$$
\left\{\begin{aligned}
x_{1} & :=x \\
x_{2} & :=\Gamma\left(x_{1}\right) \\
x_{3} & :=\Gamma\left(x_{2}\right) \\
x_{4} & :=\operatorname{fact}^{-1}\left(x_{3}\right) \\
x_{5} & :=\Gamma\left(x_{4}\right) \\
x_{6} & :=\operatorname{rem}\left(x_{5}, x_{4}\right)
\end{aligned}\right.
$$

returns positive integers $x_{1}, \ldots, x_{6}$ if and only if $(x-1)$ ! -1 is prime.

Proof. For an integer $i \in\{1, \ldots, 6\}$, let $L_{i}$ denote the set of positive integers $x$ such that the first $i$ instructions of the computation $\mathcal{L}$ returns positive integers $x_{1}, \ldots, x_{i}$. If $x \in\{1,2,3\}$, then $x_{6}=0$. Therefore, $L_{6} \subseteq \mathbb{N} \backslash\{0,1,2,3\}$. By Lemma 20 , for every integer $x \geqslant 4, x_{4}=(x-1)!-1, x_{5}=\Gamma((x-1)!-1)$, and $x_{1}, \ldots, x_{5} \in \mathbb{N} \backslash\{0\}$. By Lemma 16, for every integer $x \geqslant 4$,

$$
x_{6}=\operatorname{rem}(\Gamma((x-1)!-1),(x-1)!-1)
$$

belongs to $\mathbb{N} \backslash\{0\}$ if and only if $(x-1)!-1 \in\{4\} \cup \mathcal{P}$. The last condition equivalently expresses that $(x-1)!-1$ is prime as $(x-1)!-1 \geqslant 5$ for every integer $x \geqslant 4$. Hence,

$$
\begin{gathered}
L_{6}=(\mathbb{N} \backslash\{0,1,2,3\}) \cap\{x \in \mathbb{N} \backslash\{0,1,2,3\}:(x-1)!-1 \in \mathcal{P}\}= \\
\{x \in \mathbb{N} \backslash\{0\}:(x-1)!-1 \in \mathcal{P}\}
\end{gathered}
$$

It is conjectured that there are infinitely many primes of the form $n!-1$, see [1, p. 443] and [15].

Theorem 17. The statement $\Phi_{6}$ implies that there are infinitely many primes of the form $x!-1$.
Proof. The number $(975-1)!-1$ is prime, see [1] p. 441] and [15]. By Lemma 22, for $x=975$ the computation $\mathcal{L}$ returns positive integers $x_{1}, \ldots, x_{6}$. Since $x=975>720=h(6)$, the statement $\Phi_{6}$ guarantees that the computation $\mathcal{L}$ returns positive integers $x_{1}, \ldots, x_{6}$ for infinitely many positive integers $x$. By Lemma 22, the set $\{x \in \mathbb{N} \backslash\{0\}:(x-1)!-1 \in \mathcal{P}\}$ is infinite.

Lemma 23. For every positive integer $x$, the following computation $\mathcal{D}$

$$
\left\{\begin{aligned}
x_{1} & :=x \\
x_{2} & :=\Gamma\left(x_{1}\right) \\
x_{3} & :=\operatorname{rem}\left(x_{2}, x_{1}\right) \\
x_{4} & :=\Gamma\left(x_{3}\right) \\
x_{5} & :=\operatorname{fact}^{-1}\left(x_{4}\right) \\
x_{6} & :=\Gamma\left(x_{5}\right) \\
x_{7} & :=\operatorname{rem}\left(x_{6}, x_{5}\right)
\end{aligned}\right.
$$

returns positive integers $x_{1}, \ldots, x_{7}$ if and only if both $x$ and $x-2$ are prime.

Proof. For an integer $i \in\{1, \ldots, 7\}$, let $D_{i}$ denote the set of positive integers $x$ such that the first $i$ instructions of the computation $\mathcal{D}$ returns positive integers $x_{1}, \ldots, x_{i}$. If $x=1$, then $x_{3}=0$. Hence, $D_{7} \subseteq D_{3} \subseteq \mathbb{N} \backslash\{0,1\}$. If $x \in\{2,3,4\}$, then $x_{7}=0$. Therefore,

$$
D_{7} \subseteq(\mathbb{N} \backslash\{0,1\}) \cap(\mathbb{N} \backslash\{0,2,3,4\})=\mathbb{N} \backslash\{0,1,2,3,4\}
$$

By Lemma 16 for every integer $x \geqslant 5$, the term $x_{3}$ (which equals $\operatorname{rem}(\Gamma(x), x)$ ) belongs to $\mathbb{N} \backslash\{0\}$ if and only if $x \in \mathcal{P} \backslash\{2,3\}$. By Lemma 18 , for every $x \in \mathcal{P} \backslash\{2,3\}$, $x_{3}=x-1 \in \mathbb{N} \backslash\{0,1,2,3\}$. By Lemma 20. for every $x \in \mathcal{P} \backslash\{2,3\}$, the terms $x_{4}$ and $x_{5}$ belong to $\mathbb{N} \backslash\{0\}$ and $x_{5}=x_{3}-1=x-2$. By Lemma 16, for every $x \in \mathcal{P} \backslash\{2,3\}$, the term $x_{7}$ (which equals rem $\left(\Gamma\left(x_{5}\right), x_{5}\right)$ ) belongs to $\mathbb{N} \backslash\{0\}$ if and only if $x_{5}=x-2 \in\{4\} \cup \mathcal{P}$. From these facts, we obtain that

$$
\begin{gathered}
D_{7}=(\mathbb{N} \backslash\{0,1,2,3,4\}) \cap(\mathcal{P} \backslash\{2,3\}) \cap(\{6\} \cup\{p+2: p \in \mathcal{P}\})= \\
\{p \in \mathcal{P}: p-2 \in \mathcal{P}\}
\end{gathered}
$$

Theorem 18. The statement $\Phi_{7}$ implies that there are infinitely many twin primes.
Proof. Harvey Dubner proved that the numbers $459 \cdot 2^{8529}-1$ and $459 \cdot 2^{8529}+1$ are prime, see [17] p. 87]. By Lemma 23, for $x=459 \cdot 2^{8529}+1$ the computation $\mathcal{D}$ returns positive integers $x_{1}, \ldots, x_{7}$. Since $x>720!=h(7)$, the statement $\Phi_{7}$ guarantees that the computation $\mathcal{D}$ returns positive integers $x_{1}, \ldots, x_{7}$ for infinitely many positive integers $x$. By Lemma 23, there are infinitely many twin primes.

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