

The physical limits of computation inspire an open problem that concerns decidable sets $X \subseteq \mathbb{N}$ and cannot be formalized in mathematics understood as an a priori science because it refers to the current knowledge on X

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ABSTRACT. Algorithms always terminate. We explain the distinction between existing algorithms (i.e. algorithms whose existence is provable in *ZFC*) and known algorithms (i.e. algorithms whose definition is constructive and currently known). Assuming that the infiniteness of a set $X \subseteq \mathbb{N}$ is false or unproven, we define which elements of X are classified as known. No known set $X \subseteq \mathbb{N}$ satisfies Conditions (1)-(4) and is widely known in number theory or naturally defined, where this term has only informal meaning. (1) A known algorithm with no input returns an integer n satisfying $\text{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$. (2) A known algorithm for every $k \in \mathbb{N}$ decides whether or not $k \in X$. (3) No known algorithm with no input returns the logical value of the statement $\text{card}(X) = \omega$. (4) There are many elements of X and it is conjectured, though so far unproven, that X is infinite. (5) X is naturally defined. The infiniteness of X is false or unproven. X has the simplest definition among known sets $Y \subseteq \mathbb{N}$ with the same set of known elements. Let \mathcal{P}_{n^2+1} denote the set of primes of the form $n^2 + 1$. Conditions (2)-(5) hold for $X = \mathcal{P}_{n^2+1}$. We discuss a conjecture which implies the conjunction of Conditions (1)-(5) for $X = \mathcal{P}_{n^2+1}$. No set $X \subseteq \mathbb{N}$ will satisfy Conditions (1)-(4) forever, if for every algorithm with no input, at some future day, a computer will be able to execute this algorithm in 1 second or less. The physical limits of computation disprove this assumption. We present a table that shows satisfiable conjunctions consisting of Conditions (1)-(5) and their negations.

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1. DEFINITIONS AND THE DISTINCTION BETWEEN EXISTING ALGORITHMS AND KNOWN ALGORITHMS

Algorithms always terminate. Semi-algorithms may not terminate. Examples 1–4 and the proof of Statement 1 explain the distinction between *existing algorithms* (i.e. algorithms whose existence is provable in *ZFC*) and *known algorithms* (i.e. algorithms whose definition is constructive and currently known). A definition of an integer n is called *constructive*, if it provides a known algorithm with no input that returns n . Definition 1 applies to sets $X \subseteq \mathbb{N}$ whose infiniteness is false or unproven.

Definition 1. We say that a non-negative integer k is a known element of X , if $k \in X$ and we know an algebraic expression that defines k and consists of the following signs: 1 (one), + (addition), - (subtraction), \cdot (multiplication), / (division), \wedge (exponentiation), ! (factorial), ((left parenthesis),) (right parenthesis).

Let t denote the largest twin prime that is smaller than (((((((((9!)!)!)!)!)!)!)!)!. The number t is an unknown element of the set of twin primes.

Lemma 1. (*Wilson's theorem*, [2, p. 89]). *For every integer $x \geq 2$, x is prime if and only if x divides $(x - 1)! + 1$.*

Edmund Landau's conjecture states that the set \mathcal{P}_{n^2+1} of primes of the form $n^2 + 1$ is infinite, see [10]–[13]. Let $[\cdot]$ denote the integer part function. By Lemma 1, for every positive integer n ,

$$1 + 1 + \left((n^2)! - (n^2 + 1) \cdot \left[\frac{(n^2)!}{n^2 + 1} \right] \right) \cdot \frac{n^2 - 1}{n^2} = \begin{cases} n^2 + 1, & \text{if } n^2 + 1 \text{ is prime} \\ 2, & \text{otherwise} \end{cases}$$

Similar identities are unknown for algebraic expressions considered in Definition 1, so Definition 1 seems to be correct.

Definition 2. *Conditions (1)–(5) concern sets $X \subseteq \mathbb{N}$.*

(1) *A known algorithm with no input returns an integer n satisfying $\text{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$.*

(2) *A known algorithm for every $k \in \mathbb{N}$ decides whether or not $k \in X$.*

(3) *No known algorithm with no input returns the logical value of the statement $\text{card}(X) = \omega$.*

(4) *There are many elements of X and it is conjectured, though so far unproven, that X is infinite.*

(5) *X is naturally defined. The infiniteness of X is false or unproven. X has the simplest definition among known sets $\mathcal{Y} \subseteq \mathbb{N}$ with the same set of known elements.*

Condition (3) implies that no known proof shows the finiteness/infiniteness of X . No known set $X \subseteq \mathbb{N}$ satisfies Conditions (1)–(4) and is widely known in number theory or naturally defined, where this term has only informal meaning.

Let $\beta = (((24!)!)!)!$.

Lemma 2. $\log_2(\log_2(\log_2(\log_2(\log_2(\log_2(\log_2(\beta))))))) \approx 1.42298$.

Proof. We ask Wolfram Alpha at <https://wolframalpha.com>. □

Example 1. *The set $X = \mathcal{P}_{n^2+1}$ satisfies Condition (3).*

Example 2. *The set $X = \begin{cases} \mathbb{N}, & \text{if } [\frac{\beta}{\pi}] \text{ is odd} \\ \emptyset, & \text{otherwise} \end{cases}$ does not satisfy Condition (3) because we know an algorithm with no input that computes $[\frac{\beta}{\pi}]$. The set of known elements of X is empty. Hence, Condition (5) fails for X .*

Example 3. ([1], [7], [9, p. 9]). *The function*

$$\mathbb{N} \ni n \xrightarrow{h} \begin{cases} 1, & \text{if the decimal expansion of } \pi \text{ contains } n \text{ consecutive zeros} \\ 0, & \text{otherwise} \end{cases}$$

is computable because $h = \mathbb{N} \times \{1\}$ or there exists $k \in \mathbb{N}$ such that

$$h = (\{0, \dots, k\} \times \{1\}) \cup (\{k + 1, k + 2, k + 3, \dots\} \times \{0\})$$

No known algorithm computes the function h .

Example 4. *The set*

$$\mathcal{X} = \begin{cases} \mathbb{N}, & \text{if the continuum hypothesis holds} \\ \emptyset, & \text{otherwise} \end{cases}$$

is decidable. This \mathcal{X} satisfies Conditions (1) and (3) and does not satisfy Conditions (2), (4), and (5). These facts will hold forever.

Let Φ denote the following unproven statement:

$$\text{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq [2, \beta]$$

Landau's conjecture implies the statement Φ . Theorem 6 heuristically justifies the statement Φ . This justification does not yield the finiteness/infiniteness of \mathcal{P}_{n^2+1} .

Statement 1. *Condition (1) remains unproven for $\mathcal{X} = \mathcal{P}_{n^2+1}$.*

Proof. For every set $X \subseteq \mathbb{N}$, there exists an algorithm $\text{Alg}(X)$ with no input that returns

$$n = \begin{cases} 0, & \text{if } \text{card}(X) \in \{0, \omega\} \\ \max(X), & \text{otherwise} \end{cases}$$

This n satisfies the implication in Condition (1), but the algorithm $\text{Alg}(\mathcal{P}_{n^2+1})$ is unknown because its definition is ineffective. \square

Proving the statement Φ will disprove Statement 1. Statement 1 cannot be formalized in mathematics understood as an a priori science because it refers to the current mathematical knowledge. The same is true for Open Problems 1–3 and Statements 2–4.

Definition 3. *We say that an integer n is a threshold number of a set $X \subseteq \mathbb{N}$, if $\text{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$.*

If a set $X \subseteq \mathbb{N}$ is empty or infinite, then any integer n is a threshold number of X . If a set $X \subseteq \mathbb{N}$ is non-empty and finite, then the all threshold numbers of X form the set $[\max(X), \infty) \cap \mathbb{N}$.

2. THE PHYSICAL LIMITS OF COMPUTATION INSPIRE OPEN PROBLEM 1

Let $f(1) = 2$, $f(2) = 4$, and let $f(n+1) = f(n)!$ for every integer $n \geq 2$.

Statement 2. *The set*

$$\mathcal{X} = \{k \in \mathbb{N} : (10^6 < k) \Rightarrow (f(10^6), f(k)) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$$

satisfies Conditions (1)–(4). Condition (5) fails for \mathcal{X} .

Proof. Condition (4) holds as $\mathcal{X} \supseteq \{0, \dots, 10^6\}$ and the set \mathcal{P}_{n^2+1} is conjecturally infinite. By Lemma 2, due to known physics we are not able to confirm by a direct computation that some element of \mathcal{P}_{n^2+1} is greater than $f(10^6) > f(7) = \beta$, see [5]. Thus Condition (3) holds. Condition (2) holds trivially. Since the set

$$\{k \in \mathbb{N} : (10^6 < k) \wedge (f(10^6), f(k)) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$$

is empty or infinite, the integer 10^6 is a threshold number of \mathcal{X} . Thus \mathcal{X} satisfies Condition (1). Condition (5) fails for \mathcal{X} as the set of known elements of \mathcal{X} equals $\{0, \dots, 10^6\}$. \square

For a positive integer n , let $\theta(n)$ denote the largest divisor of $10^{10^{10}}$ not greater than n .

Lemma 3. For every integer $j \geq 10^{10^{10}}$, $\theta(j) = 10^{10^{10}}$.

Lemma 4. For every integer $j \in [6553600, 7812500)$, $\theta(j) = 6553600$.

Proof. The integers $6553600 = 2^{18} \cdot 5^2$ and $7812500 = 2^2 \cdot 5^9$ divide $10^{10^{10}}$. The following MuPAD code

```
A:=[2^a $a=0..floor(log(2,781250000))]:
B:=[5^b $b=0..floor(log(5,781250000))]:
C:={}:
for a from 1 to nops(A) do
for b from 1 to nops(B) do
C:=C union {A[a]*B[b]}:
end_for:
end_for:
C intersect {$ 6553600..7812500};
```

returns the set $\{6553600, 7812500\}$. Therefore, every integer $j \in (6553600, 7812500)$ does not divide $10^{10^{10}}$, which completes the proof. \square

Lemma 5. The function

$$\mathbb{N} \setminus \{0\} \ni n \xrightarrow{\kappa} \text{the_exponent_of_2_in_the_prime_factorization_of_}n \in \mathbb{N}$$

takes every non-negative integer value infinitely often.

Lemma 6. $\{2k+1 : k \in \mathbb{N}\} \cap [6553600, 7812500) \subseteq$

$$\{n \in \mathbb{N} \setminus \{0\} : (\theta(n) + \kappa(n))^2 + 1 \text{ is prime}\}$$

Proof. The following PARI/GP command

```
isprime(6553600^2+1, {flag=2})
```

returns %1 = 1. This command performs the APRCL primality test, the best deterministic primality test algorithm ([12, p. 226]). It rigorously shows that the number $6553600^2 + 1$ is prime. By Lemma 4, for every odd integer $j \in [6553600, 7812500)$, the number $(\theta(j) + \kappa(j))^2 + 1 = (6553600 + 0)^2 + 1$ is prime. \square

Statement 3. Conditions (1)–(5) hold for $\mathcal{X} = \{n \in \mathbb{N} \setminus \{0\} : (\theta(n) + \kappa(n))^2 + 1 \text{ is prime}\}$ except the requirement that \mathcal{X} is naturally defined.

Proof. Condition (2) holds trivially. Let δ denote $10^{10^{10}}$. By Lemmas 3 and 5, Condition (1) holds for $n = \delta - 1$. Since the statement $\mathcal{P}_{n^2+1} \cap [\delta^2 + 1, \infty) \neq \emptyset$ remains unproven, Condition (3) holds. Lemma 6 and the implication

$$\mathcal{P}_{n^2+1} \cap [\delta^2 + 1, \infty) \neq \emptyset \implies \text{card}(\mathcal{X}) = \omega$$

show that Condition (4) holds. Since there exists an integer $t \geq 10^{10}$ such that $t^2 + 1$ is prime, we cannot simplify the definition of \mathcal{X} by defining the function θ with the number 10^{10} instead of the number $10^{10^{10}}$. Thus Condition 5 holds except the requirement that \mathcal{X} is naturally defined. \square

Proving Landau's conjecture will disprove Statements 2 and 3.

Open Problem 1. Is there a set $\mathcal{X} \subseteq \mathbb{N}$ which satisfies Conditions (1)–(5)?

Theorem 1. *No set $\mathcal{X} \subseteq \mathbb{N}$ will satisfy Conditions (1)–(4) forever, if for every algorithm with no input, at some future day, a computer will be able to execute this algorithm in 1 second or less.*

Proof. The proof goes by contradiction. We fix an integer n that satisfies Condition (1). Since Conditions (1)–(3) will hold forever, the semi-algorithm in Figure 1 never terminates and sequentially prints the following sentences:

$$(T) \quad n + 1 \notin \mathcal{X}, n + 2 \notin \mathcal{X}, n + 3 \notin \mathcal{X}, \dots$$

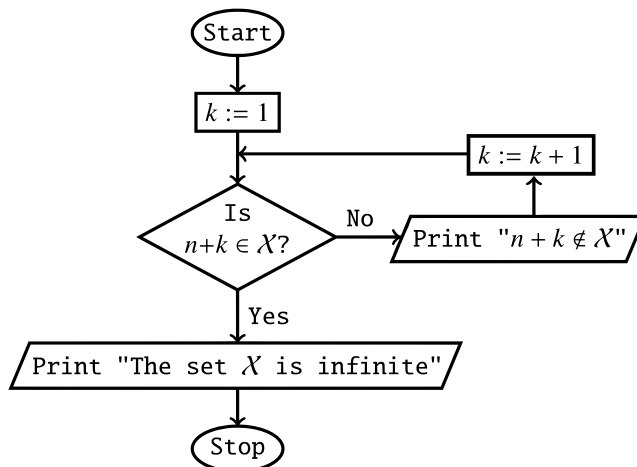


Fig. 1 Semi-algorithm that terminates if and only if \mathcal{X} is infinite

The sentences from the sequence (T) and our assumption imply that for every integer $m > n$ computed by a known algorithm, at some future day, a computer will be able to confirm in 1 second or less that $(n, m] \cap \mathcal{X} = \emptyset$. Thus, at some future day, numerical evidence will support the conjecture that the set \mathcal{X} is finite, contrary to the conjecture in Condition (4). \square

The physical limits of computation ([5]) disprove the assumption of Theorem 1.

3. NUMBER-THEORETIC STATEMENTS Ψ_n

Let \mathcal{U}_1 denote the system of equations which consists of the equation $x_1! = x_1$. For an integer $n \geq 2$, let \mathcal{U}_n denote the following system of equations:

$$\begin{cases} x_1! = x_1 \\ x_1 \cdot x_1 = x_2 \\ \forall i \in \{2, \dots, n-1\} x_i! = x_{i+1} \end{cases}$$

Lemma 7. *For every positive integer n , the system \mathcal{U}_n has exactly two solutions in positive integers x_1, \dots, x_n , namely $(1, \dots, 1)$ and $(f(1), \dots, f(n))$.*

Let B_n denote the following system of equations:

$$\{x_i! = x_k : i, k \in \{1, \dots, n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}\}$$

For every positive integer n , no known system $\mathcal{S} \subseteq B_n$ with a finite number of solutions in positive integers x_1, \dots, x_n has a solution $(x_1, \dots, x_n) \in (\mathbb{N} \setminus \{0\})^n$ satisfying $\max(x_1, \dots, x_n) > f(n)$. For every positive integer n and for every known system $\mathcal{S} \subseteq B_n$, if the finiteness/infiniteness of the set

$$\{(x_1, \dots, x_n) \in (\mathbb{N} \setminus \{0\})^n : (x_1, \dots, x_n) \text{ solves } \mathcal{S}\}$$

is unknown, then the statement

$$\exists x_1, \dots, x_n \in \mathbb{N} \setminus \{0\} ((x_1, \dots, x_n) \text{ solves } \mathcal{S}) \wedge (\max(x_1, \dots, x_n) > f(n))$$

remains unproven.

For a positive integer n , let Ψ_n denote the following statement: *if a system $\mathcal{S} \subseteq B_n$ has at most finitely many solutions in positive integers x_1, \dots, x_n , then each such solution (x_1, \dots, x_n) satisfies $x_1, \dots, x_n \leq f(n)$* . The statement Ψ_n says that for subsystems of B_n with a finite number of solutions, the largest known solution is indeed the largest possible. The statements Ψ_1 and Ψ_2 hold trivially. There is no reason to assume the validity of the statement $\forall n \in \mathbb{N} \setminus \{0\} \Psi_n$.

Theorem 2. *For every statement Ψ_n , the bound $f(n)$ cannot be decreased.*

Proof. It follows from Lemma 7 because $\mathcal{U}_n \subseteq B_n$. □

Theorem 3. *For every integer $n \geq 2$, the statement Ψ_{n+1} implies the statement Ψ_n .*

Proof. If a system $\mathcal{S} \subseteq B_n$ has at most finitely many solutions in positive integers x_1, \dots, x_n , then for every integer $i \in \{1, \dots, n\}$ the system $\mathcal{S} \cup \{x_i! = x_{n+1}\}$ has at most finitely many solutions in positive integers x_1, \dots, x_{n+1} . The statement Ψ_{n+1} implies that $x_i! = x_{n+1} \leq f(n+1) = f(n)!$. Hence, $x_i \leq f(n)$. □

Theorem 4. *Every statement Ψ_n is true with an unknown integer bound that depends on n .*

Proof. For every positive integer n , the system B_n has a finite number of subsystems. □

4. A CONJECTURAL SOLUTION OF OPEN PROBLEM 1

Lemma 8. *For every positive integers x and y , $x! \cdot y = y!$ if and only if*

$$(x + 1 = y) \vee (x = y = 1)$$

Let \mathcal{A} denote the following system of equations:

$$\left\{ \begin{array}{l} x_2! = x_3 \\ x_3! = x_4 \\ x_5! = x_6 \\ x_8! = x_9 \\ x_1 \cdot x_1 = x_2 \\ x_3 \cdot x_5 = x_6 \\ x_4 \cdot x_8 = x_9 \\ x_5 \cdot x_7 = x_8 \end{array} \right.$$

Lemma 8 and the diagram in Figure 2 explain the construction of the system \mathcal{A} .

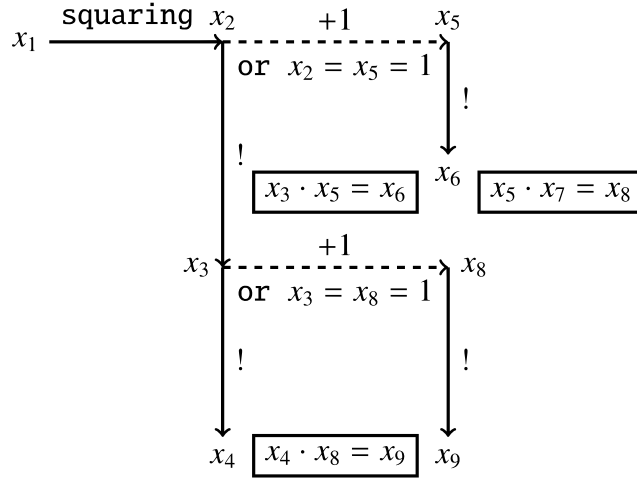


Fig. 2 Construction of the system \mathcal{A}

Lemma 9. For every integer $x_1 \geq 2$, the system \mathcal{A} is solvable in positive integers x_2, \dots, x_9 if and only if $x_1^2 + 1$ is prime. In this case, the integers x_2, \dots, x_9 are uniquely determined by the following equalities:

$$\begin{aligned}
 x_2 &= x_1^2 \\
 x_3 &= (x_1^2)! \\
 x_4 &= ((x_1^2)!)! \\
 x_5 &= x_1^2 + 1 \\
 x_6 &= (x_1^2 + 1)! \\
 x_7 &= \frac{(x_1^2)! + 1}{x_1^2 + 1} \\
 x_8 &= (x_1^2)! + 1 \\
 x_9 &= ((x_1^2)! + 1)!
 \end{aligned}$$

Proof. By Lemma 8, for every integer $x_1 \geq 2$, the system \mathcal{A} is solvable in positive integers x_2, \dots, x_9 if and only if $x_1^2 + 1$ divides $(x_1^2)! + 1$. Hence, the claim of Lemma 9 follows from Lemma 1. \square

Lemma 10. There are only finitely many tuples $(x_1, \dots, x_9) \in (\mathbb{N} \setminus \{0\})^9$, which solve the system \mathcal{A} and satisfy $x_1 = 1$. It is true as every such tuple (x_1, \dots, x_9) satisfies $x_1, \dots, x_9 \in \{1, 2\}$.

Proof. The equality $x_1 = 1$ implies that $x_2 = x_1 \cdot x_1 = 1$. Hence, $x_3 = x_2! = 1$. Therefore, $x_4 = x_3! = 1$. The equalities $x_5! = x_6$ and $x_5 = 1 \cdot x_5 = x_3 \cdot x_5 = x_6$ imply that $x_5, x_6 \in \{1, 2\}$. The equalities $x_8! = x_9$ and $x_8 = 1 \cdot x_8 = x_4 \cdot x_8 = x_9$ imply that $x_8, x_9 \in \{1, 2\}$. The equality $x_5 \cdot x_7 = x_8$ implies that $x_7 = \frac{x_8}{x_5} \in \left\{\frac{1}{1}, \frac{1}{2}, \frac{2}{1}, \frac{2}{2}\right\} \cap (\mathbb{N} \setminus \{0\}) = \{1, 2\}$. \square

Conjecture 1. The statement Ψ_9 is true when is restricted to the system \mathcal{A} .

Theorem 5. *Conjecture 1 proves the following implication: if there exists an integer $x_1 \geq 2$ such that $x_1^2 + 1$ is prime and greater than $f(7)$, then the set \mathcal{P}_{n^2+1} is infinite.*

Proof. Suppose that the antecedent holds. By Lemma 9, there exists a unique tuple $(x_2, \dots, x_9) \in (\mathbb{N} \setminus \{0\})^8$ such that the tuple (x_1, x_2, \dots, x_9) solves the system \mathcal{A} . Since $x_1^2 + 1 > f(7)$, we obtain that $x_1^2 \geq f(7)$. Hence, $(x_1^2)! \geq f(7)! = f(8)$. Consequently,

$$x_9 = ((x_1^2)! + 1)! \geq (f(8) + 1)! > f(8)! = f(9)$$

Conjecture 1 and the inequality $x_9 > f(9)$ imply that the system \mathcal{A} has infinitely many solutions $(x_1, \dots, x_9) \in (\mathbb{N} \setminus \{0\})^9$. According to Lemmas 9 and 10, the set \mathcal{P}_{n^2+1} is infinite. \square

Theorem 6. *Conjecture 1 implies the statement Φ .*

Proof. It follows from Theorem 5 and the equality $f(7) = (((24!)!)!)!$. \square

Theorem 7. *The statement Φ implies Conjecture 1.*

Proof. By Lemmas 9 and 10, if positive integers x_1, \dots, x_9 solve the system \mathcal{A} , then

$$(x_1 \geq 2) \wedge (x_5 = x_1^2 + 1) \wedge (x_5 \text{ is prime})$$

or $x_1, \dots, x_9 \in \{1, 2\}$. In the first case, Lemma 9 and the statement Φ imply that the inequality $x_5 \leq (((24!)!)!)! = f(7)$ holds when the system \mathcal{A} has at most finitely many solutions in positive integers x_1, \dots, x_9 . Hence, $x_2 = x_5 - 1 < f(7)$ and $x_3 = x_2! < f(7)! = f(8)$. Continuing this reasoning in the same manner, we can show that every x_i does not exceed $f(9)$. \square

Statement 4. *Conditions (2)–(5) hold for $X = \mathcal{P}_{n^2+1}$. The statement Φ implies that Condition (1) holds for $X = \mathcal{P}_{n^2+1}$.*

Proof. The set \mathcal{P}_{n^2+1} is conjecturally infinite. There are 2199894223892 primes of the form $n^2 + 1$ in the interval $[2, 10^{28})$, see [11]. These two facts imply Condition (4). By Lemma 2, due to known physics we are not able to confirm by a direct computation that some element of \mathcal{P}_{n^2+1} is greater than $f(7) = (((24!)!)!)! = \beta$, see [5]. Thus Condition (3) holds. Conditions (2) and (5) hold trivially. The statement Φ implies that Condition (1) holds for $X = \mathcal{P}_{n^2+1}$ with $n = \beta = (((24!)!)!)!$. \square

Proving Landau's conjecture will disprove Statement 4.

Conjecture 2. *Conditions (1)–(5) hold for $X = \mathcal{P}_{n^2+1} \wedge \Phi$.*

Conjecture 2 implies that every known proof of the statement Φ does not yield the finiteness/infiniteness of \mathcal{P}_{n^2+1} .

5. A NEW HEURISTIC ARGUMENT FOR THE INFINITENESS OF \mathcal{P}_{n^2+1}

The system \mathcal{A} contains four factorials and four multiplications. Let \mathcal{F} denote the family of all systems $\mathcal{S} \subseteq B_9$ which contain at most four factorials and at most four multiplications.

Among known systems $\mathcal{S} \in \mathcal{F}$, the following system \mathcal{C}

$$\left\{ \begin{array}{l} x_1! = x_2 \\ x_2 \cdot x_9 = x_1 \\ x_2 \cdot x_2 = x_3 \\ x_3 \cdot x_3 = x_4 \\ x_4 \cdot x_4 = x_5 \\ x_5! = x_6 \\ x_6! = x_7 \\ x_7! = x_8 \end{array} \right.$$

attains the greatest solution in positive integers x_1, \dots, x_9 and has at most finitely many solutions in $(\mathbb{N} \setminus \{0\})^9$. Only the tuples $(1, \dots, 1)$ and $(2, 2, 4, 16, 256, 256!, (256!)!, ((256!)!)!, 1)$ solve \mathcal{C} and belong to $(\mathbb{N} \setminus \{0\})^9$.

For every known system $\mathcal{S} \in \mathcal{F}$, if the finiteness of the set

$$\{(x_1, \dots, x_9) \in (\mathbb{N} \setminus \{0\})^9 : (x_1, \dots, x_9) \text{ solves } \mathcal{S}\}$$

is unproven and conjectured, then the statement

$$\exists x_1, \dots, x_9 \in \mathbb{N} \setminus \{0\} ((x_1, \dots, x_9) \text{ solves } \mathcal{S}) \wedge (\max(x_1, \dots, x_9) > ((256!)!)!)$$

remains unproven.

Let Γ denote the statement: *if the system \mathcal{A} has at most finitely many solutions in positive integers x_1, \dots, x_9 , then each such solution (x_1, \dots, x_9) satisfies $x_1, \dots, x_9 \leq ((256!)!)!$. The number $46^{512} + 1$ is prime ([6]) and greater than $256!$, see also [8, p. 239] for the primality of $150^{2048} + 1$. Hence, the statement Γ is equivalent to the infiniteness of \mathcal{P}_{n^2+1} . It heuristically justifies the infiniteness of \mathcal{P}_{n^2+1} in a sophisticated way.*

6. SATISFIABLE CONJUNCTIONS WHICH CONSIST OF CONDITIONS 1–5 AND THEIR NEGATIONS

The set $\mathcal{X} = \mathcal{P}_{n^2+1}$ satisfies the conjunction

$$\neg(\text{Condition 1}) \wedge (\text{Condition 2}) \wedge (\text{Condition 3}) \wedge (\text{Condition 4}) \wedge (\text{Condition 5})$$

The set $\mathcal{X} = \{0, \dots, f(7)\} \cup \mathcal{P}_{n^2+1}$ satisfies the conjunction

$$\neg(\text{Condition 1}) \wedge (\text{Condition 2}) \wedge (\text{Condition 3}) \wedge (\text{Condition 4}) \wedge \neg(\text{Condition 5})$$

The set $\mathcal{X} = \begin{cases} \mathbb{N}, & \text{if } (f(9^8), f(9^9)) \cap \mathcal{P}_{n^2+1} \neq \emptyset \\ \{0, \dots, 10^6\}, & \text{otherwise} \end{cases}$ satisfies the conjunction

$$(\text{Condition 1}) \wedge (\text{Condition 2}) \wedge \neg(\text{Condition 3}) \wedge (\text{Condition 4}) \wedge \neg(\text{Condition 5})$$

Open Problem 2. *Is there a set $\mathcal{X} \subseteq \mathbb{N}$ that satisfies the conjunction*

$$(\text{Condition 1}) \wedge (\text{Condition 2}) \wedge \neg(\text{Condition 3}) \wedge (\text{Condition 4}) \wedge (\text{Condition 5})?$$

The numbers $2^{2^k} + 1$ are prime for $k \in \{0, 1, 2, 3, 4\}$. It is open whether or not there are infinitely many primes of the form $2^{2^k} + 1$, see [4, p. 158] and [8, p. 74]. It is open whether or not there are infinitely many composite numbers of the form $2^{2^k} + 1$, see [4, p. 159] and [8, p. 74]. Most mathematicians believe that $2^{2^k} + 1$ is composite for every integer $k \geq 5$, see [3, p. 23].

The set

$$\mathcal{X} = \begin{cases} \mathbb{N}, & \text{if } (f(9^8), f(9^9)) \cap \mathcal{P}_{n^2+1} \neq \emptyset \\ \{0, \dots, 10^6\} \cup \{n \in \mathbb{N} : n \text{ is the sixth prime number of the form } 2^{2^k} + 1\}, & \text{otherwise} \end{cases}$$

satisfies the conjunction

$$\neg(\text{Condition 1}) \wedge (\text{Condition 2}) \wedge \neg(\text{Condition 3}) \wedge (\text{Condition 4}) \wedge \neg(\text{Condition 5})$$

Open Problem 3. *Is there a set $\mathcal{X} \subseteq \mathbb{N}$ that satisfies the conjunction*

$$\neg(\text{Condition 1}) \wedge (\text{Condition 2}) \wedge \neg(\text{Condition 3}) \wedge (\text{Condition 4}) \wedge (\text{Condition 5})?$$

It is possible, although very doubtful, that at some future day, the set $\mathcal{X} = \mathcal{P}_{n^2+1}$ will solve Open Problem 2. The same is true for Open Problem 3. It is possible, although very doubtful, that at some future day, the set $\mathcal{X} = \{k \in \mathbb{N} : 2^{2^k} + 1 \text{ is composite}\}$ will solve Open Problem 1. The same is true for Open Problems 2 and 3.

The following table shows satisfiable conjunctions consisting of Conditions (1)–(5) and their negations.

| | | |
|--|---|--|
| | (Condition 2) \wedge (Condition 3) \wedge (Condition 4) | (Condition 2) \wedge \neg (Condition 3) \wedge (Condition 4) |
| (Condition 1) \wedge (Condition 5) | Open Problem 1 (conjecturally solved with $\mathcal{X} = \mathcal{P}_{n^2+1}$) | Open Problem 2 |
| (Condition 1) \wedge \neg (Condition 5) | $\mathcal{X} = \{k \in \mathbb{N} : (10^6 < k) \Rightarrow (f(10^6), f(k)) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$ | $\mathcal{X} = \begin{cases} \mathbb{N}, & \text{if } (f(9^8), f(9^9)) \cap \mathcal{P}_{n^2+1} \neq \emptyset \\ \{0, \dots, 10^6\}, & \text{otherwise} \end{cases}$ |
| \neg (Condition 1) \wedge (Condition 5) | $\mathcal{X} = \mathcal{P}_{n^2+1}$ | Open Problem 3 |
| \neg (Condition 1) \wedge \neg (Condition 5) | $\mathcal{X} = \{0, \dots, f(7)\} \cup \mathcal{P}_{n^2+1}$ | $\mathcal{X} = \begin{cases} \mathbb{N}, & \text{if } (f(9^8), f(9^9)) \cap \mathcal{P}_{n^2+1} \neq \emptyset \\ \{0, \dots, 10^6\} \cup \{n \in \mathbb{N} : n \text{ is the sixth prime number of the form } 2^{2^k} + 1\}, & \text{otherwise} \end{cases}$ |

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