On sets $\mathcal{W} \subseteq \mathbb{N}$ whose infinity follows from the existence in $\mathcal{W}$ of an element which is greater than a threshold number computed for $\mathcal{W}$

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Abstract

We define computable functions $f, g\colon \mathbb{N} \setminus \{0\} \to \mathbb{N} \setminus \{0\}$. For a positive integer n, let Θ_n denote the following statement: if a system $\mathcal{S} \subseteq \Big\{x_i! = x_k : i, k \in \{1, \ldots, n\}\Big\} \cup \Big\{x_i \cdot x_j = x_k : i, j, k \in \{1, \ldots, n\}\Big\}$ has only finitely many solutions in integers $x_1, \ldots, x_n$ greater than 1, then each such solution $(x_1, \ldots, x_n)$ satisfies $\min(x_1, \ldots, x_n) \leqslant f(n)$. The statement Θ_9 proves that if there exists an integer $x > f(9)$ such that $x^2 + 1$ (alternatively, $x! + 1$) is prime, then there are infinitely many primes of the form $n^2 + 1$ (respectively, $n! + 1$). The statement Θ_{16} proves that if there exists a twin prime greater than $f(16) + 3$, then there are infinitely many twin primes. We formulate a statement which proves that if $2^{2^n} + 1$ is composite for some integer $n > g(13)$, then $2^{2^n} + 1$ is composite for infinitely many positive integers n.

Key words and phrases: Brocard's problem, Brocard-Ramanujan equation, composite Fermat numbers, composite numbers of the form $2^{2^n} + 1$, prime numbers of the form $n^2 + 1$, prime numbers of the form $n! + 1$, Richert's lemma, Richert's theorem, twin prime conjecture.

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1 Introduction

The following observation concerns the theme described in the title of the article.

Observation 1. *If $n \in \mathbb{N}$ and $\mathcal{W} \subseteq \{0, \ldots, n\}$, then we take any integer $m \geqslant n$ as a threshold number for $\mathcal{W}$. If $\mathcal{W} \subseteq \mathbb{N}$ and W is infinite, then we take any non-negative integer m as a threshold number for $\mathcal{W}$.*

We define the set $\mathcal{U} \subseteq \mathbb{N}$ by declaring that a non-negative integer n belongs to $\mathcal{U}$ if and only if $\sin\left(10^{10^{10^{10}}}\right) > 0$. This inequality is practically undecidable, see [7].

Corollary 1. *The set $\mathcal{U}$ equals $\emptyset$ or $\mathbb{N}$. The statement "$\mathcal{U} = \emptyset$" remains unproven and the statement "$\mathcal{U} = \mathbb{N}$" remains unproven. Every non-negative integer m is a threshold number for $\mathcal{U}$. For every non-negative integer k, the sentence "$k \in \mathcal{U}$" is only theoretically decidable.*

The first-order language of graph theory contains two relation symbols of arity 2: $\sim$ and $=$, respectively for adjacency and equality of vertices. The term first-order imposes the condition that the variables represent vertices and hence the quantifiers apply to vertices only. For a first-order sentence Λ about graphs, let $\text{Spectrum}(\Lambda)$ denote the set of all positive integers n such that there is a graph on n vertices satisfying Λ. By a graph on n vertices we understand a set of n elements with a binary relation which is symmetric and irreflexive.

Theorem 1. *([15, p. 171]). If a sentence Λ in the language of graph theory has the form $\exists x_1 \ldots x_k \, \forall y_1 \ldots y_l \, \Upsilon(x_1, \ldots, x_k, y_1, \ldots, y_l)$, where $\Upsilon(x_1, \ldots, x_k, y_1, \ldots, y_l)$ is quantifier-free, then either* $\text{Spectrum}(\Lambda) \subseteq [1, (2^k \cdot 4^l) - 1]$ *or* $\text{Spectrum}(\Lambda) \supseteq [k + l, \infty) \cap \mathbb{N}$.

Corollary 2. *The number $\left(2^k \cdot 4^l\right) - 1$ is a threshold number for* $\text{Spectrum}(\Lambda)$.

The classes of the infinite recursively enumerable sets and of the infinite recursive sets are not recursively enumerable, see [13, p. 234].

Corollary 3. *If an algorithm $\texttt{Al}_1$ for every recursive set $\mathcal{W} \subseteq \mathbb{N}$ finds a non-negative integer $\texttt{Al}_1(\mathcal{W})$, then there exists a finite set $\mathcal{M} \subseteq \mathbb{N}$ such that $\mathcal{M} \cap [\texttt{Al}_1(\mathcal{M}) + 1, \infty) \neq \emptyset$.*

Corollary 4. *If an algorithm $\texttt{Al}_2$ for every recursively enumerable set $\mathcal{W} \subseteq \mathbb{N}$ finds a non-negative integer $\texttt{Al}_2(\mathcal{W})$, then there exists a finite set $\mathcal{M} \subseteq \mathbb{N}$ such that $\mathcal{M} \cap [\texttt{Al}_2(\mathcal{M}) + 1, \infty) \neq \emptyset$.*

Let $K = \{j \in \mathbb{N} : \, 2^{\aleph_j} = \aleph_{j+1}\}$.

Theorem 2. *If ZFC is consistent, then for every non-negative integer n the sentence*

$$\text{"}n \texttt{ is a threshold number for } K\text{"}$$

is not provable in ZFC

Proof. There exists a model $\mathcal{E}$ of ZFC such that

$$\forall i \in \{0, \dots, n+1\} \; \mathcal{E} \models 2^{\aleph_i} = \aleph_{i+1}$$

and

$$\forall i \in \{n+2, n+3, n+4, \dots\} \; \mathcal{E} \models 2^{\aleph_i} = \aleph_{i+2}$$

see [5] and [8, p. 232]. In the model $\mathcal{E}$, $K = \{0, \dots, n+1\}$ and n is not a threshold number for K. □

Theorem 3. *If ZFC is consistent, then for every non-negative integer n the sentence*

"n is not a threshold number for K"

is not provable in ZFC.

Proof. The Generalized Continuum Hypothesis (GCH) is consistent with ZFC, see [8, p. 188] and [8, p. 190]. GCH implies that $K = \mathbb{N}$. Consequently, GCH implies that every non-negative integer n is a threshold number for K. □

Theorem 4. *([2, p. 35]). There exists a polynomial $D(x_1, \dots, x_m)$ with integer coefficients such that if ZFC is arithmetically consistent, then the sentences*

"The equation $D(x_1, \dots, x_m) = 0$ is solvable in non-negative integers"

and

"The equation $D(x_1, \dots, x_m) = 0$ is not solvable in non-negative integers"

are not provable in ZFC.

Let Δ denote the set of all non-negative integers k such that the equation $D(x_1, \dots, x_m) = 0$ has no solutions in $\{0, \dots, k\}^m$. Since the set $\{0, \dots, k\}^m$ is finite, the set Δ is computable. Theorem 4 implies the following corollary.

Corollary 5. *If ZFC is arithmetically consistent, then for every non-negative integer n the sentences*

"n is a threshold number for Δ"

and

"n is not a threshold number for Δ"

are not provable in ZFC.

Let $g(1) = 1$, and let $g(n+1) = 2^{2^{g(n)}}$ for every positive integer n.

Hypothesis 1. *([20]). If a system*

$$\mathcal{S} \subseteq \{x_i \cdot x_j = x_k : \ i,j,k \in \{1,\ldots,n\}\} \cup \{x_i + 1 = x_k : \ i,k \in \{1,\ldots,n\}\}$$

has only finitely many solutions in non-negative integers $x_1,\ldots,x_n$, then each such solution $(x_1,\ldots,x_n)$ satisfies $x_1,\ldots,x_n \leqslant g(2n)$.

Theorem 5. *([20]). Hypothesis 1 implies that for every $W(x_1,\ldots,x_n) \in \mathbb{Z}[x_1,\ldots,x_n]$ we can compute a threshold number $b \in \mathbb{N} \setminus \{0\}$ such that any non-negative integers $a_1,\ldots,a_n$ which satisfy*

$$(W(a_1,\ldots,a_n) = 0) \wedge (\max(a_1,\ldots,a_n) > b)$$

guarantee that the equation $W(x_1,\ldots,x_n) = 0$ has infinitely many solutions in non-negative integers.

2 Basic lemmas

Let $f(1) = 2$, $f(2) = 4$, and let $f(n+1) = f(n)!$ for every integer $n \geqslant 2$. Let $\mathcal{V}_1$ denote the system of equations $\{x_1! = x_1\}$, and let $\mathcal{V}_2$ denote the system of equations $\{x_1! = x_1, \ x_1 \cdot x_1 = x_2\}$. For an integer $n \geqslant 3$, let $\mathcal{V}_n$ denote the following system of equations:

$$\begin{cases} x_1! &= x_1 \\ x_1 \cdot x_1 &= x_2 \\ \forall i \in \{2,\ldots,n-1\} \ x_i! &= x_{i+1} \end{cases}$$

The diagram in Figure 1 illustrates the construction of the system $\mathcal{V}_n$.

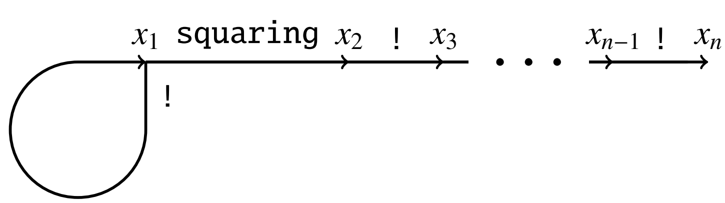


Fig. 1 Construction of the system $\mathcal{V}_n$

Lemma 1. *For every positive integer n, the system $\mathcal{V}_n$ has exactly one solution in integers greater than 1, namely $\big(f(1),\ldots,f(n)\big)$.*

Let

$$H_n = \left\{x_i! = x_k : \ i,k \in \{1,\ldots,n\}\right\} \cup \left\{x_i \cdot x_j = x_k : \ i,j,k \in \{1,\ldots,n\}\right\}$$

For a positive integer n, let Θ_n denote the following statement: if a system $\mathcal{S} \subseteq H_n$ has at most finitely many solutions in integers $x_1,\ldots,x_n$ greater than 1, then each such solution $(x_1,\ldots,x_n)$ satisfies $\min(x_1,\ldots,x_n) \leqslant f(n)$. The assumption $\min(x_1,\ldots,x_n) \leqslant f(n)$ is weaker than the assumption $\max(x_1,\ldots,x_n) \leqslant f(n)$ suggested by Lemma 1.

Lemma 2. *For every positive integer n, the system H_n has a finite number of subsystems.*

Theorem 6. *Every statement Θ_n is true with an unknown integer bound that depends on n.*

Proof. It follows from Lemma 2. □

Lemma 3. *For every integers x and y greater than* 1, $x! \cdot y = y!$ *if and only if* $x + 1 = y$.

Lemma 4. *If* $x \geqslant 4$, *then* $\dfrac{(x-1)!+1}{x} > 1$.

Lemma 5. *(Wilson's theorem, [6, p. 89]). For every integer $x \geqslant 2$, x is prime if and only if x divides* $(x-1)!+1$.

3 Brocard's problem

A weak form of Szpiro's conjecture implies that there are only finitely many solutions to the Brocard-Ramanujan equation $x! + 1 = y^2$, see [14]. It is conjectured that $x! + 1$ is a square only for $x \in \{4,5,7\}$, see [21, p. 297].

Let $\mathcal{A}$ denote the following system of equations:

$$\begin{cases} x_1! &= x_2 \\ x_2! &= x_3 \\ x_5! &= x_6 \\ x_4 \cdot x_4 &= x_5 \\ x_3 \cdot x_5 &= x_6 \end{cases}$$

Lemma 3 and the diagram in Figure 2 explain the construction of the system $\mathcal{A}$.

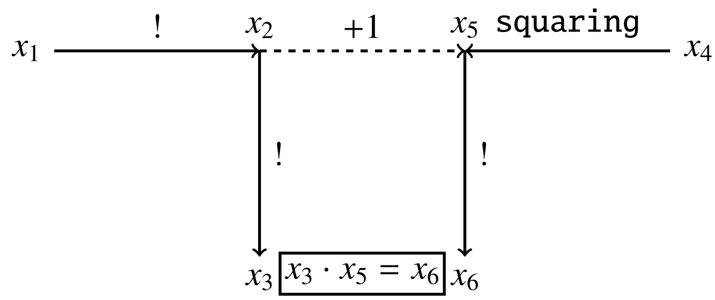


Fig. 2 Construction of the system $\mathcal{A}$

Lemma 6. *For every integers x_1 and x_4 greater than 1, the system $\mathcal{A}$ is solvable in integers x_2, x_3, x_5, x_6 greater than 1 if and only if $x_1! + 1 = x_4^2$. In this case, the integers x_2, x_3, x_5, x_6 are uniquely determined by the following equalities:*

$$\begin{aligned} x_2 &= x_1! \\ x_3 &= (x_1!)! \\ x_5 &= x_1! + 1 \\ x_6 &= (x_1! + 1)! \end{aligned}$$

and $x_1 = \min(x_1, \ldots, x_6)$.

Proof. It follows from Lemma 3. □

Theorem 7. *The statement Θ_6 proves the following implication: if the equation $x_1! + 1 = x_4^2$ has only finitely many solutions in positive integers, then each such solution (x_1, x_4) satisfies $x_1 \leqslant f(6)$.*

Proof. Let positive integers x_1 and x_4 satisfy $x_1!+1 = x_4^2$. Then, $x_1, x_4 \in \mathbb{N} \setminus \{0, 1\}$. By Lemma 6, there exists a unique tuple $(x_2, x_3, x_5, x_6) \in (\mathbb{N} \setminus \{0, 1\})^4$ such that the tuple $(x_1, \ldots, x_6)$ solves the system $\mathcal{A}$. Lemma 6 guarantees that $x_1 = \min(x_1, \ldots, x_6)$. By the antecedent and Lemma 6, the system $\mathcal{A}$ has only finitely many solutions in integers $x_1, \ldots, x_6$ greater than 1. Therefore, the statement Θ_6 implies that $x_1 = \min(x_1, \ldots, x_6) \leqslant f(6)$. □

Hypothesis 2. *The implication in Theorem 7 is true.*

Corollary 6. *Assuming Hypothesis 2, a single query to an oracle for the halting problem decides the problem of the infinitude of the solutions of the equation $x! + 1 = y^2$.*

4 Are there infinitely many prime numbers of the form $n^2 + 1$?

Edmund Landau's conjecture states that there are infinitely many primes of the form $n^2 + 1$, see [12, pp. 37–38]. Let $\mathcal{B}$ denote the following system of equations:

$$\begin{cases} x_2! &= x_3 \\ x_3! &= x_4 \\ x_5! &= x_6 \\ x_8! &= x_9 \\ x_1 \cdot x_1 &= x_2 \\ x_3 \cdot x_5 &= x_6 \\ x_4 \cdot x_8 &= x_9 \\ x_5 \cdot x_7 &= x_8 \end{cases}$$

Lemma 3 and the diagram in Figure 3 explain the construction of the system $\mathcal{B}$.

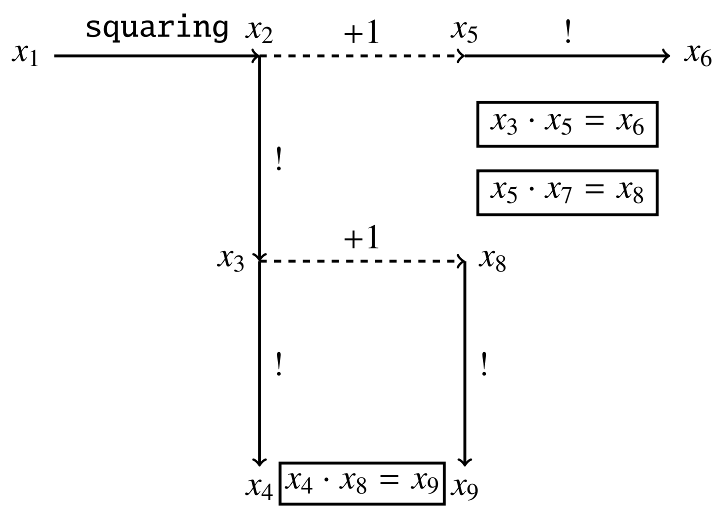


Fig. 3 Construction of the system $\mathcal{B}$

Lemma 7. *For every integer $x_1 \geqslant 2$, the system $\mathcal{B}$ is solvable in integers $x_2, \dots, x_9$ greater than 1 if and only if $x_1^2 + 1$ is prime. In this case, the integers $x_2, \dots, x_9$ are uniquely determined* *by the following equalities:*

$$
\begin{array}{rcl}
x_2 & = & x_1^2 \\
x_3 & = & (x_1^2)! \\
x_4 & = & ((x_1^2)!)! \\
x_5 & = & x_1^2 + 1 \\
x_6 & = & (x_1^2 + 1)! \\
x_7 & = & \dfrac{(x_1^2)! + 1}{x_1^2 + 1} \\
x_8 & = & (x_1^2)! + 1 \\
x_9 & = & ((x_1^2)! + 1)!
\end{array}
$$

and $\min(x_1, \ldots, x_9) = x_1$.

Proof. By Lemmas 3 and 4, for every integer $x_1 \geqslant 2$, the system $\mathcal{B}$ is solvable in integers $x_2, \ldots, x_9$ greater than 1 if and only if $x_1^2 + 1$ divides $(x_1^2)! + 1$. Hence, the claim of Lemma 7 follows from Lemma 5. □

Theorem 8. *The statement* Θ_9 *proves the following implication: if there exists an integer* $x_1 > f(9)$ *such that* $x_1^2 + 1$ *is prime, then there are infinitely many primes of the form* $n^2 + 1$.

Proof. Assume that an integer x_1 is greater than $f(9)$ and $x_1^2 + 1$ is prime. By Lemma 7, there exists a unique tuple $(x_2, \ldots, x_9) \in (\mathbb{N} \setminus \{0, 1\})^8$ such that the tuple $(x_1, x_2, \ldots, x_9)$ solves the system $\mathcal{B}$. Lemma 7 guarantees that $\min(x_1, \ldots, x_9) = x_1$. Since $\mathcal{B} \subseteq H_9$, the statement Θ_9 and the inequality $\min(x_1, \ldots, x_9) = x_1 > f(9)$ imply that the system $\mathcal{B}$ has infinitely many solutions $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0, 1\})^9$. According to Lemma 7, there are infinitely many primes of the form $n^2 + 1$. □

Hypothesis 3. *The implication in Theorem 8 is true.*

Corollary 7. *Assuming Hypothesis 3, a single query to an oracle for the halting problem decides the problem of the infinitude of primes of the form* $n^2 + 1$.

Let $\mathcal{P}$ denote the set of prime numbers. For a non-negative integer n, let $\Omega(n)$ denote the following statement: $\exists m \in \mathbb{N} \cap (n, \infty)\ m^2 + 1 \in \mathcal{P}$. By Theorem 8, assuming the statement Θ_9, we can infer the statement $\forall n \in \mathbb{N}\ \Omega(n)$ from any statement $\Omega(n)$ with $n \geqslant f(9)$. A similar situation holds for inference by the so called *"super-induction method"*, see [22]–[25]. In section 8, we present a theorem whose computer-assisted proof is based on the super-induction method.

5 Are there infinitely many prime numbers of the form $n! + 1$?

It is conjectured that there are infinitely many primes of the form $n! + 1$, see [1, p. 443] and [18]. Let $\mathcal{G}$ denote the following system of equations:

$$\begin{cases} x_1! = x_2 \\ x_2! = x_3 \\ x_3! = x_4 \\ x_5! = x_6 \\ x_8! = x_9 \\ x_3 \cdot x_5 = x_6 \\ x_4 \cdot x_8 = x_9 \\ x_5 \cdot x_7 = x_8 \end{cases}$$

Lemma 3 and the diagram in Figure 4 explain the construction of the system $\mathcal{G}$.

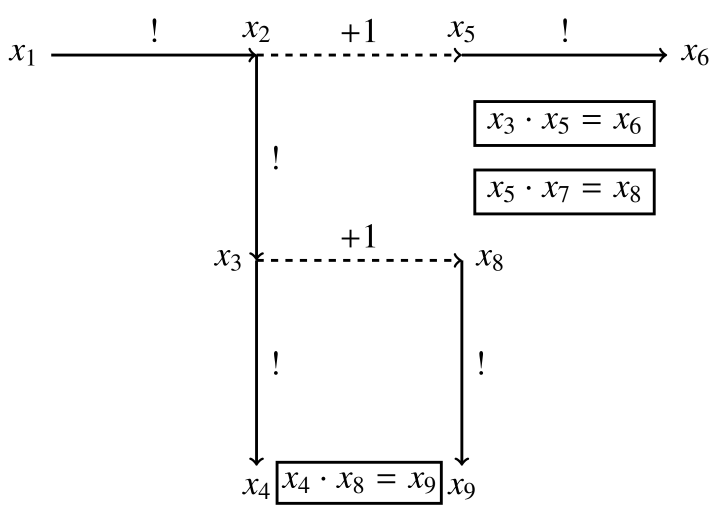


Fig. 4 Construction of the system $\mathcal{G}$

Lemma 8. *For every integer $x_1 \geqslant 2$, the system $\mathcal{G}$ is solvable in integers $x_2, \ldots, x_9$ greater than 1 if and only if $x_1! + 1$ is prime. In this case, the integers $x_2, \ldots, x_9$ are uniquely determined by* *the following equalities:*

$$
\begin{array}{rcl}
x_2 & = & x_1! \\
x_3 & = & (x_1!)! \\
x_4 & = & ((x_1!)!)! \\
x_5 & = & x_1^{!} + 1 \\
x_6 & = & (x_1! + 1)! \\
x_7 & = & \dfrac{(x_1!)! + 1}{x_1! + 1} \\
x_8 & = & (x_1!)! + 1 \\
x_9 & = & ((x_1!)! + 1)!
\end{array}
$$

and $\min(x_1, \ldots, x_9) = x_1$.

Proof. By Lemmas 3 and 4, for every integer $x_1 \geqslant 2$, the system $\mathcal{G}$ is solvable in integers $x_2, \ldots, x_9$ greater than 1 if and only if $x_1! + 1$ divides $(x_1!)! + 1$. Hence, the claim of Lemma 8 follows from Lemma 5. □

Theorem 9. *The statement* Θ_9 *proves the following implication: if there exists an integer* $x_1 > f(9)$ *such that* $x_1! + 1$ *is prime, then there are infinitely many primes of the form* $n! + 1$.

Proof. Assume that an integer x_1 is greater than $f(9)$ and $x_1! + 1$ is prime. By Lemma 8, there exists a unique tuple $(x_2, \ldots, x_9) \in (\mathbb{N} \setminus \{0, 1\})^8$ such that the tuple $(x_1, x_2, \ldots, x_9)$ solves the system $\mathcal{G}$. Lemma 8 guarantees that $\min(x_1, \ldots, x_9) = x_1$. Since $\mathcal{G} \subseteq H_9$, the statement Θ_9 and the inequality $\min(x_1, \ldots, x_9) = x_1 > f(9)$ imply that the system $\mathcal{G}$ has infinitely many solutions $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0, 1\})^9$. According to Lemma 8, there are infinitely many primes of the form $n! + 1$. □

Hypothesis 4. *The implication in Theorem 9 is true.*

Corollary 8. *Assuming Hypothesis 4, a single query to an oracle for the halting problem decides the problem of the infinitude of primes of the form* $n! + 1$.

6 The twin prime conjecture

A twin prime is a prime number that is either 2 less or 2 more than another prime number. The twin prime conjecture states that there are infinitely many twin primes, see [12, p. 39].

Let C denote the following system of equations:

$$\begin{cases} x_1! &= x_2 \\ x_2! &= x_3 \\ x_4! &= x_5 \\ x_6! &= x_7 \\ x_7! &= x_8 \\ x_9! &= x_{10} \\ x_{12}! &= x_{13} \\ x_{15}! &= x_{16} \\ x_2 \cdot x_4 &= x_5 \\ x_5 \cdot x_6 &= x_7 \\ x_7 \cdot x_9 &= x_{10} \\ x_4 \cdot x_{11} &= x_{12} \\ x_3 \cdot x_{12} &= x_{13} \\ x_9 \cdot x_{14} &= x_{15} \\ x_8 \cdot x_{15} &= x_{16} \end{cases}$$

Lemma 3 and the diagram in Figure 5 explain the construction of the system C.

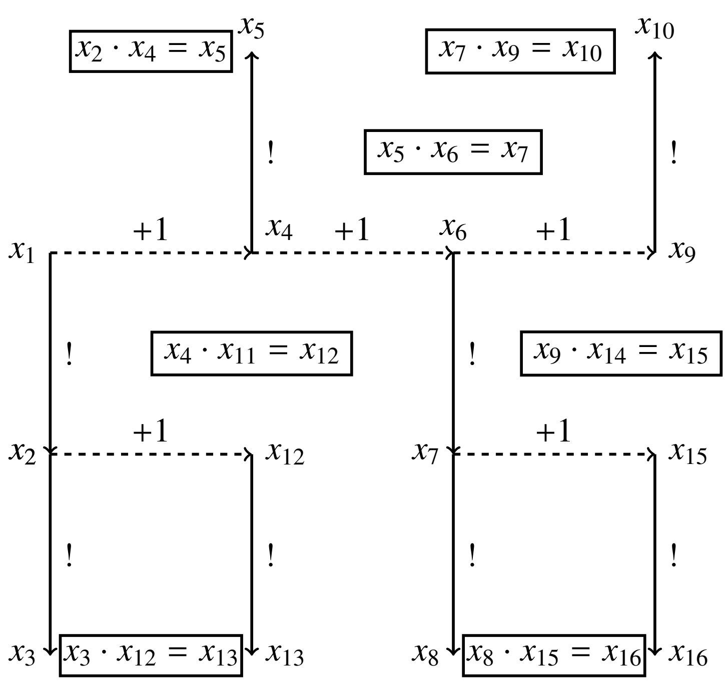


Fig. 5 Construction of the system C

Lemma 9. *If $x_4 = 2$, then the system C has no solutions in integers $x_1, \ldots, x_{16}$ greater than* 1.

Proof. The equality $x_2 \cdot x_4 = x_5 = x_4!$ and the equality $x_4 = 2$ imply that $x_2 = 1$. □

Lemma 10. *If $x_4 = 3$, then the system C has no solutions in integers $x_1, \ldots, x_{16}$ greater than* 1.

Proof. The equality $x_4 \cdot x_{11} = x_{12} = (x_4 - 1)! + 1$ and the equality $x_4 = 3$ imply that $x_{11} = 1$. □

Lemma 11. *For every $x_4 \in \mathbb{N} \setminus \{0, 1, 2, 3\}$ and for every $x_9 \in \mathbb{N} \setminus \{0, 1\}$, the system C is solvable in integers x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16} greater than* 1 *if and only if x_4 and x_9 are prime and $x_4 + 2 = x_9$. In this case, the integers x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16} are uniquely determined by the following equalities:*

$$\begin{aligned}
x_1 &= x_4 - 1 \\
x_2 &= (x_4 - 1)! \\
x_3 &= ((x_4 - 1)!)! \\
x_5 &= x_4! \\
x_6 &= x_9 - 1 \\
x_7 &= (x_9 - 1)! \\
x_8 &= ((x_9 - 1)!)! \\
x_{10} &= x_9! \\
x_{11} &= \frac{(x_4 - 1)! + 1}{x_4} \\
x_{12} &= (x_4 - 1)! + 1 \\
x_{13} &= ((x_4 - 1)! + 1)! \\
x_{14} &= \frac{(x_9 - 1)! + 1}{x_9} \\
x_{15} &= (x_9 - 1)! + 1 \\
x_{16} &= ((x_9 - 1)! + 1)!
\end{aligned}$$

and $\min(x_1, \ldots, x_{16}) = x_1 = x_9 - 3$.

Proof. By Lemmas 3 and 4, for every $x_4 \in \mathbb{N} \setminus \{0, 1, 2, 3\}$ and for every $x_9 \in \mathbb{N} \setminus \{0, 1\}$, the system C is solvable in integers x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16} greater than 1 if and only if

$$\Big(x_4 + 2 = x_9\Big) \wedge \Big(x_4 | (x_4 - 1)! + 1\Big) \wedge \Big(x_9 | (x_9 - 1)! + 1\Big)$$

Hence, the claim of Lemma 11 follows from Lemma 5. □

Theorem 10. *The statement Θ_{16} proves the following implication: if there exists a twin prime greater than $f(16)+3$, then there are infinitely many twin primes.*

Proof. Assume that the antecedent holds. Then, there exist prime numbers x_4 and x_9 such that $x_9 = x_4 + 2 > f(16) + 3$. Hence, $x_4 \in \mathbb{N} \setminus \{0,1,2,3\}$. By Lemma 11, there exists a unique tuple $(x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}) \in (\mathbb{N} \setminus \{0,1\})^{14}$ such that the tuple $(x_1, \ldots, x_{16})$ solves the system C. Lemma 11 guarantees that $\min(x_1, \ldots, x_{16}) = x_1 = x_9 - 3 > f(16)$. Since $C \subseteq H_{16}$, the statement Θ_{16} and the inequality $\min(x_1, \ldots, x_{16}) > f(16)$ imply that the system C has infinitely many solutions in integers $x_1, \ldots, x_{16}$ greater than 1. According to Lemmas 9–11, there are infinitely many twin primes. □

Hypothesis 5. *The implication in Theorem 10 is true.*

Corollary 9. *(cf. [3]). Assuming Hypothesis 5, a single query to an oracle for the halting problem decides the twin prime problem.*

7 Are there infinitely many composite Fermat numbers?

Primes of the form $2^{2^n}+1$ are called Fermat primes, as Fermat conjectured that every integer of the form $2^{2^n}+1$ is prime, see [11, p. 1]. Fermat correctly remarked that $2^{2^0}+1=3$, $2^{2^1}+1=5$, $2^{2^2}+1=17$, $2^{2^3}+1=257$, and $2^{2^4}+1=65537$ are all prime, see [11, p. 1].

Open Problem. ([11, p. 159]). *Are there infinitely many composite numbers of the form $2^{2^n}+1$?*

Most mathematicians believe that $2^{2^n}+1$ is composite for every integer $n \geqslant 5$, see [10, p. 23].

Theorem 11. *([19]). An unproven inequality stated in [19] implies that $2^{2^n}+1$ is composite for every integer $n \geqslant 5$.*

Lemma 12. *([11, p. 38]). For every positive integer n, if a prime number p divides $2^{2^n}+1$, then there exists a positive integer k such that $p = k \cdot 2^{n+1}+1$.*

Corollary 10. *Since $k \cdot 2^{n+1}+1 \geqslant 2^{n+1}+1 \geqslant n+3$, for every positive integers x, y, and n, the equality $(x+1)(y+1) = 2^{2^n}+1$ implies that $\min(n, x, x+1, y, y+1) = n$.*

Let

$$G_n = \left\{x_i \cdot x_j = x_k :\ i,j,k \in \{1,\ldots,n\}\right\} \cup \left\{2^{2^{x_i}} = x_k :\ i,k \in \{1,\ldots,n\}\right\}$$

Lemma 13. *The following subsystem of* G_n

$$\begin{cases} x_1 \cdot x_1 = x_1 \\ \forall i \in \{1, \dots, n-1\}\ 2^{2^{x_i}} = x_{i+1} \end{cases}$$

has exactly one solution $(x_1, \dots, x_n) \in (\mathbb{N} \setminus \{0\})^n$*, namely* $(g(1), \dots, g(n))$.

For a positive integer n, let Ψ_n denote the following statement: if a system $S \subseteq G_n$ has at most finitely many solutions in positive integers $x_1, \dots, x_n$, then each such solution $(x_1, \dots, x_n)$ satisfies $\min(x_1, \dots, x_n) \leqslant g(n)$. The assumption $\min(x_1, \dots, x_n) \leqslant g(n)$ is weaker than the assumption $\max(x_1, \dots, x_n) \leqslant g(n)$ suggested by Lemma 13.

Lemma 14. *For every positive integer n, the system* G_n *has a finite number of subsystems.*

Theorem 12. *Every statement* Ψ_n *is true with an unknown integer bound that depends on n.*

Proof. It follows from Lemma 14. □

Lemma 15. *For every non-negative integers b and c,* $b + 1 = c$ *if and only if* $2^{2^b} \cdot 2^{2^b} = 2^{2^c}$.

Theorem 13. *The statement* Ψ_{13} *proves the following implication: if* $2^{2^n} + 1$ *is composite for some integer* $n > g(13)$*, then* $2^{2^n} + 1$ *is composite for infinitely many positive integers n.*

Proof. Let us consider the equation

$$(x + 1)(y + 1) = 2^{2^z} + 1 \tag{1}$$

in positive integers. By Lemma 15, we can transform equation (1) into an equivalent system $\mathcal{F}$ which has 13 variables (x, y, z, and 10 other variables) and which consists of equations of the forms $\alpha \cdot \beta = \gamma$ and $2^{2^\alpha} = \gamma$, see the diagram in Figure 6.

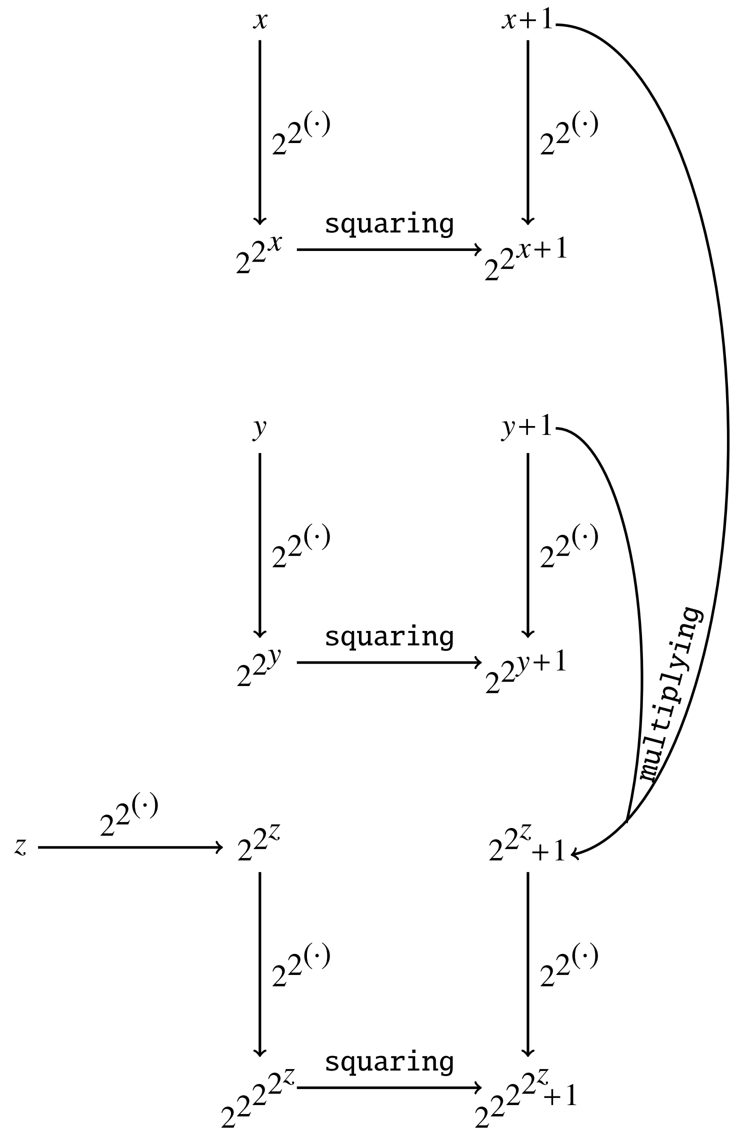


Fig. 6 Construction of the system $\mathcal{F}$

Assume that $2^{2^n}+1$ is composite for some integer $n > g(13)$. By this and Corollary 10, equation (1) has a solution $(x,y,z) \in (\mathbb{N} \setminus \{0\})^3$ such that $z = n$ and $z = \min(z,x,x+1,y,y+1)$. Hence, the system $\mathcal{F}$ has a solution in positive integers such that $z = n$ and n is the smallest number in the solution sequence. Since $n > g(13)$, the statement Ψ_{13} implies that the system $\mathcal{F}$ has infinitely many solutions in positive integers. Therefore, there are infinitely many positive integers n such that $2^{2^n}+1$ is composite. □

Hypothesis 6. *The implication in Theorem 13 is true.*

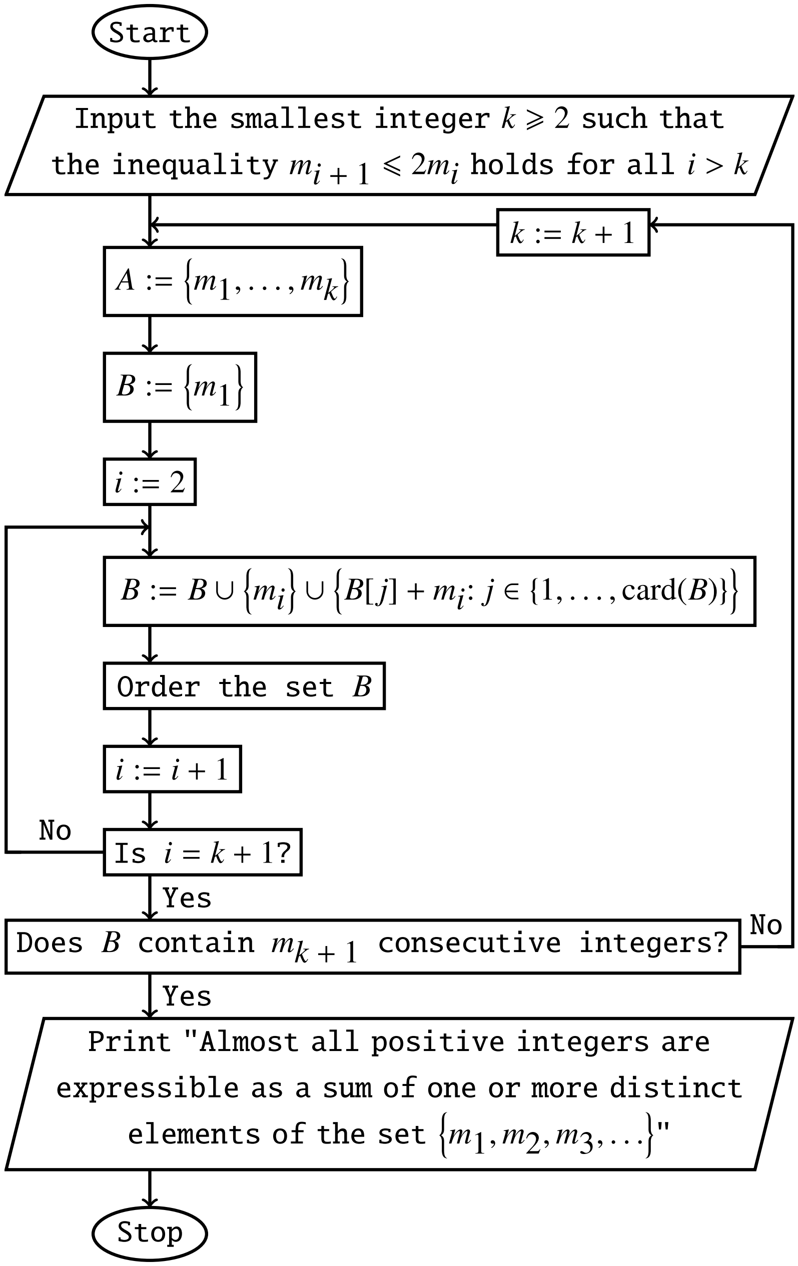
Corollary 11. *Assuming Hypothesis 6, a single query to an oracle for the halting problem decides whether or not the set of composite Fermat numbers is infinite.*

8 A computer-assisted proof that a certain subset of $\mathbb{N} \setminus \{0\}$ is cofinite

The following lemma is known as Richert's lemma or Richert's theorem.

Lemma 16. *([4], [16], [17, p. 152]). Let $\{m_i\}_{i=1}^{\infty}$ be an increasing sequence of positive integers such that for some positive integer k the inequality $m_{i+1} \leqslant 2m_i$ holds for all $i > k$. Suppose there exists a non-negative integer b such that the numbers $b+1,\ b+2,\ b+3,\ \ldots,\ b+m_{k+1}$ are all expressible as sums of one or more distinct elements of the set $\{m_1, \ldots, m_k\}$. Then every integer greater than b is expressible as a sum of one or more distinct elements of the set $\{m_1, m_2, m_3, \ldots\}$.*

Corollary 12. *If the sequence $\{m_i\}_{i=1}^{\infty}$ is computable and the algorithm in Flowchart 1 terminates, then almost all positive integers are expressible as a sum of one or more distinct elements of the set $\{m_1, m_2, m_3, \ldots\}$.*



Flowchart 1

For a large class of sequences $\{m_i\}_{i=1}^{\infty}$ the converse is true: if almost all positive integers are expressible as a sum of one or more distinct elements of the set $\{m_1, m_2, m_3, \ldots\}$, then the algorithm in Flowchart 1 terminates, see [9, Theorem 2.3].

Let $[\cdot]$ denote the integer part function. For a positive integer i, let $t_i = \frac{(i+19)^{i+19}}{(i+19)! \cdot 2^{i+19}}$, and let $m_i = [t_i]$. Let $\mathcal{T}$ denote the set of all positive integers which are expressible as a sum of one or more distinct elements of the set $\{m_i : i \in \mathbb{N} \setminus \{0\}\}$.

Lemma 17. *The inequality $m_{i+1} \leqslant 2m_i$ holds for every positive integer i.*

Proof. For every positive integer i,

$$\frac{m_i}{m_{i+1}} = \frac{[t_i]}{[t_{i+1}]} > \frac{t_i - 1}{t_{i+1}} = \left(\frac{t_i}{t_{i+1}}\right) - \left(\frac{1}{t_{i+1}}\right) \geqslant \left(\frac{t_i}{t_{i+1}}\right) - \left(\frac{1}{t_2}\right) =$$

$$2 \cdot \frac{i+20}{i+19} \cdot \left(1 - \frac{1}{i+20}\right)^{i+20} - \left(\frac{21! \cdot 2^{21}}{21^{21}}\right) > 2 \cdot \left(1 - \frac{1}{21}\right)^{21} - \left(\frac{21! \cdot 2^{21}}{21^{21}}\right) =$$

$$\frac{408715852844271520448512000}{5842587018385982521381124421}$$

The above fraction was computed by *MuPAD* and is greater than $\frac{1}{2}$. □

Theorem 14. *The algorithm in Flowchart 1 terminates for the sequence $\{m_i\}_{i=1}^{\infty}$.* 2761 *is the largest integer which does not belong to $\mathcal{T}$.*

Proof. By Lemma 17, we take $k = 2$ as the initial value of k. The first 15 lines of the *MuPAD* code below describe the algorithm in Flowchart 1. It follows from the fact that *MuPAD* automatically orders every finite set of integers.

***MuPAD* code**

```
TEXTWIDTH:=80:
k:=2:
repeat
C:={floor((i+19)^(i+19)/((i+19)!*2^(i+19))) $i=1..k+1}:
A:={floor((i+19)^(i+19)/((i+19)!*2^(i+19))) $i=1..k}:
B:={A[1]}:
for i from 2 to nops(A) do
B:=B union {A[i]} union {B[j]+A[i] $j=1..nops(B)}:
end_for:
G:={y $y=B[1]-1..B[nops(B)]+1} minus B:
H:={G[n+1]-G[n] $n=1..nops(G)-1}:
k:=k+1:
until H[nops(H)]>C[nops(C)] end_repeat:
print(Unquoted, "Almost all positive integers are expressible as a sum of"):
```

```
print(Unquoted, "one or more distinct elements of the set {m_1,m_2,m_3,...}"):
print(k-1):
print({floor((i+19)^(i+19)/((i+19)!*2^(i+19))) $i=1..k}):
print({2761} minus B):
print({2761+i $i=1..C[k]} minus B):
```

Output

```
Almost all positive integers are expressible as a sum of

one or more distinct elements of the set {m_1,m_2,m_3,...}

15

{41, 54, 72, 96, 128, 170, 227, 303, 404, 540, 722, 966, 1293, 1730, 2317, 3105}

{2761}

{}
```

The output indicates that the algorithm in Flowchart 1 terminates and 15 is the smallest value of k which satisfies the condition of Lemma 16. The last three instructions list the following sets:

$\{m_i : i \in \{1, \dots, 16\}\}$,
$\{2761\} \setminus B$,
$\{2761 + 1,\ \dots,\ 2761 + m_{16}\} \setminus B$,

where B denotes the set of all positive integers which are expressible as a sum of one or more distinct elements of the set $\{m_i : i \in \{1, \dots, 15\}\}$. Since the set $\{2761\}$ equals $\{2761\} \setminus B$, we conclude that $2761 \notin B$. By this and the inequality $m_{16} = 3105 > 2761$, we conclude that $2761 \notin \mathcal{T}$. Since the empty set $\{\}$ equals

$$\{2761 + 1,\ \dots,\ 2761 + m_{16}\} \setminus B$$

we conclude that each of the integers

$$2761 + 1,\ 2761 + 2,\ 2761 + 3,\ \dots,\ 2761 + m_{16}$$

belongs to B. Consequently, in virtue of Lemma 17, we can apply Lemma 16 with $k = 15$ and $b = 2761$ to confirm that every integer greater than 2761 belongs to $\mathcal{T}$. □

MuPAD is a general-purpose computer algebra system. The commercial version of *MuPAD* is no longer available as a stand-alone product, but only as the *Symbolic Math Toolbox* of *MATLAB*. Fortunately, the presented code can be executed by *MuPAD Light*, which was offered for free for research and education until autumn 2005.

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